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ABSTRACT and CONTENTS

The comprehensive diagnostic for all non-branching ITP instructions is described. All information necessary for operating the program is (hopefully) given.

ITP DIAGNOSTIC

This program tests all ITP instructions and features except

branches

UPS, UPOT, UPIN, XCA, CCA, EXU

pops

addressing (but the program uses all addressing features except absolute indirection)

traps (illegal opcodes and memory protection)

input/output

It works in the following way.

Instructions are described by a table called INSTAB. Each entry in this tables describes one instruction and contains 8 words:

- 0) the instruction under test. If it addresses memory, its address field contains NOP2.
- 1) the address of the result table for the instruction
- 2) the address of a word containing the value which should be in A after execution of the instruction
- 3) same for B
- 4) same for X
- 5) same for NOP2
- 6) same for skip. The value is zero if the instruction does not skip, -1 if it does.
- 7) the number of words in a result table entry for this instruction.

The contents of A,B,X and NOP2 before instruction execution are contained in SA, SB, SX, and OP2; the addresses of these locations are therefore frequent occupants of words 2-5 of an instruction table entry. For example, the

instruction table entry for STA is

STA NOP2

Ø

SA

SB

SX

SA

=Ø

Ø

indicating that STA has no result table, does not change ABX, never skips and leaves the old contents of A in the location addressed by it.

The result table takes care of the possibility that one of the results of an instruction might not be the same as any of its inputs. In this case, the result table contains an entry for each set of arguments which may be presented to the function. The address of the entry is put into X before the instruction table is accessed, so that an indexed address in words 2-6 will refer to the result table entry. The table is set up by running the diagnostic with the SKE instruction which does the comparison replaced by an STA. Here is the instruction table entry for SKR

SKR NOP2

XXXX

SA

SB

SX

2B7

2B7+1

2

It has two result table words, one for the contents of the effective address and one for skip/noskip. The address of the current result table entry is initialized from word 1 of the instruction table and is incremented by word 7 after each execution of the instruction under test. Its current value is kept in RESPTR.

Arguments are presented to instructions (the same arguments for all instructions) from a table called ARGTAB. The four words of this table addressed by ARGPTR are used to set up A, B, X and NOP2 before the instruction under test is executed. ARGPTR is initialized to ARGTAB when testing of an instruction is started and is incremented by one after each trial. NRES trials are made for each instruction. Note that each word of ARGTAB is used successively as A, B, X and NOP2 (except the first and last 3).

The following flags and pointers control execution of the diagnostic. Numbers in parentheses are the values they are initialized to

BEGINS (INSTAB) address of first INSTAB entry to test
ENDINS (EINST-8) address of last INSTAB entry to test
BEGCHM (\emptyset) address of first memory location included
in checksum
ENDCHM (10000B) address of last memory location
included in checksum
EWAITF (-1) if -1, wait for carriage return to be
typed after an error before proceeding; if \emptyset ,
proceed at once.
DWAITF (-1) if -1, wait for carriage return to be
typed after testing all instructions before looping
to test them again; if \emptyset , proceed at once.

IWAITF (\emptyset) if -1, wait after testing each instruction before looping to test the next one; if \emptyset , proceed at once.

TRPTF (\emptyset) if -1, repeat the current test with the current operands indefinitely; if \emptyset , sequence through tests normally.

IRPTF (\emptyset) if -1, request the test of the current instruction indefinitely; if \emptyset , sequence through instructions normally.

CHKSF (-1) if -1, memory is checksummed after each instruction has been tested to see if execution of that instruction (or foulups in the diagnostic) improperly altered memory.

CPRINT (-1) if -1, prints a message after a compare error; if \emptyset , does not.

SPRINT (-1) if -1, prints a message after a checksum error; if \emptyset , does not.

The following locations contain useful information:

INSPTR address of INSTAB entry for instruction under test

RESPTR address of result table entry for operands under test

ARGPTR address of 4 words used to initialize ABX and NOP2

ERRCNT number of errors detected

LCTR number of complete loops performed

ICTR number of instructions tested in current loop

TCTR number of trials of current instruction

The following locations contain entry points to the various loops:

ESTART start program

EMTEST test all instructions

EMTLOOP test one instruction

EMTSTLP test one case of one instruction

There are two kinds of errors: comparison and checksum. A comparison error occurs when the new contents of the register, and effective address after execution of an instruction are not what the diagnostic thinks they ought to be. What is printed is:

```
COMPARE ON instruction AFTER nØ:n1:n2
old a  old b  old x  old memory
new a  new b  new x  new memory
```

where nØ is the number of complete cycles, n1 is the number of instructions tested in the current cycle, and n2 the number of trials of the current instruction.

A checksum error (see CHKSF above) prints
CHECKSUM ON instruction AFTER n1:n2
only.

After an error a wait occurs if the proper flag is set. Execution then continues by retrying the instruction which failed. In the case of a compare error, a checksum is tried first, which may elicit a checksum error.

The program is loaded at 10000B in the ITP address space. Since self-fill starts the ITP address space at 10000B in memory, the program itself starts at 20000B. The attached table gives the addresses of all the magic cells named above relative to 20000B.

PAGE SYSTEM (LAMPSON)/LISTING OF DITP 12/14/69 0345:33

00000 LIST
00000 =1 940F EQU =1

* MAIN INSTRUCTION DIAGNOSTIC FOR ITP

* LOCATIONS OF INTEREST TO THE USER, DOCUMENTED IN THE MANUAL

00000	ESTART	BRU	START
00001	EMTEST	BRU	MTEST
00002	EMTSTLP	BRU	MTSTLP
00003	EMTLOOP	BRU	MTLOOP
00004	BEGINS	ZR0	INSTAB
00005	ENDINS	ZR0	EINST=8
00006	BEGCHM	ZR0	0
00007	ENDCHM	ZR0	10000B

00010	EWAITF	DATA	=1
00011	DWAITF	DATA	=1
00012	IWAITF	DATA	0
00013	TRPTF	DATA	0
00014	IRPTF	DATA	0
00015	CHKSF	DATA	=1
00016	CPRINT	DATA	=1
00017	SPRINT	DATA	=1

00020	INSPTR	ZR0	0
00021	RESPTR	ZR0	0
00022	ARGPTR	ZR0	0
00023	ERRCNT	ZR0	0
00024	LCTR	ZR0	0
00025	ICTR	ZR0	0
00026	TCTR	ZR0	0

* MACROS

MSG MACRO D,G,1

LDA **2

BRU G(1)

MSTR (D(1))

G(1) BRM PMSG

ENDM

MSTR MACRO D

IF 940F

ZR0 **1

ASC 'D(1)'/

ELSE

ZR0 **1,3

Q1 NCHR (D(1))

RPT (Q2=1,Q1)

DATA ' D(Q2)'+240B