

1620 USERS GROUP  
WESTERN REGION  
MINUTES OF THE MEETING  
JUNE 17-19, 1964  
DENVER, COLORADO

ROBERT R. WHITE  
WESTERN REGION SECRETARY

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1620 USERS GROUP                      WESTERN REGION  
 SUMMER MEETING  
 JULY 17, 18, 19, 1964                      DENVER, COLORADO

ROSTER OF ATTENDEES

1032	MRS. BETTY CILSICK GEORGIA STATE COLL. ATLANTA, GA.	1084	LONA HECKART AEROSPACE CORP. PATRICK AFB., FLA.
1118	NANCY PAQUIN PUBLIC HEALTH SERV. ROCKVILLE, MD.	1118	JAMES E. DALY PUBLIC HEALTH SERV. ROCKVILLE, MD.
1170	LANNY L. HOFFMAN GUGGENHEIM LABS PRINCETON, N.J.	1177	LEON P. GOLDBERG PRINCETON U. PROJ ACC PRINCETON, N.J.
1238	RENE SEUIGNY JR. HAYES INTERNATIONAL HUNTSVILLE, ALA.	1238	W. S. HAMER HAYES INTERNATIONAL HUNTSVILLE, ALA.
1258	FRANCES K. DURKAN U.S.A.E.C. NEW YORK, N.Y.	1273	JUDITH KOERNER ARGONNE NAT. LAB. IDAHO FALLS, IDAHO
1290	ROBERT D. WEEMS N.C. STATE COLLEGE RALEIGH, N.C.	1352	WALKER R. HURD DPI, BOARD OF EDUC. MEMPHIS, TENN.
3016	THOMAS P. SODANO ITT FEDERAL LABS FORT WAYNE, IND.	3041	ARTHUR P. WOODS JR. ARMCO STEEL CORP MIDDLETOWN, OHIO
3053	EUGENE C. EWING NAT COOP REFINERY AS MC PHERSON, KAN.	3055	BARNEY T. WATSON VA HOSPITAL OMAHA, NEB.
3059	J. RICHARD BURROWS H.D.+R., ENGINEERS OMAHA, NEB.	3082	PAUL A. BICKFORD OU MED RES COMP CNTR OKLAHOMA CITY, OKLA.
3082	E. N. BRANDT JR., M.D. OU MED RES COMP CNTR OKLAHOMA CITY, OKLA.	3082	LEE HENDERSON OU MED RES COMP CNTR OKLAHOMA CITY, OKLA.
3089	CLARENCE B. GERMAIN CARTE CORP ST. PAUL, MINN.	3096	FLOYD M. CORE SEISMOGRAPH SERVICE TULSA, OKLA.
3107	S. THOMAS PARKER KANSAS STATE UNIV. MANHATTAN, KAN.	3108	LEORA THOMAS PIONEER HI-BRED CORN DES MOINES, IOWA
3118	J. D. PENDERGRASS SOUTHWESTER POWER AD TULSA, OKLA.	3136	EVERETT L. COOK UNIV. OF WICHITA WICHITA, KAN.
3206	LOLAFAYE COYNE MENNINGER FOUNDATION TOPEKA, KAN.	3216	J. ROBERTS BRITTON HAWAIIAN ELEC. CO. HONOLULU, HAWAII

3220	LOUIS BOVITZ HIBBING AREA TECH. HIBBING, MINN.	3220	WILLIAM J. MCGRAW HIBBING AREA TECH. HIBBING, MINN.
3227	JEROME F. FOCKE K.C. BOARD OF EDUC. KANSAS CITY, MO.	3238	WALTER G. ELWELL NEB. WESLEYAN UNIV. LINCOLN, NEB.
3242	RICHARD V. ANDREE UNIV. OF OKLAHOMA NORMAN, OKLA.	3242	GEORGE A. SPRADLING UNIV. OF OKLAHOMA NORMAN, OKLA.
3242	JACK L. MORRISON UNIV. OF OKLAHOMA NORMAN, OKLA.	3260	JOHN C. HARVEY IBM ROCHESTER, MINN.
3273	CHARLES WEISS USAF CHART + INFO ST. LOUIS, MO.	3277	JAMES L. GRISELL PHD LAFAYETTE CLINIC DETROIT, MICH.
3277	ROGER GUDOBBA LAFAYETTE CLINIC DETROIT, MICH.	3283	JOE D. PEGRAM MEMPHIS STATE UNIV. MEMPHIS, TENN.
3299	LEO DOUGLAS ARGONNE NAT. LAB. ARGONNE, ILL.	3302	ARNE GARNES CONCORDIA COLLEGE MOORHEAD, MINN.
3304	CAROL J. BILLINGHAM LEONARD REFINERIES ALMA, MICH.	3319	RAYMOND T. MCNAMARA DEFENSE SUB SUPPLY CHICAGO, ILL.
3326	JACK T. DUNN AVCO CORPORATION HUNTSVILLE, ALA.	3337	G. ENYEDY DIAMOND ALKALI CO. PAINESVILLE, OHIO
5001	W. W. JONES ACF INDUSTRIES ALBUQUERQUE, N.M.	5014	DAVE NIELSON GD/ASTRONAUTICS SAN DIEGO, CAL.
5018	JAMES W. HUNTER LA COUNTY ENGINEERS LOS ANGELES, CAL.	5020	MARILYN DOIG COLORADO STATE UNIV. FORT COLLINS, COLO.
5020	DR FRANKLIN GRAYBILL COLORADO STATE UNIV. FORT COLLINS, COLO.	5020	LARRY JACKSON COLORADO STATE UNIV. FORT COLLINS, COLO.
5021	JAMES C. IRVIN US ARMY, CORPS ENGR. ALBUQUERQUE, N.M.	5026	MARVIN J. CARR DOUGLAS AIRCRAFT CO. SANTA MONICA, CAL.
5027	MARVIN RUBENSTEIN ELECTRO OPTICAL SYS PASADENA, CAL.	5028	B. P. DUNCAN STEARNS-ROGER CORP. DENVER, COLO.
5028	D. I. CRANICHER STEARNS-ROGER CORP. DENVER, COLO.	5028	ED JURACEK STEARNS-ROGER CORP. DENVER, COLO.

5032	BOB H. MANNING GOODYEAR AEROSPACE LITCHFIELD PARK, ARIZ	5032	D. H. O'MERREN GOODYEAR AEROSPACE LITCHFIELD PARK, ARIZ
5032	MRS. N. A. KUFFEL GOODYEAR AEROSPACE LITCHFIELD PARK, ARIZ	5033	JOHN A. FERLING CLAREMONT MEN'S COLL CLAREMONT, CAL.
5041	RONALD E. WILDER MOTOROLA SEMICONDUCT PHOENIX, ARIZ.	5043	GLENN R. INGRAM MONTANA STATE COLL. BOZEMAN, MONT.
5058	L. L. KOPPIN SUNDSTRAND AVIATION DENVER, COLO.	5058	TED J. MCKENNA SUNDSTRAND AVIATION DENVER, COLO.
5060	CHARLES A. BETTINGER UNIV. OF TEXAS AUSTIN, TEX.	5073	PAUL KLEWEIN CAMPUS HIGH SCHOOL WICHITA, KAN.
5075	WENDELL T. BEYER UNIV. OF OREGON EUGENE, ORE.	5076	JAMES N. BOLES U.C. BERKELEY BERKELEY, CAL.
5077	PAUL BROWNE UNION OIL CO. BREA, CAL.	5086	HARRY CASTLE JR. PIONEER NATURAL GAS AMARILLO, TEX.
5086	B. G. GRANT PIONEER NATURAL GAS AMARILLO, TEX.	5089	DAVID BAER U.S.P.H.S. LAS VEGAS, NEV.
5095	LOWELL A. RASMUSSEN MEYERHAUSER CO. TACOMA, WASH.	5096	BOYD C. NORRIS BUREAU OF REC. SACRAMENTO, CAL.
5098	BURTON L. WILLIAMS WHITE SANDS M.R. WHITE SANDS, N.M.	5104	DR. J.R. GUINN TEXAS COLL. OF A+I KINGSVILLE, TEX.
5108	ELTON W. CHASE JR. CLARK COLLEGE VANCOUVER, WASH.	5117	ASTRIK BEIRMENDJIAN NAT CTR ATMOS RES BOULDER, COLO.
5126	WENDELL L. POPE UTAH STATE UNIV. LOGAN, UTAH	5130	A.A.J. HOFFMAN TEXAS CHRISTIAN UNIV FORT WORTH, TEX.
5131	ROGER HOFFMAN MARTIN CO. DENVER, COLO.	5131	E.E. EVANS MARTIN CO. DENVER, COLO.
5131	A.G. BISENIWS MARTIN CO. DENVER, COLO.	5133	RALPH D. PERRINE MASON + HANGER AMARILLO, TEX.
5133	LEWIS R. NOMBLE MASON + HANGER AMARILLO, TEX.	5143	GERALD W. LOCKE TEXAS TECH. COLLEGE LUBBOCK, TEX.

5144	JOSEPH A. STRAHL U.S. WEATHER BUREAU SACRAMENTO, CAL.	5145	ROBERT C. STEINBACH GROSSMONT COLLEGE SPRING VALLEY, CAL.
5147	RICHARD ROSANOFF NORTH AMERICAN AVIA. DOWNEY, CAL.	5150	JOHN M. GOODE HALLIBURTON CO. DUNCAN, OKLA.
5150	EUGENE BAKER HALLIBURTON CO. DUNCAN, OKLA.	5150	GEORGE V. COPLAND HALLIBURTON CO. DUNCAN, OKLA.
5150	GEORGE A. LARCADE HALLIBURTON CO. DUNCAN, OKLA.	5152	CARL J. REICH MONTEREY PEN. COLL. MONTEREY, CAL.
5162	JAMES W. BRUCE SAN DIEGO CO RD DEPT SAN DIEGO, CAL.	5164	KENNETH W. JONES COLO DEPT HIGHWAYS DENVER, COLO.
5166	JAMES M. ANDERSON LOCKHEED PROPULSION REBLANDS, CAL.	5168	WARD CROWLEY UNIV. OF IDAHO MOSCOW, IDAHO
5169	JOHN LAFLEUR SURVEY RES CTR U.C. BERKELEY, CAL.	5171	OLIVER BENSON UNIV. OF OKLAHOMA NORMAN, OKLA.
5171	CHARLES MAUDLIN JR. UNIV. OF OKLAHOMA NORMAN, OKLA.	5171	DAVE ASHBAUCHER JR. UNIV. OF OKLAHOMA NORMAN, OKLA.
5172	BERNARD BURGER HYDRA ELECTRIC CO. BURBANK, CAL.	5173	RICHARD G. SCHERER CENTRALIA COLLEGE CENTRALIA, WASH.
5174	L. PIERCE LOBERG COMPUTERMAT INC. LOS ANGELES, CAL.	5177	R.G. WAEDEMON TEXACO INC. PORT ARTHUR, TEX.
5181	ROBERT R. WHITE LA DEPT WATER POWER LOS ANGELES, CAL.	5189	L.E. HARVEY FOOTHILL COLLEGE LOS ALTOS HILLS, CAL.
5192	NGEL T. SMITH INDIANA STATE COLL. TERRE HAUTE, IND.	5200	AGNES W. PRUSZKA USAA DEFENSE BD FT. BLISS, TEX.
5200	LT. JACK KEEN USAA DEFENSE BD FT. BLISS, TEX.	5201	JOHN F. PEARSON SPRINGFIELD SCH DIST SPRINGFIELD, MO.
5202	ALFRED L. TAYLOR N.E. OKLA. A+M COLL. MIAMI, OKLA.	5202	WALTER E. MOORE N.E. OKLA. A+M COLL. MIAMI, OKLA.
5205	M.E. NETHERINGTON WESTERN STATE COLL. GUMMISON, COLO.	5214	WILLIAM C. PIQUETTE OTERO JUNIOR COLLEGE LA JUNTA, COLO.

5215 ROSEMARY PETERSEN  
WDPC-UCLA  
LOS ANGELES,CAL.

5217 MAURICE H. WITTEN  
FT HAYS KANS ST COLL  
HAYS,KAN.

5219 GEORGE G. TOWN  
SEATTLE UNIVERSITY  
SEATTLE,WASH.

5223 STANLEY WISNIEWSKI  
NORTHROP CORP.  
BEVERLY HILLS,CAL.

5226 B.S. SHANNON JR  
WADLEY RESEARCH INST  
DALLAS,TEX.

7021 P.J. REDBERGER  
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EDMONTON,ALBERTA

PAUL S. CHAN  
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RUTH C. MOSSMAN  
IBM  
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DAVID R. DYE  
IBM  
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RONALD R. LAKE  
IBM  
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IBM  
SAN JOSE,CAL.

DICK WILLIAMS  
IBM  
LOS ANGELES,CAL.

AUBREY D. WOOD  
IBM  
OKLAHOMA CITY,OKLA.

5215 JOHN W. RETTENMAYER  
WDPC-UCLA  
LOS ANGELES,CAL.

5217 JOHN B. O'LOUGHLIN  
FT HAYS KANS ST COLL  
HAYS,KAN.

5222 ROBIN YOUNG  
UNIV. OF NEW MEXICO  
ALBUQUERQUE,N.M.

5225 VICTOR W. HOFFMAN  
US WEATHER BUREAU  
FORT WORTH,TEX.

5226 J.M. HILL M.D.  
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PAUL L. TUAN  
UNIVERSITY OF UTAH  
SALT LAKE CITY,UTAH

REX TURCO  
US BUREAU OF RECL.  
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GERALD W. CALL  
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IBM  
DENVER,COLO.

STEVE LOPEZ  
IBM  
WHITE PLAINS,N.Y.

GORDON W. GOESCH  
IBM  
SAN JOSE,CAL.

DOROTHY MCGOWAN  
IBM  
LOS ANGELES,CAL.

CHARLES E. BERRY  
IBM  
LOS ANGELES,CAL.



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5408 S. UNIVERSITY AVENUE  
CHICAGO, ILLINOIS 60637

MINUTES OF THE MEETING - JUNE 17-19, 1964

The general meeting was opened on Wednesday, June 17, 1964, at the Brown Palace Hotel in Denver, Colorado, by the Western Region President, Paul Bickford. After opening remarks by Paul and the IBM representatives, the meeting was turned over to presentation of papers and workshops in parallel technical sessions. New users were welcomed to the group in a session held Wednesday evening. This was followed by the regular sound-off session which is reported on elsewhere in the minutes.

The technical sessions were held as scheduled in the agenda with the exception of changes in rooms. Copies of those papers available are included with these minutes.

The keynote speaker at the luncheon on Thursday, June 18, was Dr. E. N. Brandt, Jr., M.D., Ph.D. The subject of the very well-received talk was, "The Future of Computers in Medicine."

At the last general session on June 19, Dick Williams, Dave Dye, and Chuck Berry of IBM responded to the sound-off session. Paul Bickford announced that the next meeting will be a joint meeting of the Western and Mid-western regions and will be held on the campus of the University of Oklahoma in Norman, Oklahoma, on November 9, 10, and 11, 1964. The meeting was then adjourned.

ROBERT R. WHITE  
Western Region Secretary



## SOUND-OFF SESSION

The sound-off session was held Wednesday, June 17, 1964, with Chuck Maudlin as Chairman. The following is a report on most of the points made and the response from IBM representatives on Friday, June 19.

The question was raised about the speed of the IBM supplied FORTRAN compilers. It was stated that the old Bell Interpretive System for the 650 was faster than current compiler systems for the 1620. IBM's reply was that the compilers are relatively slow because they incorporate extensive source language error checking at compilation. The FORTRAN/FORMAT compiler is the fastest of the current IBM Fortrans for the 1620 because most of the error checking is done by the pre-compiler. Changes have been made to some of the FORTRAN compilers, notably II-D, Version 2, to provide for in-line compilation and allow error checking of mathematical operations at object time.

A request was made for an advanced monitor workshop which would supply enough information so that users can modify the Monitor program. Dick Williams stated that it might be possible to hold such a class for advanced systems programmers on a limited enrollment basis. Coincident with this request was a strong request for much more complete documentation on the software systems. Dick Williams stated that an operator's guide for changing Monitor will probably be available from his office and suggested that the Users Group bring more pressure to bear on IBM, through the Users Group Executive Council, to supply better documentation.

The complaint was made that some Systems Engineers are not filling out APAR forms when requested. It was stated that this is part of their job and if it is not done, the user should contact his branch manager.

Inquiries were made regarding possible future software systems and the answers were as follows:

- FORTTRAN IV - not being planned.
- Report Program Generator - not being planned.
- A non-disk loader for Monitor programs - not likely because of complexity.
- A disk-independent SPS with all disk and printer op codes - not being planned.
- Network Analysis - near future release.
- Linear Programming Package - release in July.
- Non-variable subroutines for FORTRAN II - now available on FTN II-D, Version II.

Requests for hardware additions were made and the answers to these were as follows:

Reset on Typewriter and 1622 - RPQ  
Punch and Read start on console - RPQ  
Programmable time clock - RPQ  
Tele-processing - RPQ 834308; approximate price -  
\$600/mo. including a 1026 adapter which can  
handle up to 26 - 1050 consoles on line.

There was some comment about the 1620 failing to operate properly outside of a temperature range of 70-80 degrees F. IBM representatives requested that users notify their C.E. so that corrective action can be taken.

In answer to a request for a flexible rental schedule for the first year of use, the reply from IBM was that no changes from present pricing policies are planned.

It was explained that the machine characteristic of not destroying the contents of the read buffer on a reset operation was designed on purpose. It was felt that the best method was to require positive action to destroy this information.

Inquiries about excessive down-time on the 1443 or the punch-read feed feature were answered by other users. Apparently this is not a general problem.

IBM announced the availability of COGO-ID. It is described in Bulletin H 20-4219-0. The program number is 1620-UG-05X and it requires a 20K 1620 with one 1311 disk pack.

This concluded the sound-off session.

WEDNESDAY - JUNE 17, 1964

8:00 Registration Outside Central City Room

8:30 Welcome and Opening Remarks. . . Paul Bickford, Regional President  
IBM Announcements. . . . . Dick Williams, Chuck Berry, Angelo Arena  
R. P. Paterniti  
Index Registers. . . . . Miss Dorothy McGowan, IBM

10:00 Coffee

Session A  
Central City Room

Session B  
Leadville & Silver Plume Rooms

10:30 A Least Squares Solution for A  
Range Measuring Instrumentation  
System -- Oliver Lee Kingsley  
and Burton L. Williams (5098)

IBM Monitor I Workshop

11:00 Boundary Value Problems in  
Ordinary Differential Equations  
with Constant Coefficients --  
Richard Rosanoff and Gordon Mah  
(5157)

IBM Monitor I Workshop  
(Continued)

11:30 (1) A program to check elementary  
machine language laboratory  
exercises  
(2) A 519 Simulator --  
R. C. Steinbach (5145)

IBM Monitor I Workshop  
(Continued)

12:00 Lunch ---On your own

1:15 Introduction to Matrices --  
C. E. Maudlin, Jr. (5171)

IBM Monitor I Workshop

3:00 Coffee

3:30 Simultaneous Linear Equations  
with Complex Coefficients --  
Nancy C. Kuffel (5032)

IBM Monitor I Workshop  
(Continued)

4:00 Applications of Numerical  
Filters in the Power Spectral  
Analysis of Stationary Time  
Series -- A. J. Hoffman (5130)

IBM Monitor I Workshop  
(Continued)

4:30 The Acquisition and Utilization  
of an IBM 1620 Computer in the  
U. S. Public Health Service --  
James A. Daly and Nancy A.  
Paquin (1118)

IBM Monitor I Workshop  
(Continued)

5:00 Adjournment of Day's Sessions

7:30 New Users Session. . . . . Paul Bickford. . . . . Stratton Room & )

8:15 Sound-off Session. . . . . C. E. Maudlin. . . . . Tabor Room )

THURSDAY - JUNE 18, 1964

8:00	Late Registration	Promenade Area
	<u>Session A</u> <u>Ballroom B</u>	<u>Session B</u> <u>Tabor &amp; Stratton Rooms</u>
8:30	How the IBM 1620 Assists Student Counselors at Junior College -- Paul S. Chan, IBM	Generalized Filter Network A/C Steady Analysis Program -- Davis H. O'Herren (5032)
9:00	(Continued)	FORTTRAN II - Debugging Techniques and Aids -- Leon P. Goldberg (1177)
9:15	AD-APT -- Bill Rogers, IBM	FORTTRAN II (Continued)
10:00	Coffee	
10:30	1620 Computer Utilization in a Wind Tunnel Data Acquisition System -- Stanley Wisniewski (5223)	FORTTRAN II (Continued)
10:50	Linear Programming -- Dr. S. T. Parker (3107)	FORTTRAN II (Continued)
12:00	Luncheon . . . . .Ballroom	
	Keynote Address: Future of Computers in Medicine. . .Dr. E. N. Brandt, Jr., M. D., Ph. D.	
1:45	IPL V -- Wendell T. Beyer (5075)	IBM Disk File Applications
3:00	Coffee	
3:30	Petroleum Exploration and Production Applications for the IBM 1620 and Plotter -- Jack L. Morrison (5171)	FORTTRAN II and The 1443 -- L. Hoffman (1177)
4:00	A Control Systems Approach to Automatic Jet Engine Testing -- Aubrey D. Wood, IBM	FORTTRAN II and The 1443 (Continued)
5:00	Adjournment of Day's Sessions	

FRIDAY - JUNE 19, 1964

8:00 Registration

Session A  
Ballroom B

Session B  
Onyx Room

8:30 SPS Tutorial Workshop --  
Clarence B. Germain

Panel on Education:

Chairman: Richard V. Andree  
University of Oklahoma

Members: Charles A. Bettinger  
University of Texas

Noel T. Smith  
Indiana State College

Donald L. Ferguson  
Campus High School

10:00

Coffee

10:30 SPS Tutorial (Continued)

Automatic Processing of AUTOSPOT  
and AUTOMAP Programs with the  
1620-131 Disc System --  
Jack T. Dunn (3326)

12:00

Lunch. . .On your own

1:00 FORTRAN Teaching Techniques --  
Wendell L. Pope

Workshop on 10:30 Paper

1:20 A Load and Go SPS with Monitor  
Control -- Kenneth M. Lochner  
and Glenn R. Ingram (5043)

Workshop on 10:30 Paper

1:55 How The 1620 is Used at  
Colorado State University --  
Franklin A. Graybill

Workshop on 10:30 Paper

2:30 IBM Reply From Sound-Off Session. . . . .Ballroom B

3:15

Meeting Adjournment



"A Least Squares Solution  
for a  
Range Measuring Instrumentation System"

by

Oliver Lee Kingsley  
and  
Burton L. Williams

Range Instrumentation Systems Office  
White Sands Missile Range  
New Mexico

June 17, 1964

A Least Squares Solution  
for a  
Range Measuring Instrumentation System

ABSTRACT

A brief development of the least squares equations for an instrumentation system that measures a set of ranges to obtain an estimate of the space position of a space vehicle. In addition, a method for obtaining the precision of that space position is included.

I INTRODUCTION

At White Sands Missile Range, the instrument that measures the radial range to an object in space is called a Distance Measuring Equipment or simply DME. An instrumentation system capable of giving Euclidean three space position estimates is termed as a DME/DME system. A typical DME/DME system usually consists of three non-colinear instrument or equipment sites used to measure range.

The typical solution equations are the classical deterministic set that rejects the minor image solution. A four or more DME system presents a problem because a slight error in any range measurement will not produce a set of homogeneous space position estimates by the classical approach. There is a need for a good method to combine

the set of overdetermined measurements into a single space point estimate. The least squares method will provide the required space point estimate if the set of measurements are unbiased.

## II LEAST SQUARES DATA REDUCTION EQUATIONS

The least squares equations developed minimize the sums of squares of the error set of range measurements.

The observational equation from which the error equation is derived is written:

$$(1) \quad R_{mi} = R_i + E_{mi}$$

where  $R_{mi}$  denotes the measured range from the  $i$ -th DME to the tracked vehicle.

$R_i$  denotes the true range from the  $i$ -th DME to the tracked vehicle.

$E_{mi}$  denotes the measurement error associated with the observation taken from the  $i$ -th DME.

The error equation is easily obtained from the observational equation, thus:

$$(2) \quad E_{mi} = R_{mi} - R_i$$

The true value needs to be replaced by some suitable approximation.

Later, the true value will be estimated by the final solution or

reduction equations. The assumption is made that the true value can be represented by the linear terms of a Taylor series about some nearby point  $R_o$ , where:

$$(3) R_o = R_o(X_o, Y_o, Z_o)$$

From the  $i$ -th DME, the range to the point  $(X_o, Y_o, Z_o)$  is written:

$$(4a) R_{oi} = \sqrt{(X_o - X_i)^2 + (Y_o - Y_i)^2 + (Z_o - Z_i)^2}$$

For any space position  $(X, Y, Z)$ , the equation becomes:

$$(4b) R_i = \sqrt{(X - X_i)^2 + (Y - Y_i)^2 + (Z - Z_i)^2}$$

The linear Taylor series representation is written:

$$(5) R_i = R_{oi} + \left. \frac{\partial R_i}{\partial x} \right|_o (X - X_o) + \left. \frac{\partial R_i}{\partial y} \right|_o (Y - Y_o) + \left. \frac{\partial R_i}{\partial z} \right|_o (Z - Z_o)$$

where

$$\left. \frac{\partial R_i}{\partial x} \right|_o = (X_o - X_i)/R_{oi}$$

$$\left. \frac{\partial R_i}{\partial y} \right|_o = (Y_o - Y_i)/R_{oi}$$

$$\left. \frac{\partial R_i}{\partial z} \right|_o = (Z_o - Z_i)/R_{oi}$$

The weighted error sums of squares for  $k$  measurements from  $k$  DME sites is written:

$$(6a) \sum_{i=1}^k E_{mi}^2 = \sum_{i=1}^k W_i [R_{mi} - R_i]^2$$

where  $w_i$  is the weight given to each measurement. The errors sums of squares for equally weighted measurements is written:

$$(6b) \quad \sum_{i=1}^k E_{mi}^2 = \sum_{i=1}^k [R_{mi} - R_i]^2$$

Generally, the components of the instrumentation system are near enough alike in performance and behavior that equation (6b) is applicable.

For a system that consists of heterogeneous distance measuring equipment (system components), each weight that would be inversely proportional to the equipments range variance would be appropriate.

The error sums of squares are minimized with respect to the range measurement parameter which is a function of the three orthogonal components X, Y and Z. The three resulting equations thus formed are called the normal equations from which the estimates  $\hat{X}$ ,  $\hat{Y}$  and  $\hat{Z}$  are obtained:

$$(7a) \quad \frac{\partial}{\partial X} \left[ \sum_{i=1}^k (R_{mi} - R_i)^2 \right] = 0$$

$$(7b) \quad \frac{\partial}{\partial Y} \left[ \sum_{i=1}^k (R_{mi} - R_i)^2 \right] = 0$$

$$(7c) \quad \frac{\partial}{\partial Z} \left[ \sum_{i=1}^k (R_{mi} - R_i)^2 \right] = 0$$

The constant terms not involving  $(X-X_0)$ ,  $(Y-Y_0)$  and  $(Z-Z_0)$  are placed on the right hand side of the equation:

$$(8a) \left\{ \sum_{R_{oi}} \frac{1}{2} \left[ (X_0 - X_i)(X_0 - X_i) \Delta X + (Y_0 - Y_i)(X_0 - X_i) \Delta Y + (Z_0 - Z_i)(X_0 - X_i) \Delta Z \right] \right\} = C_1$$

$$(8b) \left\{ \sum_{R_{oi}} \frac{1}{2} \left[ (X_0 - X_i)(Y_0 - Y_i) \Delta X + (Y_0 - Y_i)(Y_0 - Y_i) \Delta Y + (Z_0 - Z_i)(Y_0 - Y_i) \Delta Z \right] \right\} = C_2$$

$$(8c) \left\{ \sum_{R_{oi}} \frac{1}{2} \left[ (X_0 - X_i)(Z_0 - Z_i) \Delta X + (Y_0 - Y_i)(Z_0 - Z_i) \Delta Y + (Z_0 - Z_i)(Z_0 - Z_i) \Delta Z \right] \right\} = C_3$$

where

$$C_1 = \sum_{i=1}^k \left[ (R_{mi} - R_{oi})(X_0 - X_i) / R_{oi} \right]$$

$$C_2 = \sum_{i=1}^k \left[ (R_{mi} - R_{oi})(Y_0 - Y_i) / R_{oi} \right]$$

$$C_3 = \sum_{i=1}^k \left[ (R_{mi} - R_{oi})(Z_0 - Z_i) / R_{oi} \right]$$

$$\Delta X = \hat{X} - X_0$$

$$\Delta Y = \hat{Y} - Y_0$$

$$\Delta Z = \hat{Z} - Z_0$$

The equations can be written in compact form by matrix notation:

$$(9a) \quad A \begin{pmatrix} \Delta X \\ \Delta Y \\ \Delta Z \end{pmatrix} = \begin{pmatrix} C_1 \\ C_2 \\ C_3 \end{pmatrix}$$

or

$$(9b) \quad A \Delta = C$$

Solving for  $\Delta$  from equation (9b):

$$(10) \quad \Delta = A^{-1} C$$

The final solution for X, Y, and Z can be written:

$$(11) \quad \begin{pmatrix} X \\ Y \\ Z \end{pmatrix} = \begin{pmatrix} X_0 \\ Y_0 \\ Z_0 \end{pmatrix} + A^{-1} C$$

The necessary start point  $(X_0, Y_0, Z_0)$  can be obtained from a deterministic solution for three range measurements\*. The region for convergent solutions has not been fully explored at this time.

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\* Armijo, Larry, "Determination of Trajectories Using Range Data from Three Non-Collinear Radar Stations", Technical Memorandum 766, USASMSA, Sept. 1960, WSMR, N.M.

### III. A METHOD FOR ESTIMATING INSTRUMENTATION SYSTEM PRECISION

The term precision estimate refers to the standard deviation estimates for the coordinate data X, Y and Z from the instrumentation system. If there exists a common range measurement variance ( $\sigma_m^2$ ), then by use of the relative variance-covariance matrix,  $A^{-1}$ , from equation (10) estimates of the component variances can be obtained. The diagonal elements from the inverse matrix  $A^{-1}$  are used to estimate the component variance:

$$(12) \begin{pmatrix} \hat{\sigma}_X^2 \\ \hat{\sigma}_Y^2 \\ \hat{\sigma}_Z^2 \end{pmatrix} = \begin{pmatrix} A_{11} & 0 & 0 \\ 0 & A_{22} & 0 \\ 0 & 0 & A_{33} \end{pmatrix} \hat{\sigma}_m^2$$

where  $A_{11}$ ,  $A_{22}$  and  $A_{33}$  are from the matrix  $A^{-1}$ .

The total variance estimate is defined by the equation:

$$(13) \hat{\sigma}_T^2 = \hat{\sigma}_X^2 + \hat{\sigma}_Y^2 + \hat{\sigma}_Z^2$$

In terms of the matrix A, the total variance estimate is written:

$$(14) \hat{\sigma}_T^2 = \hat{\sigma}_m^2 \text{Tr } A^{-1}$$

where Tr denotes the trace of the inverse of the matrix A.



# THE BOUNDARY VALUE PROBLEM IN ORDINARY DIFFERENTIAL EQUATIONS WITH CONSTANT COEFFICIENTS

By Richard Rosanoff and Gordon Mah\*

## INTRODUCTION

In the process of solving seventh and eleventh order boundary value problems in closed form, the authors experienced an improvement in their understanding of the requirements of good formulation for a digital computer. Better control of numerical error and easier programming and debugging were the products of improved formulation.

The specific problems programmed were boundary value problems in ordinary differential equations with constant coefficients. The general solution of the ordinary differential equations with constant coefficients was known in 1739. It is mathematically simple. Its nature has been thoroughly explored. Perhaps just because of this, it provides an excellent vehicle for the study of the consequences of restrictions imposed by automatic digital computation.

To put together a good computer program, one must consider several ideas and points of view. Thus, although we shall be concerned with simple material, the structure of our paper is mildly complex. We include, therefore, the following outline of the material to be covered:

1. Mathematical preliminaries
  - a. The finite nature of the computer number system
  - b. The implications of the finite number system
  - c. A definition of linear dependence
  - d. Linear dependence in a finite number system
2. The computationally significant features of the physical problems used as illustrations. The equations and boundary conditions to be solved are also presented.

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3. Characteristic polynomials and their roots
4. Basis functions and boundary conditions
  - a. Definition of a basis and the mathematical requirements for their selection
  - b. The Wronskian
  - c. Boundary conditions in terms of the Wronskian
5. Selecting a basis
  - a. A group of bases
  - b. Physical criteria
  - c. Linear independence
    - (1) Wronskian
    - (2) Gramian
  - d. The effect on conditioning number of a dual coordinate system
6. Remarks and conclusions

## I. MATHEMATICAL PRELIMINARIES

The set of numbers which can be represented in a digital computer is not the set of real numbers. It is not even the set of rational numbers. In fact, it is a finite set. It is bounded from above and below, and it is nowhere dense on the real number system. It is not a field. For brevity, let us speak of the computer number set to mean the set of numbers available on a particular computer. For the real number system, it is true that if:

$$A \times B = C$$

and if A and B belong to the real number system, so does C. This is the property of closure. Closure does not hold for the computer number set.

$$10^{51} \times 10^{51} = 10^{102}$$

Both factors on the left are members of the set of numbers with which the IBM 1620 Computer operates. Their product is not a member of the set.

The equation:

$$A + 0 = A$$

Defines a unique zero in the real number system. But notice with eight significant figure arithmetic defining the computer number set:

$$e^{9.3} + e^{-9.3} = 10938.019 + .000091424231 = 10938.019 = e^{9.3}$$

yet  $e^{-9.3}$  is not zero because

$$e^{9.3} \times e^{-9.3} = .99999997$$

It will be seen that this difference in the number systems lies at the root of many numerical difficulties. The properties of functions must be reviewed in light of the number system available for the calculation.

One of the properties of sets of functions which is important to our discussion is the property of linear independence. By way of review we state that  $N$  functions  $\phi_i$   $i = 1, n$  are said to be linearly dependent if there exists a set of  $N$  constants  $C_i$   $i = 1, n$  not all zero such that:

$$\sum_{i=1}^n C_i \phi_i = 0$$

If no such set of  $C_i$  exist the functions  $\phi_i$  are said to be linearly independent.

As an example, we might consider three vectors in a physical space of three dimensions. If the vectors are not coplanar, then they are linearly independent. There are, at most, three linearly independent vectors in the space. Therefore, any fourth vector in the space can be completely described in terms of three linearly independent bases or coordinate vectors in the space. If, by some accident, a coplanar set of basis vectors were taken, we should be unable to completely represent a general vector in the space. We have a deficiency perpendicular to the plane of our base vectors.

The geometrical ideas are carried beyond a physical three space to  $N$  dimensional spaces in a very useful manner (Ref's. 5, 17). Further, a function itself, defined in some region, may be considered as a vector in a space of infinite dimensions, i. e., a function space. The directions in the space are associated with the real numbers in the region (Ref's. 18, 19).

Linear independence is essentially the same idea in these extensions to higher spaces. The complete description of some desired function as a linear combination of the members of a set of functions requires that these functions "span" the space. If any one of them can be represented by a finite linear combination of the others, it adds nothing to the description of the space.

On the other hand, linear independence over the real number system does not insure linear independence in a finite number system. For example:

$$C_1 \sinh(X) + C_2 \cosh(X) = \delta$$

$\delta$  may not be zero in the real number system, but  $\delta = 0$  on a computer, if  $X$  is large enough, and

$$C_1 = -C_2 \neq 0$$

## II. THE PHYSICAL PROBLEM

We have seen that the number system available to us for computation is limited and lacks some of the characteristics of the real number system in which we think. Later, we shall be faced with a choice of basis functions between bases which are completely equivalent in the mathematics of the set of real and complex numbers. Although we shall find adequate numerical arguments to justify a choice, there is also a physical argument. It is our point of view that the formulation which represents the physical problem most directly will probably be the best formulation from every point of view.

With this objective in mind, we need not spend the time required for the derivation of equations. We have a seventh order problem and an eleventh order problem which have been solved on our 1620. The seventh order problem arose in the calculation of stresses developed in the glue line of a lap joint such as is shown in Figure 1.

Two metal plates, A and B, (see Figure 1) of unequal thickness are joined by gluing in a lap joint of Length  $L$ . If the plates are pulled by forces at some distance from the joint the load must be transmitted from one plate to the other by shear stresses in the glue. Further, the eccentricity of the joint will introduce some bending, and the resulting rotation of the joint will result in a transverse component of the force at the edge of the joint. If the plates were completely rigid, one would expect the loads to be uniformly distributed along the glue line. In fact, however, the plates are not rigid but will elongate and bend in response to the loads they are carrying. Here lies the crux of the problem. At Point 1, Plate A is elongated. Plate B is not. We have almost a discontinuity in deformation. This requires a relative motion which is resisted by the glue. The resistance in the glue results in the transfer of load from Plate A to Plate B. The transfer of load eventually brings the plates to the same extension at some interior point. The situation at Point 2 is identical to that at 1 except that the roles of the plates are interchanged.

The important thing to remember, as we construct the solution, is that stresses arise at the boundaries due to the elastic deformation of the plates. These stresses are damped out as one moves to the interior. Once the two plates have developed equal extensions, the solution should be very similar to the solution for infinitely rigid plates.

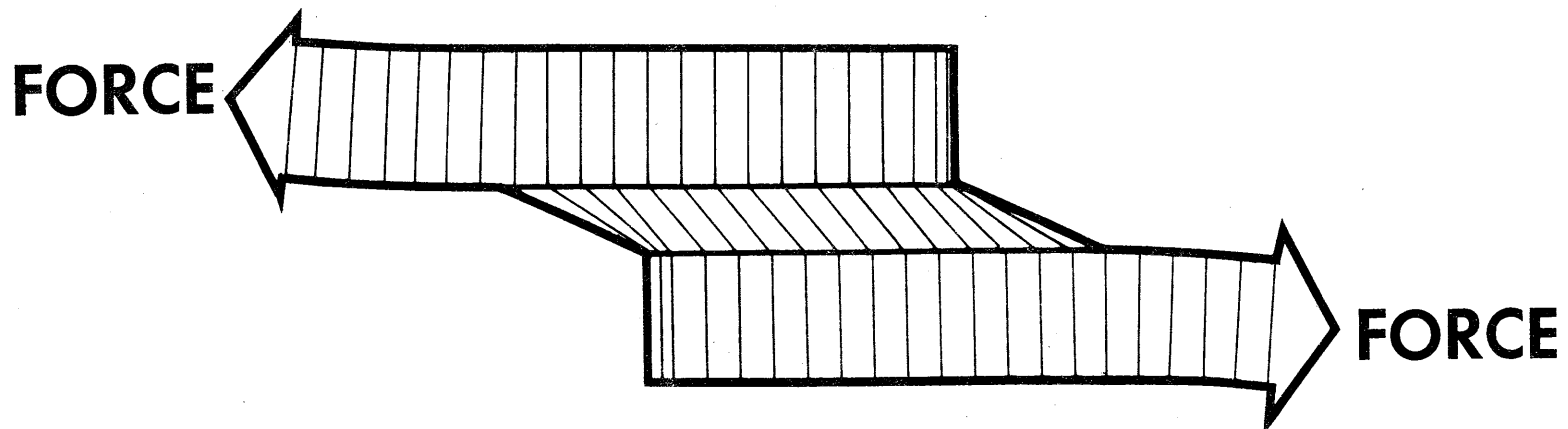
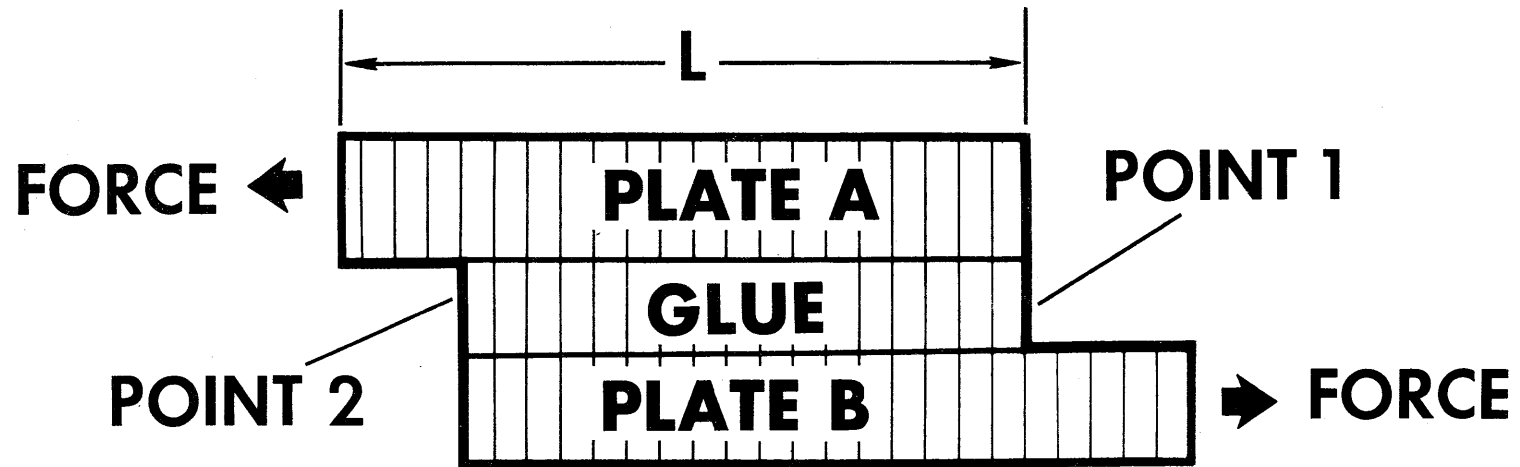


Figure 1. Physical Problem Represented by Seventh Order Equation

The shear stress in the glue is labeled  $\tau$ . The extension stress, transverse to the joint is labeled  $\sigma$ .

The analysis is an extension of a classical paper by Goland and Reissner (Ref. 2). The equations of equilibrium (under some simplifying assumptions) are:

$$\begin{aligned} \frac{d^3 \tau}{dx^3} - \xi_1 \frac{d\tau}{dx} - \xi_2 \sigma &= 0 \\ \frac{d^4 \sigma}{dx^4} + \xi_3 \sigma + \xi_4 \frac{d\tau}{dx} &= 0 \end{aligned} \quad (2.1)$$

which together give the homogenous differential equation:

$$\frac{d^7 \tau}{dx^7} - \xi_1 \frac{d^5 \tau}{dx^5} + \xi_3 \frac{d^3 \tau}{dx^3} - (\xi_1 \xi_3 - \xi_2 \xi_4) \frac{d\tau}{dx} = 0 \quad (2.2)$$

The parameters  $\xi_i$  are related to the elastic properties of the plates and glue line. We do not wish to burden the discussion with definitions of the notation. Instead, permit us to state the fact about the  $\xi_i$ , which is significant to the subject at hand. All of the  $\xi_i$  represent purely passive elements. They are geometric and elastic quantities.

The boundary conditions are given in the form of six inhomogenous differential equations of equilibrium which must be met, three at each boundary and one integral condition of equilibrium to be met over the domain. These may be written as:

$$\sum b_{i,j} \frac{d^j \tau(x)}{dx^j} = F_i$$

and

$$\int_0^L \tau(x) dx = \text{EXTERNAL SHEAR} \quad (2.3)$$

We shall summarize the 11th order problem more rapidly. The physical problem is once again the determination of stresses in a glued joint. For this problem however, we consider a section of honeycomb

sandwich material. It consists of a core with face plates A, shown in Figure 2. The section is assumed spliced to a similar honeycomb section through splice plates equal in thickness to the face plate.

The three equations of equilibrium are:

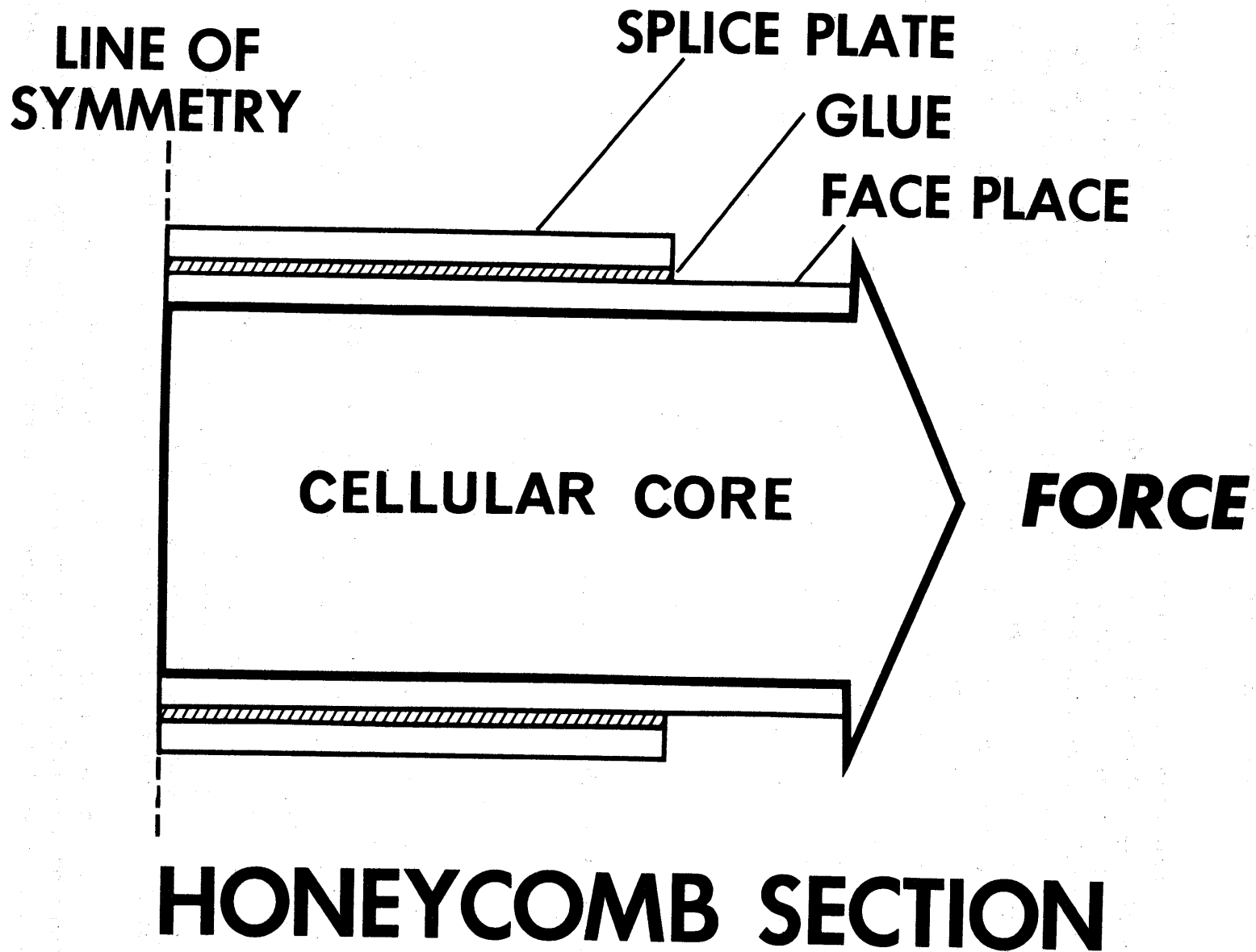
$$\begin{aligned}\sigma_c &= \frac{2h}{P_1} \frac{d^3 \tau}{dx^3} - \frac{4h}{3} \frac{d\tau}{dx} \\ \sigma_b &= \frac{1}{P_3} \frac{d^4 \sigma_c}{dx^4} + \sigma_c + \frac{h}{2} \frac{d\tau}{dx} \\ \frac{d^4 \sigma_b}{dx^4} + 2P_2 \sigma_b - P_2 \sigma_c &= 0\end{aligned}\quad (2.4)$$

A consequence of these equations is:

$$\begin{aligned}\frac{d^{11} \tau}{dx^{11}} - \frac{2P_1}{3} \frac{d^9 \tau}{dx^9} + (2P_2 + P_3) \frac{d^7 \tau}{dx^7} - \frac{P_1}{3} \left( 4P_2 + \frac{5P_3}{4} \right) \frac{d^5 \tau}{dx^5} \\ + P_2 P_3 \frac{d^3 \tau}{dx^3} - \frac{P_1 P_2 P_3}{6} \frac{d\tau}{dx} = 0\end{aligned}\quad (2.5)$$

The boundary conditions are similar in form to those of the 7th order problem.





# HONEYCOMB SECTION

Figure 2. Physical Problem Represented by Eleventh Order Equation

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### III. THE CHARACTERISTIC POLYNOMIAL

From the differential equations it is possible to write characteristic polynomials of seventh and eleventh order, respectively. As will be seen later, the roots of the characteristic polynomial provide a wealth of information concerning the nature of the solution, the relative importance of the various solution components in the derivative, hence, the boundary conditions. Even measures of the maximum stress and physical realism of the solution are provided in a qualitative way. They provide the best link between the physical parameters which occur in the coefficients of the differential equation and the solution.

Unfortunately, (Ref's. 3, 6 and 11) polynomials are subject to ill-conditioning and their roots may be poorly determined. In Reference 3 it is shown that the relative error in a root  $\alpha$  of a polynomial  $P(x) = \sum a_i X^i$  due to an error  $\delta a_s$  in the coefficient  $a_s$  is given by

$$\frac{\delta \alpha}{\alpha} = \frac{\alpha^{s-1} a_s}{P'(\alpha)} \frac{\delta a_s}{a_s}$$

Thus, it is only necessary for

$$\frac{\alpha^{s-1} a_s}{P'(\alpha)}$$

to be large to have a great magnification of relative error in the root as compared to the coefficient. There is a great proliferation of methods for obtaining polynomial roots (Ref's. 3, 4, 6, 7, 11, 12, and 13). The reason for this great number of methods is that none of them are generally satisfactory.

For the two problems solved, use was made of the easily variable precision of the 1620 version of Fortran. In the first place, the polynomials were studied throughout the data range of the problem, and solutions at different levels of machine precision were compared.

The zero root of the two polynomials is, of course, precisely known and removed without any contribution to error. The reduced polynomials both consist of even ordered terms only so that we may solve polynomials

of fifth and third rather than tenth and sixth order. The constant coefficient is negative so that the fifth and third order polynomial have at least one positive real root each. Thus, our original equations are known to have a pair of real roots which are equal in magnitude and opposite in sign. Following Lanczos (Ref. 5), we obtain this root, either directly or as a reciprocal, between zero and one. The method used was to repeatedly evaluate the polynomial by synthetic division and test to isolate the root in one of the half intervals.

To minimize contributions to error in subsequent roots due to the removal of the real roots, we requested a special synthetic division subroutine from our systems programmers. We wish to acknowledge the excellent job done by John Sherman of our Division in this respect. Mr. Sherman provided us with a compact subroutine which performs the entire synthetic division iteration at 28 decimal digit floating point. In this routine the numbers stored at 28 digits include only those currently needed in the calculation to provide a full high precision iteration so that any order polynomial may be handled without additional storage. The net storage required was less than required for the Fortran version of the same routine. The routine is now part of our 1620 Fortran library.

Investigation throughout the data range showed the remaining roots to be complex. We obtained the remaining roots of the 11th order polynomial by the use of Muller's Method, which, so far in our experience, has never failed to converge (Ref's. 6 and 7). We would welcome comments based on the experience of the audience in this regard. Because the roots may be obtained as the square roots of roots of a lower ordered polynomial, they are symmetrically distributed about the origin as shown in Figure 3.

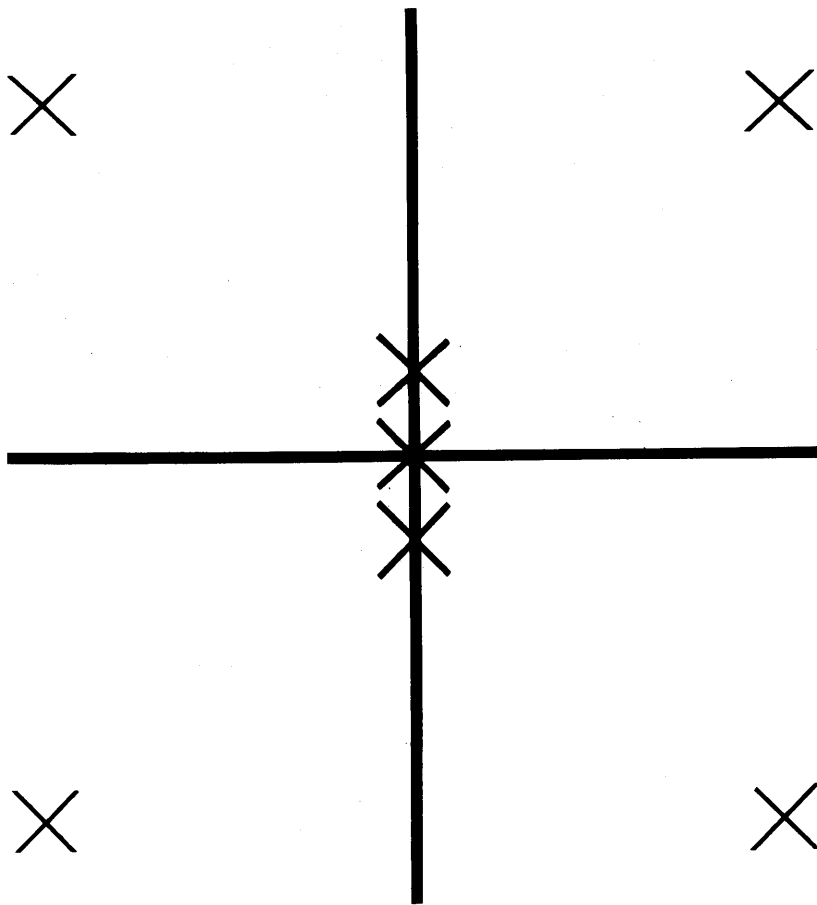


Figure 3. Map of Roots in the Complex Plane - Seventh Order Equation

#### IV. BASIS FUNCTIONS AND BOUNDARY CONDITIONS

The solution of a homogenous fourth order ordinary differential equation with constant coefficients whose characteristic polynomial has the roots  $\pm \alpha \pm i\beta$  can be represented as a linear combination of the functions.

$$\left\{ e^{(\alpha + i\beta)x}, e^{(\alpha - i\beta)x}, e^{(-\alpha + i\beta)x}, e^{(-\alpha - i\beta)x} \right\} \quad (4.1)$$

Of course, this is not the only set which could have been selected. Such a set is called a system or basis of functions if the member functions are linearly independent. A test for the linear independence is provided by the Wronskian determinant. This is a determinant whose first row is the system itself and whose  $j$ th row is made up of  $J-1$ st derivative of the function in the corresponding column.

We shall have a great deal more to say about linear independence in the next section of the paper. For the present, we wish to show how the programming and debugging of the problem are simplified by writing the derivative boundary conditions in terms of the Wronskian. Recall the form of the derivative boundary conditions.

$$\sum_{i=1}^{n-1} b_{i,j} \frac{d^{j-1} \tau(x)}{dx^{j-1}} = F_i \quad (4.2)$$

$j = 2, n \quad \text{at } x = 0 \quad \text{and } x = L$

We see that

$$\frac{d^{j-1} \tau(x)}{dx^{j-1}} \quad (4.2a)$$

is a column vector. We also see that recognition of the boundary conditions in the form (4.2) permits us to deal separately with the specification of boundary conditions ( $b_{ij}$ ) and the determination of the vector (4.2a). For our 11th order problem the matrix ( $b_{ij}$ ) was computed from a coded input. This provided a flexibility which was most useful when numerical difficulties were seen in the boundary conditions themselves. Such a scheme suggests the possibility of writing a generalized program.

Notice the vector (4.2a) is a function of X and exists in an N space. On the other hand, the matrix  $(b_{i,j})$  can be written as N linearly independent conditions at either  $X = 0$  or  $X = L$ . Thus, for the two-point boundary value, problem N conditions must be selected from 2N derivative conditions and possibly some integral conditions. This draws attention to the fact the boundary conditions selected must be such as to specify a unique solution. We shall not give adequate coverage to this problem in this paper. Let us now examine the vector (4.2a). It has a physical meaning without regard to the basis or coordinate system in which the solution is written. That is to say it could be written at various values of X as a table of numbers which would be independent of the manner in which it was obtained. To analyze it symbolically, however, we must assume a basis of functions, say  $\{\phi_i\}$ . The first element of (4.2a)

( $j = 1$ ) is the zeroth derivative, or the function

$$\tau(x) = \sum \phi_k(x) a_k \quad (4.3)$$

where the  $a_k$  are the constants of integration to be determined by the equations 4.2. But:

$$\frac{d^{j-1} \tau(x)}{dx^{j-1}} = \sum_{k=1}^n \frac{d^{j-1} \phi_k(x)}{dx^{j-1}} a_k = W(x) a_k \quad (4.4)$$

where  $W(x)$  is the matrix of functions from the Wronskian determinant.

The recognition of these matrix products is the key to relieving the program of unmanageable detail. As will be seen later, a factor in the choice of the basis is the ease with which the Wronskian may be developed by a simple set of do loops. The  $(b_{i,j})$  matrix may be checked separately. Any linear independence may be displayed in easier-to-recognize form. The program has pattern.

As we have indicated, the specification of boundary conditions which are sufficient and compatible involves more than we can discuss here. The reader is referred to Ref. 19.

## V. SELECTION OF THE SYSTEM OR BASIS FUNCTIONS FOR THE SOLUTION

Given an ordinary differential equation whose characteristic polynomial contains the non-repeated roots  $\pm \alpha \pm i\beta$  there are many choices of functions for the solution. Consider, for example, these four bases or systems:

$$\left\{ e^{(\alpha + i\beta)x}, e^{(\alpha - i\beta)x}, e^{(-\alpha + i\beta)x}, e^{(-\alpha - i\beta)x} \right\} \quad (5.1)$$

$$\left\{ \begin{aligned} &\text{Cosh}(\alpha x) \text{Cos}(\beta x), \text{Sinh}(\alpha x) \text{Cos}(\beta x), \text{Cosh}(\alpha x) \text{Sin}(\beta x), \\ &\text{Sinh}(\alpha x) \text{Sin}(\beta x) \end{aligned} \right\} \quad (5.2)$$

$$\left\{ e^{-\alpha x} \text{cos}(\beta x), e^{-\alpha x} \text{sin}(\beta x), e^{\alpha x} \text{cos}(\beta x), e^{\alpha x} \text{sin}(\beta x) \right\} \quad (5.3)$$

and:  $Z = L - X$

$$\left\{ e^{-\alpha x} \text{cos}(\beta x), e^{-\alpha x} \text{sin}(\beta x), e^{-\alpha z} \text{cos}(\beta z), e^{-\alpha z} \text{sin}(\beta z) \right\} \quad (5.4)$$

The question arises, is there a choice? If these functions are mathematically equivalent, which they are, can one set be superior to another for digital programming? The answer is yes. Let us first dispose of (5.1) on the arbitrary basis that we prefer not to perform complex arithmetic if we can avoid it.

The usual textbook treatment is to point out that the functions must be linearly independent. For the mathematician, this condition is met for all of the bases under discussion. For the digital programmer, however, this situation may be quite different. If the domain  $L$  of the solution is large enough, all of the sets of functions expressed in the computer number set become linearly dependent:

$$\begin{aligned} \lim_{\alpha L \rightarrow \infty} \frac{\sinh(\alpha L)}{\cosh(\alpha L)} &= 1 \\ \lim_{\alpha L \rightarrow \infty} e^{-\alpha L} \text{cos}(\beta L) &= \lim_{\alpha L \rightarrow \infty} e^{-\alpha L} \text{sin}(\beta L) = \lim_{\alpha L \rightarrow \infty} e^{-\alpha L \pm i\beta L} = 0 \end{aligned} \quad (5.5)$$

In the machine, this breakdown in linear independence becomes exact. That is to say, we may have  $\cosh\alpha L = \sinh\alpha L$  to the last digit. Of course one may expect computational difficulties long before the loss of the last tragic digit.

Is there any way out of this dilemma? Again the answer is yes. The moment we realize that the finite number of digits in the calculation limits our ability to produce the "exact" solution we begin to consider analogies with approximate methods. We seek functions to represent our solution which "look like" the solution. Clearly basis (5.4) is greatly superior in this light. Only (5.4) of the bases considered, directly represents a function which arises at disturbances at the boundaries and is damped as it proceeds to the interior of the region.

Comparing basis (5.4) with basis (5.2) it is seen that the functions of (5.2) approach each other in exactly the range in which they become large. The functions of (5.4), on the other hand, approach each other in exactly the range in which they drop out of the solution.

It is not surprising that the basis (5.4) is superior when we take note of the fact that it contains more physical information. Only basis (5.4) is cognizant of the location of the disturbance caused by mismatching strains at  $X = L$ .

Let us inquire into the physical significance of positive real parts of characteristic polynomial roots. In the initial value problem, with time as the independent variable, positive real parts have been used as criteria of stability. When space is made the independent variable, and the problem is formulated as an initial value problem, a numerical instability is quickly associated with positive real parts. Consider for example, the differential equation

$$EI \frac{d^4 y}{dx^4} = -ky \quad (5.6)$$

which is the homogenous equation for the deflection of a railroad track. In Ref. 15, the solution is shown in several forms including the method of initial conditions.



$$\begin{aligned}
y(x) = & y_0 \cosh(\lambda x) \cos(\lambda x) \\
& + \frac{\theta_0}{2\lambda} \left\{ \cosh(\lambda x) \sin(\lambda x) + \sinh(\lambda x) \cos(\lambda x) \right\} \\
& - \frac{M_0}{2\lambda^2 EI} \sinh(\lambda x) \sin(\lambda x) - \frac{Q_0}{4\lambda EI} \left\{ \cosh(\lambda x) \sin(\lambda x) - \sinh(\lambda x) \cos(\lambda x) \right\}
\end{aligned}
\tag{5.7}$$

where

$y_0$  = displacement at origin

$\theta_0$  = slope at origin

$M_0$  = bending moment at origin

$Q_0$  = shear at origin

In fact, the solution is damped as one moves away from a local load or disturbance. The method of initial conditions, however, requires that this damped function be composed of a linear combination of rapidly expanding functions.

In the finite arithmetic of the digital computer, this means that if a fly settles on the railroad track in Denver and the initial conditions are determined to the full capacity of the computer number set some place this side of Los Angeles, the railroad tracks will be ripped and torn in the most terrible carnage since World War II.

With time as the independent variable, the model is realistic. Positive exponentials mean positive feedback. The rate of growth of the function is proportional to the function with a positive coefficient of proportionality. "It takes money to make money." "Population growth is explosive in a favorable environment"—"chemical reactions become explosive if the rate of the reaction increases with the reaction products" (including the final energy as a product). But time moves on. Functions do not have causes today and develop into effects yesterday.

Space coordinates, however, are quite arbitrary. One does not need to establish a coordinate system to pull a glued joint. Thus, the physical meaning of the positive exponential in our glue line problem is seen to be a stereotyped pattern of coordinate system specification.

It is well known that the choice of a coordinate system can make the solution of a problem easier or more difficult (Ref. 21). In at least some problems it has great numerical consequences (Ref. 5). Interestingly enough, since the boundary conditions supply the information for the evaluation of the constants of integration, we see that they contain vital information about the origin of the coordinate system. One of the problems in specifying boundary conditions is to make sure they produce a unique solution in a defined coordinate system.

Stress analysts have traditionally handled positive exponentials by testing such a parameter as  $\alpha$  over the domain. If  $\alpha L$  is large (say  $\alpha L > 6$ ), he uses basis (5.3) as a special basis. He sets the arbitrary constants associated with positive exponentials equal to zero (when  $\alpha L > 6$ ) and calls his solution the semi-infinite case. We think the choice of basis (5.4) is superior for several reasons. For one thing, the characteristic polynomial may contain more than one set of complex roots with positive real parts. One set may correspond to a semi-infinite case and another to a short case. This situation actually occurs in our 11th order problem. Additionally, we should like to pick one basis and avoid programming more than one basis of functions.

If our mission in this paper were only to show that the basis (5.4) is superior for our problem, we could certainly bring this discussion to a close now. Our interest, however, is broader. We wish to explore the relative merits of bases in hopes that we may obtain criteria which help us in the solution of some other problem.

We have identified part of our numerical difficulties as arising from a breakdown in linear independence due to the finite nature of the computer number set. Linear independence is established if there are non-zero values of the Wronskian and Gramian. For basis (5.3) or (5.4), we rediscovered an ancient device for writing the Nth derivative which will render our discussion easier. Consider one of the functions, say  $e^{-\alpha x} \cos(\beta x)$

$$\phi = e^{-\alpha x} \cos(\beta x)$$

$$\frac{d\phi}{dx} = -\alpha e^{-\alpha x} \cos(\beta x) - \beta e^{-\alpha x} \sin(\beta x) \quad (5.8)$$

But identify

$$\alpha + i\beta = R(\cos(\theta) + i \sin(\theta))$$

$$\begin{aligned} \frac{d\phi}{dx} &= e^{-\alpha x} R \left( -\cos(\theta) \cos(\beta x) - \sin(\theta) \sin(\beta x) \right) \\ &= -e^{-\alpha x} R \cos(\beta x - \theta) \end{aligned} \quad (5.9)$$

It can be shown that:

$$\begin{aligned} \frac{d^j e^{-\alpha x} \cos(\beta x)}{dx^j} &= e^{-\alpha x} (-R)^j \cos(\beta x - j\theta) \\ \frac{d^j e^{-\alpha x} \sin(\beta x)}{dx^j} &= e^{-\alpha x} (-R)^j \sin(\beta x - j\theta) \end{aligned} \quad (5.10)$$

And since

$$Z = L - X \quad \frac{dz}{dx} = -1$$

$$\frac{d^j e^{-\alpha z} \cos(\beta z)}{dx^j} = e^{-\alpha z} R^j \cos(j\theta + \beta z) \quad (5.11)$$

$$\frac{d^j e^{-\alpha z} \sin(\beta z)}{dx^j} = e^{-\alpha z} R^j \sin(j\theta + \beta z)$$

In the seventh order problem we used the basis

$$\left\{ \begin{array}{l} e^{-\alpha_1 x} \cos(\beta_1 x), e^{-\alpha_1 x} \sin(\beta_1 x), e^{-\alpha_2 x}, 1, \\ e^{-\alpha_1 z} \cos(\beta_1 z), e^{-\alpha_1 z} \sin(\beta_1 z), e^{-\alpha_2 z} \end{array} \right\} \quad (5.12)$$

The jth row of the Wronskian determinant for this basis is:

$$\left\{ \begin{array}{l} e^{-\alpha_1 x} (-R_1)^{j-1} \cos(\beta_1 x - (j-1)\theta_1), e^{-\alpha_1 x} (-R_1)^{j-1} \sin(\beta_1 x - (j-1)\theta_1), \\ (-R_2)^{j-1} e^{-\alpha_2 x}, \frac{d^{j-1}(1)}{dx^{j-1}}, e^{-\alpha_1 z} (-R_1)^{j-1} \cos(\beta_1 z + (j-1)\theta_1), \\ e^{-\alpha_1 z} (-R_1)^{j-1} \sin(\beta_1 z + (j-1)\theta_1), R_2^{j-1} e^{-\alpha_2 z} \end{array} \right\} \quad (5.13)$$

Consider this Wronskian if L is large. When  $X = L$ , the first three columns of the determinant become very small. If  $X = 0$ , the last three become very small. It is shown in Ref. 22 that for an Nth order differential equation which does not contain the (N - 1st) derivative, the Wronskian is constant through the domain. To see that this is the fact for this Wronskian, it is helpful to factor the Wronskian into the three factors.

$$[W(X)] = [R] [T(X)] [E(X)] \quad (5.14)$$

Where  $[R]$  is the diagonal determinant whose elements are  $R_1^j$ ,  $[E(X)]$  is the diagonal determinant whose elements are the exponential terms

$$\left\{ e^{-\alpha_1 x}, e^{-\alpha_1 x}, e^{-\alpha_2 x}, e^{-\alpha_1 z}, e^{-\alpha_1 z}, e^{-\alpha_2 z} \right\}$$

and  $T(x)$  is the remaining determinant whose jth row may be written

$$\begin{aligned} & (-1)^j \cos(\beta_1 x - j\theta_1), (-1)^j \sin(\beta_1 x - j\theta_1), \\ & \left( -\frac{R_2}{R_1} \right)^j, \cos(\beta_1 z + j\theta_1), \sin(\beta_1 z + j\theta_1), \left( \frac{R_2}{R_1} \right)^j \end{aligned} \quad (5.15)$$

then

$$\begin{aligned} [R] &= R_1^{(1+2+3+4+5+6)} = R_1^{21} \\ [E(X)] &= e^{-(\alpha_1 + \alpha_1 + \alpha_2)(x+z)} = e^{-(2\alpha_1 + \alpha_2)L} \end{aligned} \quad (5.16)$$

With considerable manipulation of rows, the  $[T(X)]$  determinant can be seen to contain factors of the form

$$\cos^2(\beta x) + \sin^2(\beta x)$$

The Wronskian of basis (5.2) is less manageable, as may be seen from the single term:

$$\begin{aligned} w_{1,7} = & (\alpha^5 - 10\alpha^3\beta^2 + 5\alpha\beta^4) \sinh(\alpha x) \cos(\beta x) \\ & - (5\alpha^4\beta - 10\alpha^2\beta^3 + \beta^5) \cosh(\alpha x) \sin(\beta x) \end{aligned} \quad (5.17)$$

Numerical difficulties arise because columns 1 and 2 and columns 5 and 6 of the complete basis may be obtained from each other by replacing  $\sinh \alpha x$  with  $\cosh \alpha x$  or vice versa. Thus, the linear independence depends upon the ability to distinguish the hyperbolic functions at large values of  $X$  with a computer number set. Through considerable algebra the Wronskian may be shown to contain the factors  $(\cosh^2(\alpha x) - \sinh^2(\alpha x))$  and  $(\cos^2(\beta x) + \sin^2(\beta x))$ .

The difficulties with basis (5.3) interestingly enough do not arise in the same manner as with (5.2). Except for a minor phase difference in angles  $\beta x$  and  $\beta z$  the most significant difference between (5.3) and (5.4) lies in their respective  $[E(X)]$  determinants. For basis (5.3), this becomes

$$e^{(-\alpha_1 - \alpha_1 - \alpha_2 + \alpha_1 + \alpha_1 + \alpha_2)x} = e^0 = 1$$

so that on this consideration the basis is seen to be at least as good as basis (5.4). A similar pattern is observed if we consider the Gramian determinant.

For a basis of functions  $(\phi_i)$ , the Gramian determinate is given by

$$G_{i,j} = \int_0^L \phi_i \phi_j dx \quad (5.18)$$

we prefer to normalize the Gramian as to

$$G_{i,j}^* = \frac{\int_0^L \phi_i \phi_j dx}{\sqrt{\int_0^L \phi_i^2 dx} \sqrt{\int_0^L \phi_j^2 dx}} \quad (5.19)$$

In this form we may think of the functions  $\phi_i$  and  $\phi_j$  as being coordinate vectors in a function space. The terms  $G_{i,j}$  are then cosines of angles between the coordinate vectors. One may also think of them as simple correlation coefficients between the base functions.

Now consider the functions  $e^{-\alpha x} \cos(\beta x)$  and  $e^{-\alpha z} \cos(\beta z)$ .

Then

$$\begin{aligned} G_{1,3}^* &= \frac{\int_0^L e^{-\alpha x} \cos(\beta x) e^{-\alpha z} \cos(\beta z) dx}{\sqrt{\int_0^L e^{-2\alpha x} \cos^2(\beta x) dx} \sqrt{\int_0^L e^{-2\alpha z} \cos^2(\beta z) dx}} \\ &= \frac{e^{-\alpha x} \int_0^L \cos(\beta x) \cos(\beta z) dx}{\int_0^L e^{-2\alpha x} \cos^2(\beta x) dx} \quad (5.20) \end{aligned}$$

If  $L$  is great the definite integral in the denominator will be approximated by the lower limit so that the order of magnitude of  $G_{1,3}$  is  $e^{-\alpha L}$ . Thus, in this basis the functions separate into nearly orthogonal sets associated with the respective boundaries.

A similar situation holds for basis (5.3). The exponentials fall out of the numerator before integration and the denominator contains  $e^{\alpha L}$ . For basis (5.2), however:

$$\begin{aligned}
 G_{1,2}^* &= \frac{\int_0^L \cosh(\alpha x) \cos(\beta x) \sinh(\alpha x) \cos(\beta x) dx}{\sqrt{\int_0^L \cosh^2(\alpha x) \cos^2(\beta x) dx} \sqrt{\int_0^L \sinh^2(\alpha x) \cos^2(\beta x) dx}} \\
 &= \frac{\int_0^L (e^{2\alpha x} - e^{-2\alpha x}) \cos^2(\beta x) dx}{\sqrt{\int_0^L (e^{2\alpha x} + 2 + e^{-2\alpha x}) \cos^2(\beta x) dx} \sqrt{\int_0^L (e^{2\alpha x} - 2 + e^{-2\alpha x}) \cos^2(\beta x) dx}} \quad (5.21)
 \end{aligned}$$

How may we distinguish between basis (5.3) and (5.4)? Let us recall that our solution is in the form

$$\tau(x) = \sum a_i \phi_i(x) \quad (5.22)$$

This can be seen to be the vector  $\tau(x)$  in a function space, written in terms of its components  $a_i$  along the coordinate axes  $\phi_i(x)$ . We have seen that basis (5.2) is unattractive because the axes  $\phi_i$  become parallel. For the basis (5.3), the problem is the scaling of the axes. The length of the coordinate vectors is exactly the integral in the denominator of the Gramian. Thus, for a basis with distinct axes, such as (5.3), the choice of a second coordinate origin essentially normalizes the coordinate vectors.

Let us now consider the full matrix for the derivative boundary conditions. We may partition this matrix into four partitions as

$$\begin{array}{c}
 \left| \begin{array}{cc}
 S_{11} & S_{12} \\
 \hline
 S_{21} & S_{22}
 \end{array} \right| \quad (5.23)
 \end{array}$$

Where above the horizontal line we write the conditions to be met at  $X = 0$  and below the conditions to be met at  $X = L$ . To the left of the vertical line we write solution components associated with exponentials with negative real parts and to the right exponentials with positive real parts (or for basis 5.4 solution components in  $Z$ ).

We have two matrix equations (at  $X = 0$  and  $X = L$ ). They may be represented as:

$$[B] [R] [T(X)] [E(X)] = [S] \quad (5.24)$$

The matrices  $[B]$  are rectangular. They and the matrices  $[R]$  are unchanged from basis (5.3) to (5.4). The changes in  $[T(X)]$  are only differences in phase angle of the trigonometric functions. The matrix  $[E(X)]$ , however, is very different. At  $X = 0$ , the  $E$  matrix for basis (5.3) is the identity matrix. The  $E$  matrix for basis (5.4) is the diagonal matrix

$$\left\{ 1, 1, 1, e^{-\alpha_1 L}, e^{-\alpha_1 L}, e^{-\alpha_2 L} \right\}.$$

This same matrix multiplies the  $E$  matrix of (5.3) at  $X = L$ .

To give the  $E$  matrix of basis (5.4)

$$\left\{ e^{-\alpha_1 L}, e^{-\alpha_1 L}, e^{-\alpha_2 L}, 1, 1, 1 \right\}$$

If we denote the complete matrix for the boundary conditions for basis (5.3) as  $S$  and the diagonal transformation matrix

$$\left\{ 1, 1, 1, e^{-\alpha_1 L}, e^{-\alpha_1 L}, e^{-\alpha_2 L} \right\}$$

as  $C$ , we may write the boundary conditions in basis (5.4) as

$$[S] [C] [A] = [F]$$

except for the differences in phase angle noted above.

If we chose to pre-multiply  $[A]$ , rather than post-multiply  $[C]$ , it is seen that  $C$  represents a scaling of the coefficients  $[A]$ . This problem is identified by Lanczos (Ref. 5) as artificial ill-conditioning. His recommendation is just the sort of rescaling accomplished by the choice of basis (5.4).



## VI. REMARKS AND CONCLUSIONS

Using basis (5.4), we finally see the boundary condition equations in the partitioned form of the previous section in a physical light. The matrix  $S_{11}$  represents the semi-infinite problem at  $X = 0$ . The matrix  $S_{22}$  represents the semi-infinite problem at  $X = L$ . The matrices  $S_{12}$  and  $S_{21}$  represent cross-coupling between the semi-infinite solutions.

The final four figures show the solution and components for the seventh order problem solved with basis (5.4). Figure 4 is a plot of the shear stress for a typical joint. Figure 5 shows the components of the shear stress for this particular problem. The components associated with complex roots were too small to show on the same scale. Notice that the physical problem – two transients moving in from the boundaries, the solution, the solution components and the partitioned form of the matrix – all reflect the same pattern.

Figure 6 shows the peel stress solution. Figure 7 shows the components of the peel stress.  $R_1$  is the absolute value of the complex root  $\alpha_1 + i\beta_1$ .  $R_2$  is the value of  $\alpha_2$ . Recall with basis (5.4) the derivatives contain successively higher powers of the moduli  $R_i$ . Now  $R_1$  is 6.78, whereas,  $R_2$  is 1.77. One is not too surprised, then, when  $\sigma$  which is given by

$$\sigma(X) = \frac{1}{\xi_2} \left( \frac{d^3 \tau(X)}{dX^3} - \xi_2 \frac{d\tau(X)}{dx} \right)$$

contains larger components or the functions associated with the roots  $\pm\alpha_1 \pm i\beta_1$ .

Notice, also, that the magnitude of the real parts serve to determine how local the effects will be. Notice the components associated with  $\pm\alpha_1 \pm i\beta_1$  are much more rapidly damped than the components associated with  $\alpha_2$ .

This same problem was solved using basis (5.2). An interesting comparison of the numerical difficulties is provided by comparison of the conditioning number of the matrices which had to be inverted. The conditioning number is the ratio of the largest Eigenvalue to the lowest.

The logarithm of the conditioning number provides an estimate of the number of digits which will be lost in the inversion. For the basis (5.2) solution, a  $6 \times 6$  matrix was inverted. The conditioning number was  $2.6 \cdot 10^8$ . For the basis (5.4), two  $3 \times 3$  matrices are inverted (solution by partitioning Ref. 10). Their conditioning numbers were  $1.10 \cdot 10^2$ .

Our conclusions follow Hamming's beautiful statement (Ref. 1): "The Purpose of computing is insight, not numbers." Recognition of matrix products added tremendously to our insight and provided an unusual opportunity to see the nature of our numerical difficulties.

We confess to a strong interest in writing more solvable equations. The work which has been done on best approaches to the problem of solving poor equations, while very useful, has already run its course. Nothing but more digits will improve on the best methods available.

The problem of writing better equations is certainly not simple. Nor do we feel that we now know how. We do believe that the close imitation of the physical problem is a good clue. Further, for this problem we identified two mechanisms which could affect the equations. The basis (5.2) led to badly skew axes, the basis (5.3) to badly scaled axes. The double coordinate system improved the scaling. It is interesting to note that the normal equations of a least square approximation problem become highly skewed if the coordinate origin is very distant from the center of gravity of the function being approximated.

The possibility of writing a generalized program for the class of problems treated here looks good. If we were to write it, the first thing we should like to do is be very sure of our polynomial root routines. Another difficulty would be choosing the boundary conditions before the nature of the characteristic polynomial was established.

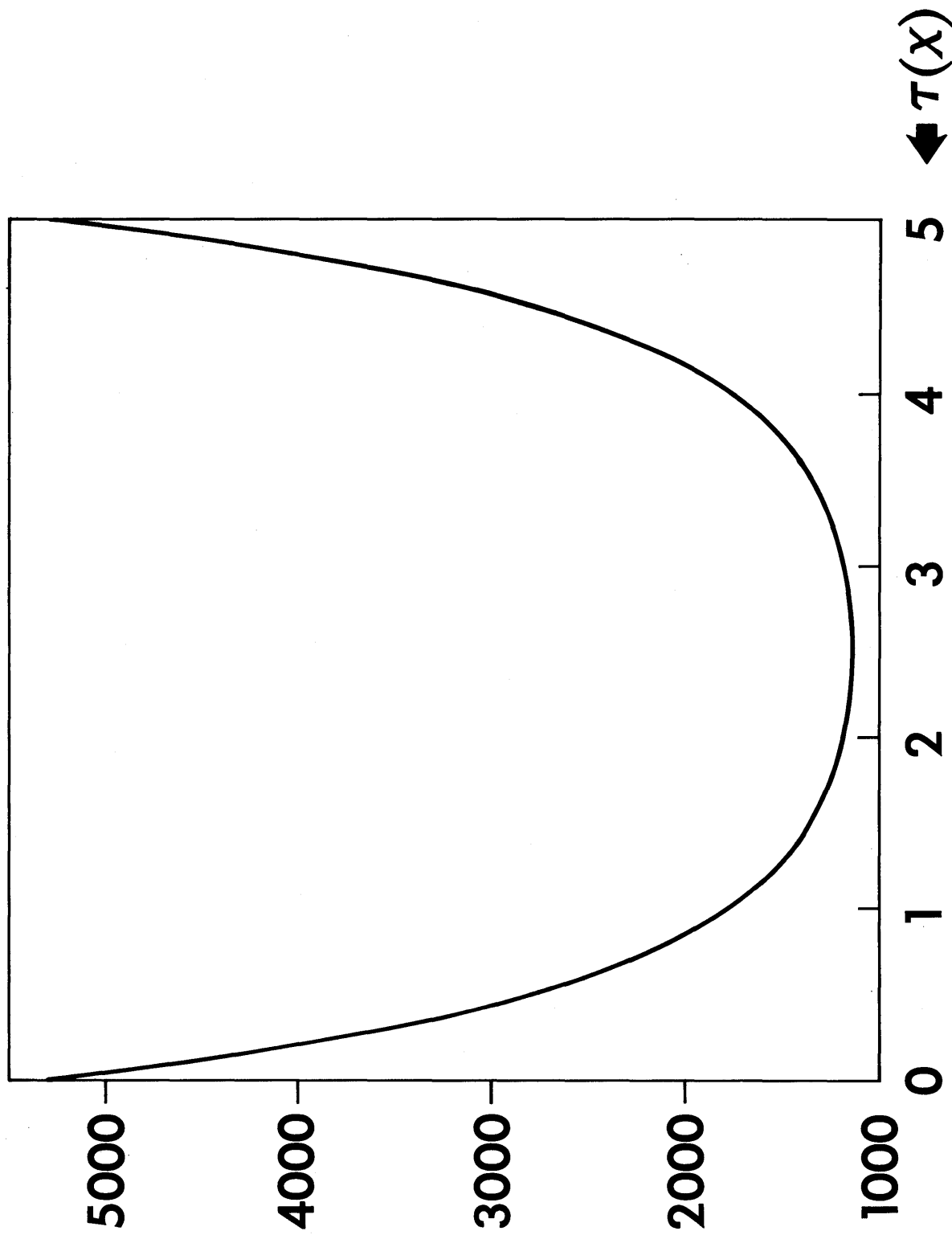


Figure 4. Shear Solution of the Seventh Order Problem

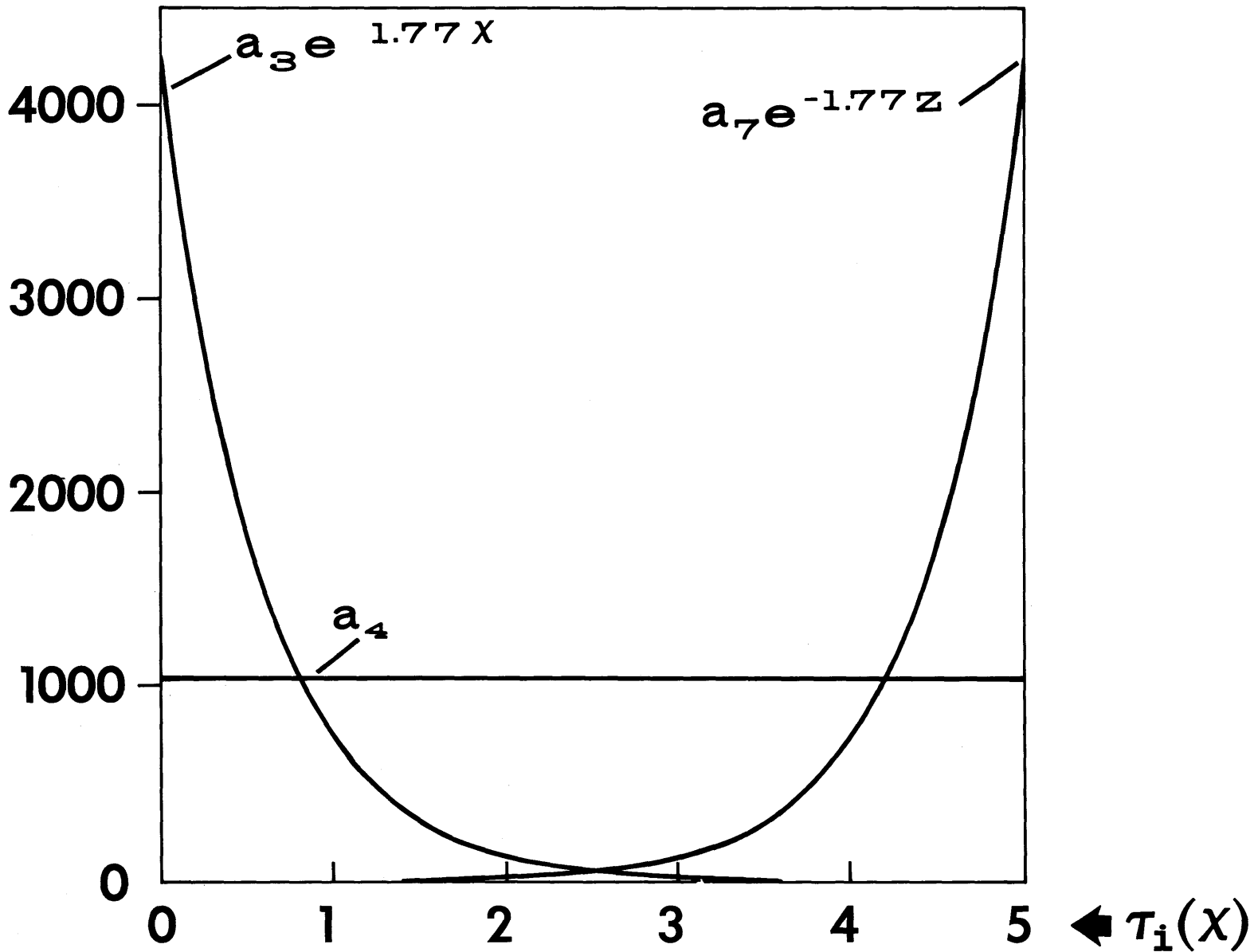


Figure 5. Components of Shear Solution (Note: Complex Components Were Too Small for Scale)

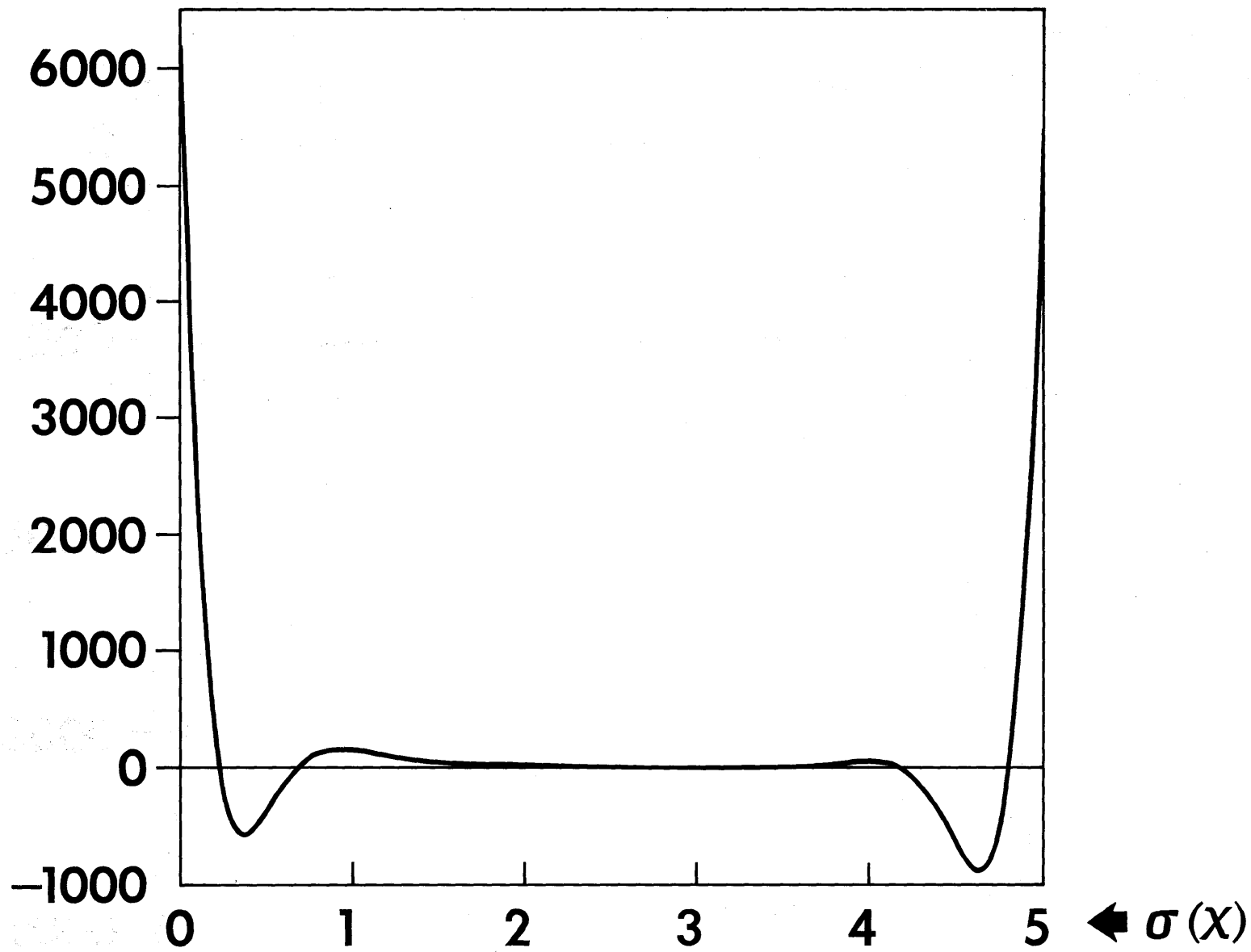
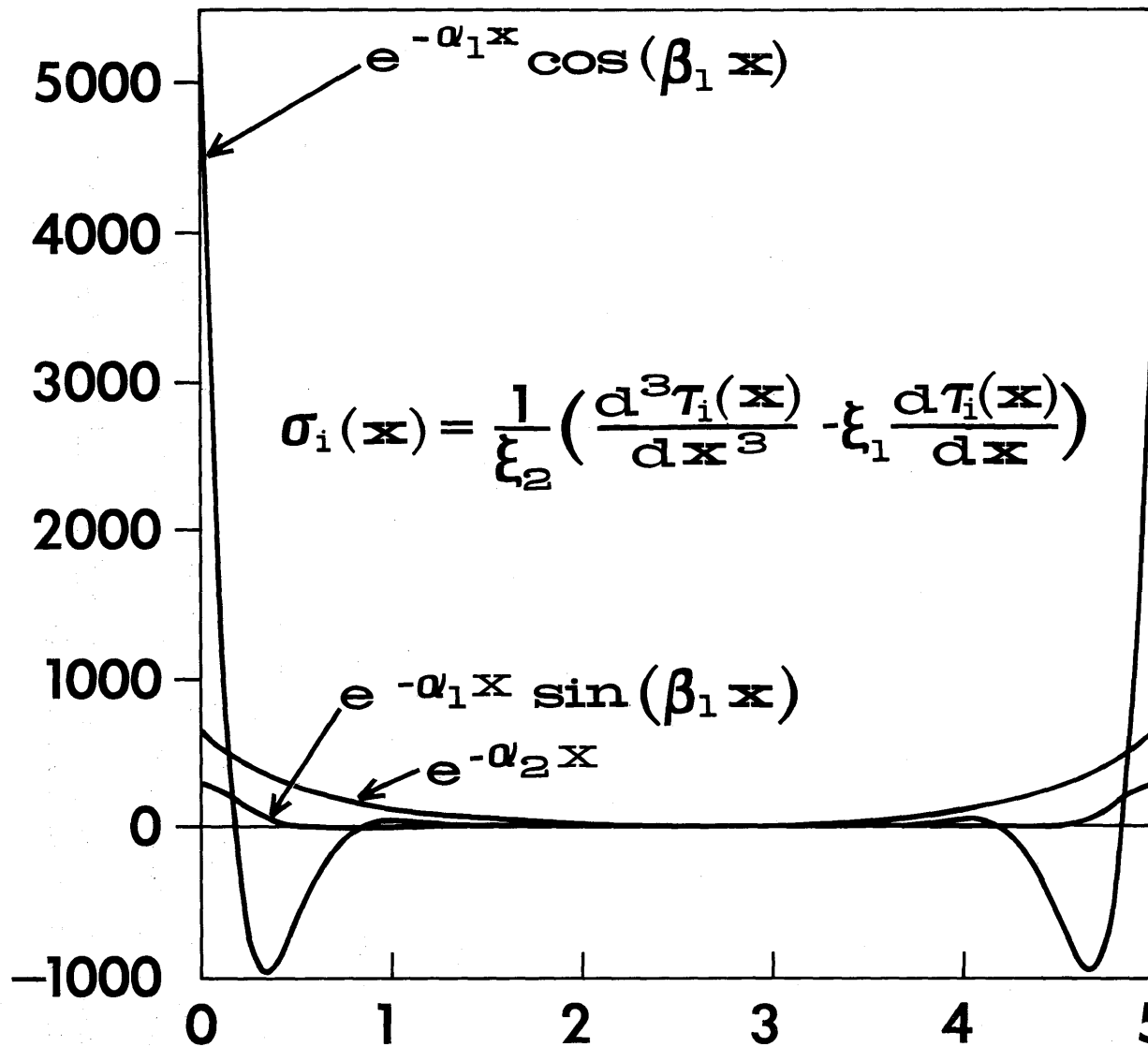


Figure 6. Peel Stress Solution Curve for Seventh Order Problem



$R_1 = 6.78$   
 $R_2 = 1.77$   
 $\alpha_1 = 4.796$   
 $\beta_1 = 4.794$   
 $\alpha_2 = 1.77$

Figure 7. Components of Peel Stress Solution, Seventh Order Problem\*

\*Curves are labeled with corresponding functions from the basis of the shear solution. As can be seen from peel stress expression, the true functions of this diagram are linear combinations of high order derivatives of the respective shear basis functions.

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## READER

A program to read and execute elementary  
machine language laboratory exercises  
R. C. Steinbach (5145)

### Introduction

Grossmont College is one of California's many public two-year colleges. These colleges provide three educational programs: (1) General education courses for the community, (2) Technical-vocational courses, (3) Transfer courses for students going on to four year institutions. Within the technical-vocational area Grossmont College has a data processing program containing a one year (four units per semester) computer programming course which begins with machine language. Students are capable of writing miniature machine language programs after approximately two lecture hours. The program described here monitors the student programs, allowing the student to see his program executed and relieving the instructor of the job of reading machine language programs.

### Student Program Format

During the first six weeks of the programming course the students are assigned specific problems to code. Examples of these problems can be found at the end of this paper. For each problem, each student hands in a deck of cards as follows: (See Figure 1) The first card or Header Card is used to identify each student's program. This card contains the student's name beginning in column one and ending with a record mark. It also contains a five digit identification number beginning in column 75 and a record mark in column 80. Reader uses this latter record mark to recognize the header card. The Program Cards follow the header card. The student machine language program is punched 72 digits (6 instructions) per card into as many cards as is necessary to a maximum of ten. A record mark in column 73 of a program card indicates

that column one of the next card follows column 72 of the card just read. Thus the last card (it may be the first and hence the only card of the program) has no record mark in column 73.

All programs return control to READER with a branch to 00000. This allows a manual restart (INSERT, RELEASE, START) if the student program hangs up and has not destroyed the READER program.

#### OPTIONS

During the time that the student has no knowledge of input/output instructions READER outputs the work area so that the student (and the instructor) may check the program results. This output may be suppressed using console switch 3 after the student is familiar with output instructions. The output device, either card punch or typewriter, for READER may be selected using console switch 4. This latter option allows remarks from READER to be output on the same device required of the student in a given problem.

#### A TYPICAL RUN

For each problem, the programs written by the students form a single deck which follows the READER object deck and four special data cards. (See Figure 2)

The first special data card contains program identification, console switch settings and tabulator information for the operator. The next three cards contain data for the student work area, e.g. numbers to add or subtract, negative numbers to count.

It is advisable to add an instructor written solution to this deck of 4 special data cards so that the students can see the right answers and see one way of writing the program. As far as READER is concerned, this is the first student program. Note that the 4 special data cards and the instructor

written program form a package which separates the reader object deck from the deck of student programs and which is easy to include for any given assignment.

READER types the program identification and operator message and halts. It then reads the three data cards, initializes the student work area and reads and executes the student programs as follows:

1. Search for Header Card. (Go to 3 when found; go to 2 on last card indicator.)
2. Type "All programs read" and halt. Press start to read next 4 special data cards and new batch of student programs.
3. Type student identification number.
4. Input student program, output student name and number of cards required for program.
5. Branch to student program. Return to 6 is automatic by student or manual by operator.
6. Output work area if switch 3 is on.
7. Initialize student work area.
8. Go to 1.

#### REMARKS

One should list the student program deck before doing anything else so that there is a permanent record of who turned in what. This is at least a partial defense against a charge of deck shuffling at execution time.

A clumsy student can wipe out core with a TF or TR. The only thing to do is reload the READER, but at least you have his identification number on the typewriter.

A loop, checkstop, or bad operation code can be noted by hand on the typewriter output and the READER restarted manually.

It is possible for a student to read the next student's program as data. As soon as this is obvious, a STOP, INSERT, R/S, will restart the READER. A comparison of the initial listing and the run listing will determine who was left out and his (their) program(s) can be placed at the end of the student program deck.

Conclusion

I would appreciate comments and criticism from any interested person. I do not plan to submit this to the Users Group Library until at least one more class has tried the system; they may think up new ways of giving the READER trouble.

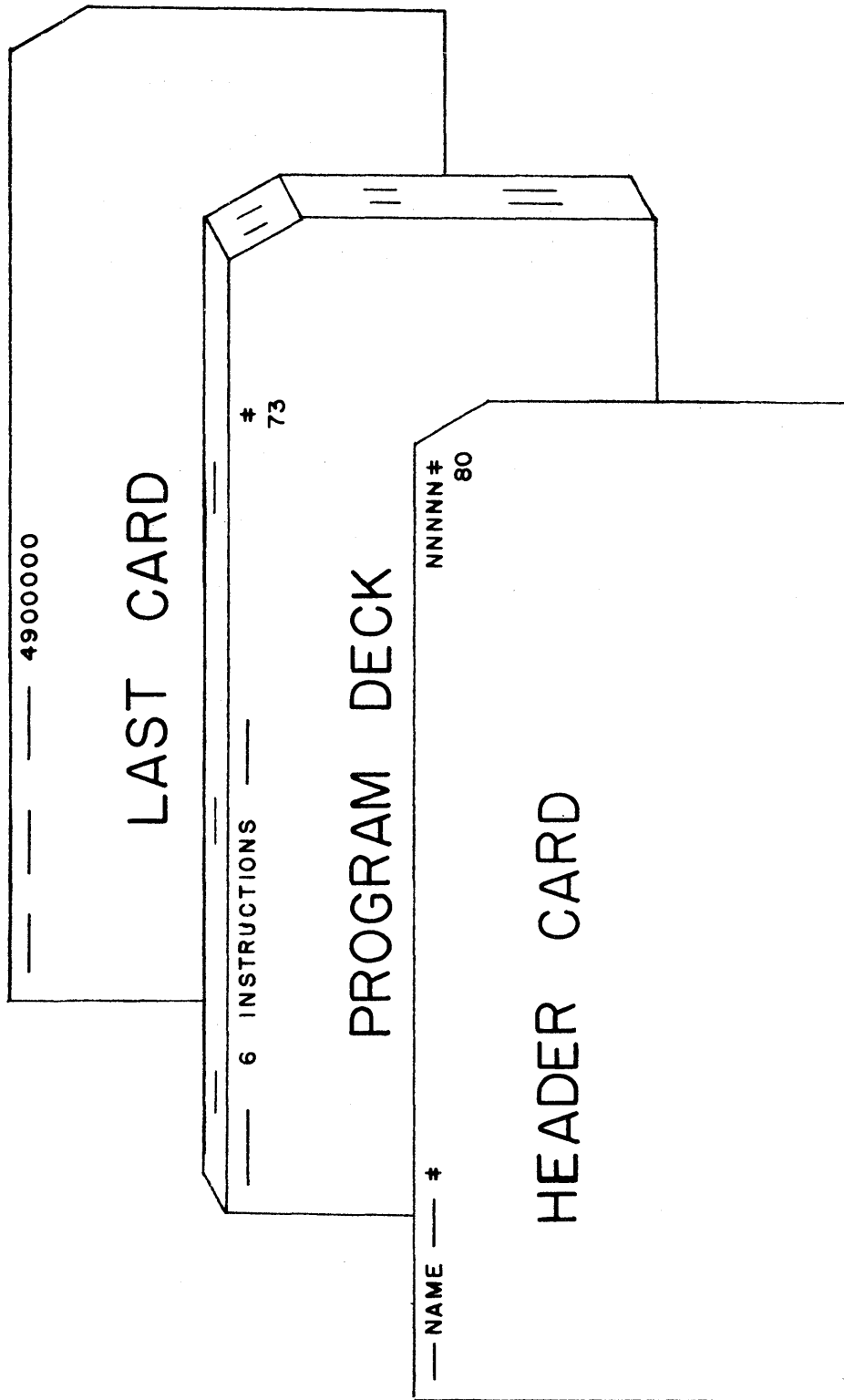


FIGURE 1

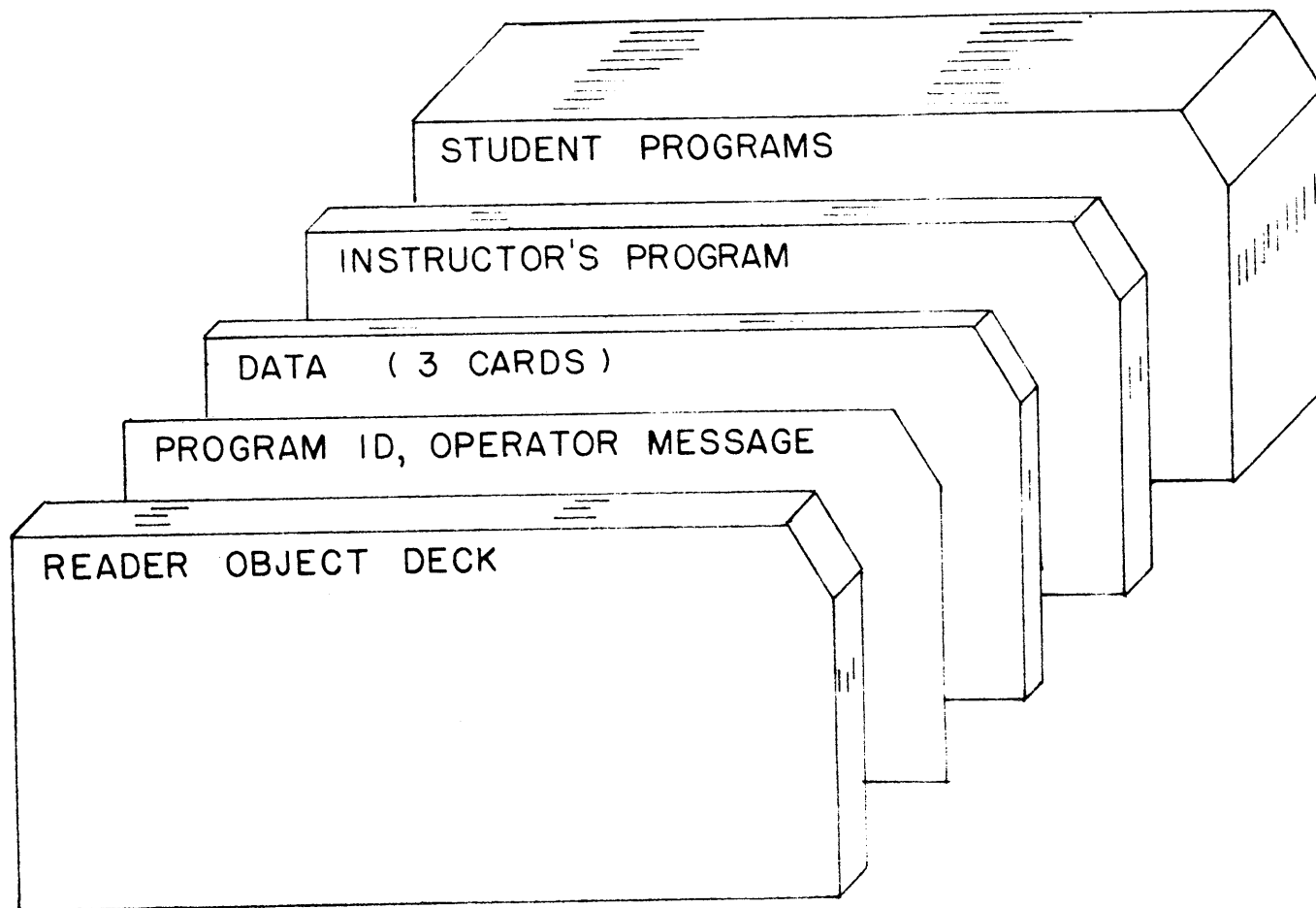


FIGURE 2

## LAB EXERCISE I

Numbers, described below, are in storage with the most significant digit flagged.

Number	Number of Digits	Address of least significant digit
A	6	7006
B	2	7010
C	4	7016
D	3	7021

1.1 Assume no overflow, numbers are integers.

Replace A by  $A+B$   
Replace C by  $C-B$   
Replace D by  $D-658$

1.2 Assume no overflow, numbers are integers.

Replace A by the integer  $A-2B+C-D$

1.3 Assume no overflow. Assume decimal locations as follows:

A = xxx.xxx  
B = .xx  
C = x.xxx  
D = xx.x

Replace A by  $A - D$   
Replace D by  $C + D$   
Replace C by  $C + 2.93$

### LAB EXERCISE 3

Note: Memory addresses above 11000 are available for your use. The first digit of your program is in 07300.

3.1 Return the carriage on the typewriter. Type out the numerical contents of 7001 - 7009, space the typewriter, type out the alphameric contents of 7030 - 7047. Return the carriage and type the numeric contents of 7030 - 7047. There are no record marks in place.

3.2 Return the carriage, type your name (25 character maximum), tabulate and type your code number.

3.3 As input to your program have one card with your name beginning in col. 1, and the words "1620 I/O PROGRAM" in col. 32-47, and a second card with 5 zeros, 5 ones, 5 twos, etc., and 5 nines in col. 1-50. Duplicate the two cards.

3.4 I will supply you with 3 cards which you will use as input to your program. Each card will have the following format:

A five digit number A in col. 6 - 10.  
A nine digit number B in col. 17 - 25.

You are to punch out three cards with the following format:

A and B as above  
A+B with low order digit in col. 40  
A.B with low order digit in col. 60

There are no flags on the input cards, and there should be no flags on the output cards.



## LAB EXERCISE 5

5.1 Type a message to turn on console switch 2 and then halt. If the switch is not set properly repeat the message and halt. Continue this process until the switch is on.

5.2 Two flagged 4 digit integers have their units position in 7005, and 7010 respectively. If the n-th integer is

less than 2222  
equal to 2222  
greater than 2222

} put a  $\begin{cases} 1 \\ 2 \\ 3 \end{cases}$  in 7011 + n

5.3 35 flagged 4 digit integers have their units position in 7004, 7008, ..., 7000 + 4n, ..., 7140. Tabulate the type-writer and type the number of negative numbers in the list.

5.4 Three 5 digit integers are located in 7005, 7010, and 7015 respectively; arrange them in ascending order in locations 7020, 7025, 7030.

A 519 Simulator  
R. C. Steinbach (5145)

Introduction

Card reproduction on the 1620 is not new; the most straight forward approach is to insert 371111100500 391111100400 4900000 R/S. The problem becomes slightly more complex if information is to be deleted, the columns permuted, sequence numbers added, and/or information gang punched into the cards. This paper describes one method of handling these other possibilities.

Method

During the first phase, the simulator sets up a table of source addresses. The first entry in the table is the address of the two digit field to be placed in column one of the output deck; the second entry addresses the source field for column two; etc. During the second phase, a card is read into an input buffer, 80 two-digit fields are transmitted from the appropriate source (the source table is addressed indirectly) sequentially into an output buffer. A card is punched and the next card read, and so on.

Format Cards

The deck to be reproduced is preceded by three format cards called INPUT, OUTPUT, and EMIT. All three cards must be there, however, the INPUT and EMIT cards may be blank. The input format card identifies the source of characters from the deck to be reproduced; the output format card identifies the destination of all characters to be punched in the new deck; the emit format card contains characters to be gang punched into all cards of the new deck.

The simulator produces the source table by scanning the output format

card. All columns that are blank in the output format card will be blank in the new, or output, deck. A field of 1's in the OUTPUT card indicates that the source is the same field on the old, or input, deck. A field of 2's (up to 5) indicates a sequence number field on the output deck. Note that this requires the OUTPUT card to be scanned from right to left. A field of 3's indicates that characters are to be emitted from the corresponding columns of the EMIT card. If a field of any other character, e.g. AAA or )))), is encountered on the OUTPUT card, then the INPUT card is searched for a corresponding field. The location of the field on the INPUT card determines the columns to be picked up in the old deck; the location of the field on the OUTPUT card determines the destination in the new deck. If the OUTPUT card contains a character other than the four special characters (blank, 1, 2, 3), that same character must appear on the input format card; furthermore, the field length defined must be the same. If either of these conditions fail, "Format card mismatch" is typed and the program will then accept new format cards. Figure 1 shows an example of the three format cards.

### Anomalies

Although it is not immediately obvious, the method chosen to set up the source table allows one field of the input deck to be placed in more than one field of the output deck. To accomplish this, a field indication on the INPUT card appears in several (non-adjacent) fields of the OUTPUT card. Two non-adjacent fields on the input card designated by the same non-special character will not be correctly interpreted.

Sequence numbers (even of different length) may also be punched in several non-adjacent fields.

### Modifications

Often, one wishes to change the emit characters whenever a master card is

detected. The variety of ways in which a master card may be indicated, and the number of possible reactions to a master card suggests one of the following manual solutions to the problem rather than a fully automated system.

If there are just a few decks headed by master cards, the same INPUT and OUTPUT cards may be used with a different EMIT card. The Master card may be used for an EMIT card if the master card is not to be duplicated and the characters to be emitted are in the correct columns.

If there are many master cards in a particular run, they may be detected using a compare or compare immediate after each card is read. A special routine is then added to the source deck to transmit characters from the Master card to the EMIT card image. The bulk of the routine can be instructions of the form  $TF\ EMIT-2+2*ecn, IN-2+2*mcn$  where  $ecn$  stands for emit column number and  $mcn$  stands for master column number. With the detect routine and transmit routine added, the source deck is reassembled.

### Conclusions

Any suggestions on ways to improve this program will be greatly appreciated. It will be submitted to the Users Group Library after these improvements are incorporated.

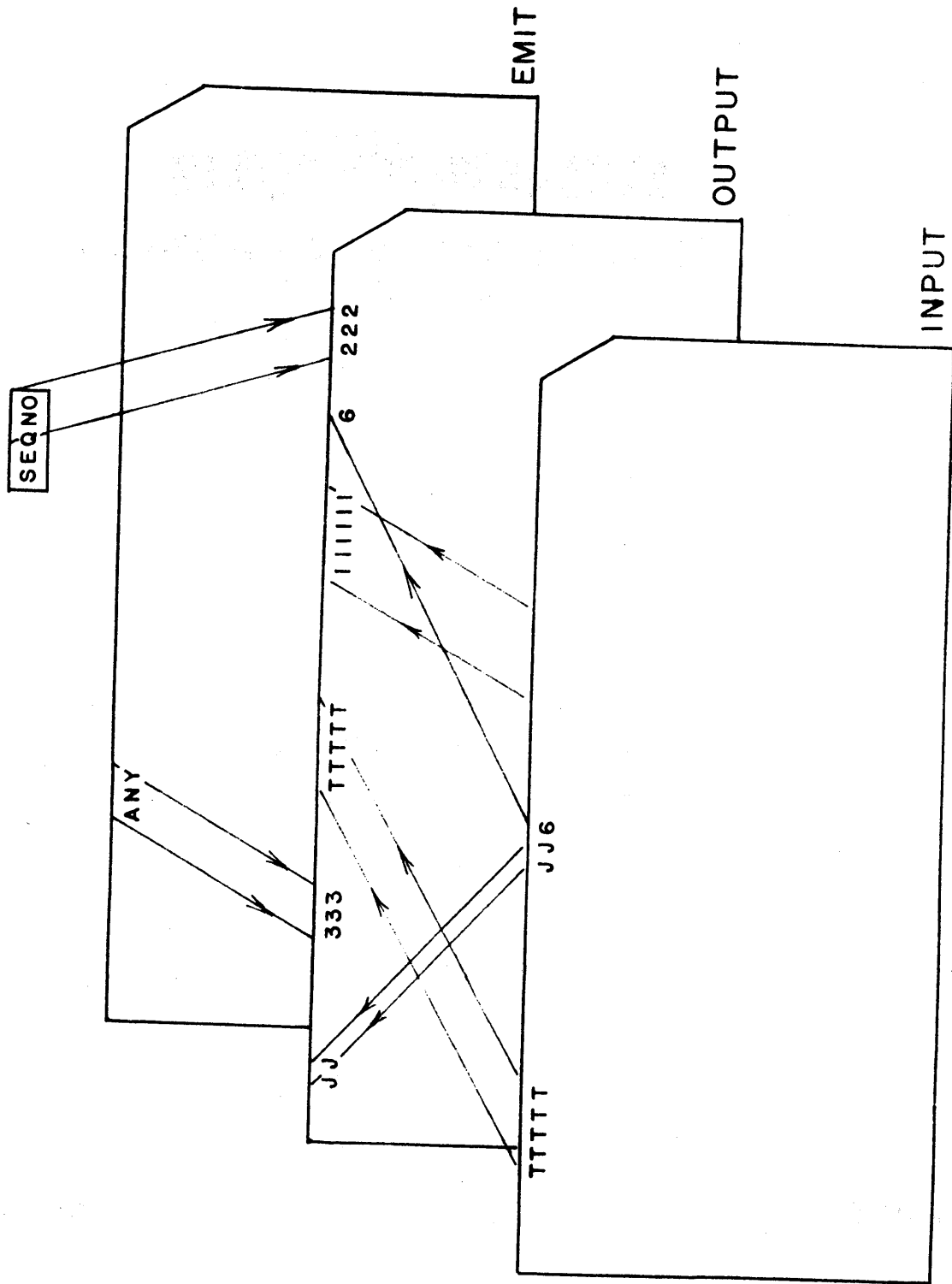


FIGURE 1

**GOOD YEAR**  
**GOODYEAR AEROSPACE**  
CORPORATION  
ARIZONA DIVISION  
LITCHFIELD PARK, ARIZONA

**SIMULTANEOUS LINEAR EQUATIONS  
WITH COMPLEX COEFFICIENTS**

**By**

**N. Kuffel**

**AAP-18906**

**May 1, 1964**

## SIMULTANEOUS LINEAR EQUATIONS

### WITH COMPLEX COEFFICIENTS

By

N. Kuffel

#### INTRODUCTION

This program solves simultaneous linear equations with complex coefficients resulting in complex roots. It was originally developed to solve large systems and has applications in mechanical and electrical engineering problems.

Of the numerous programs available for matrix inversion and simultaneous equations, very few take into account the under-and-overflow problems encountered on large matrix systems. There are no programs published at the present time for the 1620 for solutions of complex simultaneous equations, and very few available even for other machines. Several programs are available on the 1620 for real systems.

This program will solve up to 20 simultaneous linear equations with complex coefficients. Two forms of output results,  $A+jB$  and  $Ke^j$ , are available for either a specified limited number of unknowns, or for all unknowns up to 20. The program is written in Fortran with Format and requires 40 K memory.

#### GENERAL

Given a system of  $N$  simultaneous linear equations, in  $N$  unknowns, with complex (or real) coefficients, the program solves for the desired number of unknowns in terms of complex numbers. In certain situations, only a few of numerous unknowns are needed. Those desired can be rearranged to appear first in the equations. By specifying the number desired, only that number will be solved for, saving considerable computer time in the case of large systems.

## SIMULTANEOUS LINEAR EQUATIONS WITH COMPLEX COEFFICIENTS

Those equations to be solved are set in determinants of the form:

$$|Z| = |a| + j |b| \quad \text{where } a \text{ and } b \text{ are the coefficients.}$$

The application of Cramers' Rule gives us:

$$\begin{aligned} \frac{\omega}{|Z|} &= \frac{|c| + j |d|}{|a| + j |b|} \cdot \frac{|a| - j |b|}{|a| - j |b|} \\ &= \frac{(|a| |c| + |b| |d|) + j (|a| |d| - |b| |c|)}{|a|^2 + |b|^2} \end{aligned}$$

where  $Z$  is the determinant of the coefficients and  $\omega$  is the same determinant with the coefficients of the desired unknown replaced by the constant terms.

All determinants are evaluated by the triangular method, in which all elements to one side of the leading diagonal are computed to be zero. The determinant is equal to the product of the elements in the leading diagonal of the triangular determinant. This method of evaluation is preferable to that of expansion in terms of minors or the pivotal method because of the storage and time problem involved in the large complex systems.

Previous programs have made it necessary to do a manual rearrangement of data when a zero element is encountered on the diagonal, resulting either from the original coefficients or from subsequent computations. This program will check elements in the same column of the remaining rows of the determinant for a non-zero element. If such a value is found, a row interchange is performed, changing also the sign of the determinant. If no non-zero element is found we have the case of a zero determinant. If this occurs for the coefficient determinant, a message is typed out and a different method of solution must be found for this case of a nonsingular solution. A zero numerator determinant evaluates an unknown equal to zero, which is the correct result.

Over and underflow problems are quite common in matrix problems when doing accumulative operations, such as computing the product of the diagonal elements of the determinant. A scaling procedure has eliminated such difficulties in this program. Before multiplying, each diagonal element is scaled to the range between .1 and 1.0, storing an accumulative characteristic



## SIMULTANEOUS LINEAR EQUATIONS WITH COMPLEX COEFFICIENTS

(or power of ten) for the determinant, which is output with the product and then applied in the final division of determinants so that the end results have the correct magnitude.

Especially in the case of large systems, this program has been found to be as efficient even for real systems as most existing programs, particularly because of the row interchange and scaling procedures.

As many as 20 equations in 20 unknowns may be handled by this program on a 40K machine, which is minimum core for the program. The largest system run up to this time has been 18 equations, but no difficulties can be foreseen on any larger problems because of the scaling procedure.

The results are indicated in two forms. The actual outputs are the real and imaginary parts of the solution, as well as the magnitude and phase angle. These will give results in the forms:

$$A + jB \text{ and } Ke^{j\psi} \quad \text{where}$$

A = real part

B = imaginary part

K = magnitude

$\psi$  = phase angle in degrees

$$K = \sqrt{A^2 + B^2} \quad \psi = \tan^{-1} \frac{B}{A}$$

### SUMMARY

This program has been used numerous times for several months now, on systems from 3 equations to 18, both partial and complete solutions. Execution times on the 1620 MOD II have run:

3rd and 4th order - 1 min.

15th order - 20 min.

17th order - 29 min.

It should be noted that these times are dependent on the original set up of the coefficients and how many row interchanges are necessary.

## SIMULTANEOUS LINEAR EQUATIONS WITH COMPLEX COEFFICIENTS

The program is written in Fortran with Format and uses an ABSOLUTE VALUE subroutine. This can be easily changed in the source program if the subroutine is not readily available. Although the program presently begins at 6600, there is ample storage to recompile with a starting position of 8300 for other machine configurations. It would be a simple matter to change input and output modes to fit other needs and equipment. No sense switches are used. Sample input and output data follow in Appendix A and a program listing is in Appendix B.

## APPENDIX A

Sample input and output data listing follow. Input data follows the same sequence for all programs although Case 1 will be the only one described.

Case 1 - 3<sup>rd</sup> order complex system, complete solution

### Input

1<sup>st</sup> Card - NSOL = 3 (number of solutions desired) - I3 format  
Note statements 500 and 101 in program listing (Appendix B)

2<sup>nd</sup> Card - N = 3 (order of system) - I3 format

N X N (9) Cards - AR and AI (real and imaginary parts of the coefficients) - both values are on the same card in E14.8 format and are entered row-wise.  
Note statement 100 in Appendix B.

N(3) Cards - FR and FI (real and imaginary parts of the constants) - both values on the same card as were the coefficients.

### Output

Real and imaginary parts of the input coefficients

Real and imaginary parts of the input constants

Real and imaginary diagonal products, value of the coefficient determinant, scale factors for the products and the determinant, phase angle and magnitude.

Real and imaginary diagonal products, value of the determinant and scale factors for NSOL(3) solutions which include real and imaginary parts (A and B), phase angle ( $\psi$ ) and magnitude (K).

Case 2 - 4<sup>th</sup> order real system, complete solution

Case 3 - 4<sup>th</sup> order real system, partial solution

Case 4 - 3<sup>rd</sup> order real system, zero determinant

3  
3  
+.20110300E+04+.13140000E+03  
- .20550000E+04- .22700000E+01  
+ .00000000E-99+ .00000000E-99  
- .20550000E+04- .22700000E+01  
+ .16102980E+05+ .21747000E+03  
- .14170000E+05- .18500000E+03  
+ .00000000E-99+ .00000000E-99  
- .14170000E+05- .18500000E+03  
+ .22498000E+05+ .18500000E+03  
+ .00000000E-99+ .00000000E-99  
+ .00000000E-99+ .00000000E-99  
+ .83300000E+04+ .00000000E-99

SOLUTION OF SIMULTANEOUS LINEAR EQUATIONS  
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PROG. 223-63

ORDER 3

REAL	IMAGINARY
.20110300E+04	.13140000E+03
-.20550000E+04	-.22700000E+01
.00000000E-99	.00000000E-99
-.20550000E+04	-.22700000E+01
.16102980E+05	.21747000E+03
-.14170000E+05	-.18500000E+03
.00000000E-99	.00000000E-99
-.14170000E+05	-.18500000E+03
.22498000E+05	.18500000E+03

CONSTANTS

.00000000E-99	.00000000E-99
.00000000E-99	.00000000E-99
.83300000E+04	.00000000E-99

REAL PROD.	IMAG PROD.	DETERMINANT	
.22940795E-01	.25515200E-02	.53279032E-03	0
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			13
MULTIPLY DETERMINANT BY 1.0E			26

PHASE ANGLE = .63464644E+01 DEGREES  
MAGNITUDE = .23082251E-01 \* 1.0E 13

REAL PROD.	IMAG PROD.	DETERMINANT	
.24256068E-01	.34347995E-03	.58847480E-03	1
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			13
MULTIPLY DETERMINANT BY 1.0E			26

REAL PROD.	IMAG PROD.	DETERMINANT	
.23717163E-01	.18609040E-02	.56596678E-03	2
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			13
MULTIPLY DETERMINANT BY 1.0E			26

REAL PROD.	IMAG PROD.	DETERMINANT	
.23433940E-01	.21191024E-02	.55364013E-03	3
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			13
MULTIPLY DETERMINANT BY 1.0E			26

SOLUTIONS OF THE SIMULTANEOUS LINEAR EQUATIONS

ORDER 3

REAL	IMAGINARY	PHASE ANGLE	MAGNITUDE
------	-----------	-------------	-----------

.10460585E+01	-.10137221E-00	-.55351769E+01	.10509588E+01
.10301213E+01	-.33454431E-01	-.18600954E+01	.10306643E+01
.10191628E+01	-.20980622E-01	-.11793317E+01	.10193786E+01

4  
4  
+.30000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
- .10000000E+01+.00000000E-99  
+.10000000E+01+.00000000E-99  
+.10000000E+01+.00000000E-99  
- .10000000E+01+.00000000E-99  
- .20000000E+01+.00000000E-99  
+.40000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
+.30000000E+01+.00000000E-99  
+.10000000E+01+.00000000E-99  
- .20000000E+01+.00000000E-99  
+.50000000E+01+.00000000E-99  
- .20000000E+01+.00000000E-99  
+.30000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
+.10000000E+01+.00000000E-99  
+.30000000E+01+.00000000E-99  
- .20000000E+01+.00000000E-99  
+.00000000E-99+.00000000E-99

SOLUTION OF SIMULTANEOUS LINEAR EQUATIONS  
WITH COMPLEX COEFFICIENTS  
PROG. 223-63

ORDER 4

REAL	IMAGINARY
.30000000E+01	.00000000E-99
.20000000E+01	.00000000E-99
-.10000000E+01	.00000000E-99
.10000000E+01	.00000000E-99
.10000000E+01	.00000000E-99
-.10000000E+01	.00000000E-99
-.20000000E+01	.00000000E-99
.40000000E+01	.00000000E-99
.20000000E+01	.00000000E-99
.30000000E+01	.00000000E-99
.10000000E+01	.00000000E-99
-.20000000E+01	.00000000E-99
.50000000E+01	.00000000E-99
-.20000000E+01	.00000000E-99
.30000000E+01	.00000000E-99
.20000000E+01	.00000000E-99

CONSTANTS	
.10000000E+01	.00000000E-99
.30000000E+01	.00000000E-99
-.20000000E+01	.00000000E-99
.00000000E-99	.00000000E-99

REAL PROD.	IMAG PROD.	DETERMINANT		
.49999993E-02	.00000000E-99	.24999993E-04	0	
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			1.0E	4
MULTIPLY DETERMINANT BY 1.0E			1.0E	8

PHASE ANGLE = .00000000E-99 DEGREES  
MAGNITUDE = .49999992E-02 \* 1.0E 4

REAL PROD.	IMAG PROD.	DETERMINANT		
.19000000E-00	.00000000E-99	.36100000E-01	1	
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			1.0E	2
MULTIPLY DETERMINANT BY 1.0E			1.0E	4

REAL PROD.	IMAG PROD.	DETERMINANT		
-.28999992E-02	.00000000E-99	.84099953E-05	2	
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			1.0E	4
MULTIPLY DETERMINANT BY 1.0E			1.0E	8

REAL PROD.	IMAG PROD.	DETERMINANT		
-.50999986E-02	.00000000E-99	.26009985E-04	3	
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			1.0E	4
MULTIPLY DETERMINANT BY 1.0E			1.0E	8



REAL PROD.           IMAG PROD.           DETERMINANT  
.48999998E-02       .00000000E-99       .24009998E-04       4  
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E   -3  
MULTIPLY DETERMINANT BY 1.0E       -6

SOLUTIONS OF THE SIMULTANEOUS LINEAR EQUATIONS

ORDER   4

REAL	IMAGINARY	PHASE ANGLE	MAGNITUDE
.38000004E-00	-.00000000E-99	.00000000E-99	.38000003E-00
-.57999987E-00	-.00000000E-99	.00000000E-99	.57999986E-00
-.10199998E+01	-.00000000E-99	.00000000E-99	.10199997E+01
.98000006E-07	-.00000000E-99	.00000000E-99	.98000005E-07

4  
4  
+.30000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
- .10000000E+01+.00000000E-99  
+.10000000E+01+.00000000E-99  
+.10000000E+01+.00000000E-99  
- .10000000E+01+.00000000E-99  
- .20000000E+01+.00000000E-99  
+.40000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
+.30000000E+01+.00000000E-99  
+.10000000E+01+.00000000E-99  
- .20000000E+01+.00000000E-99  
+.50000000E+01+.00000000E-99  
- .20000000E+01+.00000000E-99  
+.30000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
+.10000000E+01+.00000000E-99  
+.30000000E+01+.00000000E-99  
- .20000000E+01+.00000000E-99  
+.00000000E-99+.00000000E-99

SOLUTION OF SIMULTANEOUS LINEAR EQUATIONS  
WITH COMPLEX COEFFICIENTS  
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ORDER 4

REAL	IMAGINARY
.30000000E-02	.00000000E-99
.20000000E-02	.00000000E-99
-.10000000E-02	.00000000E-99
.10000000E-02	.00000000E-99
.10000000E-02	.00000000E-99
-.10000000E-02	.00000000E-99
-.20000000E-02	.00000000E-99
.40000000E-02	.00000000E-99
.20000000E-02	.00000000E-99
.30000000E-02	.00000000E-99
.10000000E-02	.00000000E-99
-.20000000E-02	.00000000E-99
.50000000E-02	.00000000E-99
-.20000000E-02	.00000000E-99
.30000000E-02	.00000000E-99
.20000000E-02	.00000000E-99

CONSTANTS

.10000000E-02	.00000000E-99
.30000000E-02	.00000000E-99
-.20000000E-02	.00000000E-99
.00000000E-99	.00000000E-99

REAL PROD.	IMAG PROD.	DETERMINANT	
.49999993E-02	.00000000E-99	.24999993E-04	0
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			-8
MULTIPLY DETERMINANT BY 1.0E			-16

PHASE ANGLE = .00000000E-99 DEGREES  
MAGNITUDE = .49999992E-02 \* 1.0E -8

REAL PROD.	IMAG PROD.	DETERMINANT	
.19000000E-02	.00000000E-99	.36100000E-05	1
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			-8
MULTIPLY DETERMINANT BY 1.0E			-16

REAL PROD.	IMAG PROD.	DETERMINANT	
-.28999992E-02	.00000000E-99	.84099953E-05	2
MULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E			-8
MULTIPLY DETERMINANT BY 1.0E			-16

SOLUTIONS OF THE SIMULTANEOUS LINEAR EQUATIONS

ORDER 4

REAL	IMAGINARY	PHASE ANGLE	MAGNITUDE
.38000004E-00	-.00000000E-99	.00000000E-99	.38000003E-00
-.57999987E-00	-.00000000E-99	.00000000E-99	.57999986E-00

3  
3  
+.10000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
+.30000000E+01+.00000000E-99  
+.10000000E+01+.00000000E-99  
+.10000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
+.20000000E+01+.00000000E-99  
+.10000000E+02+.00000000E-99  
+.50000000E+01+.00000000E-99  
+.15000000E+02+.00000000E-99

SOLUTION OF SIMULTANEOUS LINEAR EQUATIONS  
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PROG. 223-63

ORDER 3

REAL	IMAGINARY
.10000000E+01	.00000000E-99
.20000000E+01	.00000000E-99
.20000000E+01	.00000000E-99
.30000000E+01	.00000000E-99
.10000000E+01	.00000000E-99
.10000000E+01	.00000000E-99
.20000000E+01	.00000000E-99
.20000000E+01	.00000000E-99
.20000000E+01	.00000000E-99

CONSTANTS

.10000000E+02	.00000000E-99
.50000000E+01	.00000000E-99
.15000000E+02	.00000000E-99

ZERO DETERMINANT - USE DIFFERENT METHOD OF SOLUTION

APPENDIX B

PAGE 01

SOURCE PROGRAM

07000 C SOLUTION OF SIMULTANEOUS LINEAR EQUATIONS WITH COMPLEX  
07000 C COEFFICIENTS USING CRAMERS RULE  
07000 C PROG NO 223-63  
07000 C DETERMINANTS EVALUATED BY THE TRIANGULAR METHOD  
07000 C ELEMENTS ENTERED ROW-WISE  
07000 C  
07000 C AR-REAL PART OF DETERMINANT  
07000 C AI-IMAGINARY PART OF DETERMINANT  
07000 C REAL AND IMAGINARY PARTS OF A COEFFICIENT ENTERED ON SAME CARD  
07000 C FR-REAL PART OF CONSTANT TERM  
07000 C FI-IMAGINARY PART OF CONSTANT TERM  
07000 C N-ORDER OF THE SYSTEM  
07000 C NSOL - NUMBER OF SOLUTIONS DESIRED (EQUAL TO OR LESS THAN N)  
07000 C DIMENSION AR(20,20),AI(20,20),FR(20),FI(20),WR(20,20),WI(20,20)  
07000 C DIMENSION XR(20), XI(20)  
07000 C CONV = 180./3.14159265  
07048 500 READ 101, NSOL  
07072 READ 101, N  
07096 101 FORMAT (I3)  
07118 PUNCH 104  
07142 104 FORMAT (//15X41HSOLUTION OF SIMULTANEOUS LINEAR EQUATIONS)  
07296 PUNCH 105  
07320 105 FORMAT (22X, 25HWITH COMPLEX COEFFICIENTS/29X12HPRG. 223-63//)  
07552 PUNCH 119, N  
07576 119 FORMAT(5HORDER,I3//5X,4HREAL,11X,9HIMAGINARY)  
07712 C INPUT AND PUNCH MATRIX

PAGE 02

```
07712      DO 1 I = 1, N
07724      DO 1 J = 1, N
07736      READ 100, AR(I,J), AI(I,J)
07892 100  FORMAT (E14.8, E14.8)
07920      PUNCH 116, AR(I,J), AI(I,J)
08076 C      SET UP WORKING MATRIX
08076      WR(I,J) = AR(I,J)
08232      WI(I,J) = AI(I,J)
08388      1 CONTINUE
08460      PUNCH 103
08484 103  FORMAT (/9HCONSTANTS)
08532 C      INPUT AND PUNCH CONSTANTS
08532      DO 2 I = 1, N
08544      READ 100, FR(I), FI(I)
08628      2 PUNCH 116, FR(I), FI(I)
08748      MN = 0
08784      LIM = N-1
08832      50 SIGN = 1.0
08868 C      DIAGONALIZATION OF DETERMINANT
08868      DO 23 I = 1,LIM
08880      NUM = I
08916      L = I+1
08964      18 DEN = WR(NUM,I)*WR(NUM,I)+WI(NUM,I)*WI(NUM,I)
09300      IF(DEN) 14, 15, 14
09356      14 IF(NUM-I) 914, 914, 24
09424      15 NUM = NUM+1
```



PAGE 03

```
09472      IF (NUM-N) 18, 18, 53
09540      53 IF(MN) 4, 4, 5
09596      4 PUNCH 110
09620      PRINT 110
09644      110 FORMAT(//42HZERO DETERMINANT - USE DIFFERENT METHOD OF9H SOLUTION)
09788      STOP
09836      24 DO 16 J = 1, N
09848      WRT = WR(I,J)
09944      WIT = WI(I,J)
10040      WR(I,J) = WR(NUM,J)
10196      WI(I,J) = WI(NUM,J)
10352      WR(NUM,J) = WRT
10448      16 WI(NUM,J) = WIT
10580 C      CHANGE SIGN OF DETERMINANT IF ROWS ARE INTERCHANGED
10580      SIGN = -SIGN
10628      914 DO 23 J = L, N
10640      WRT = WR(J,I)
10736      WIT = WI(J,I)
10832      DO 23 K = I, N
10844      X1 = WRT*WR(I,K)-WI(I,K)*WIT
11072      X2 = WR(I,K)*WIT+WI(I,K)*WRT
11288      WR(J,K) = WR(J,K)-(WR(I,I)*X1+WI(I,I)*X2)/DEN
11660      WI(J,K) = WI(J,K)-(WR(I,I)*X2-WI(I,I)*X1)/DEN
12020      23 CONTINUE
12128 C      ADJUST MAGNITUDE TO AVOID OVER OR UNDERFLOW
12128      5 IS = 0
```

```
12164      DO 200 I = 1, N
12176      220 ABWR = ABSF(WR(I,I))
12272      IF(ABWR) 200, 200, 213
12328      213 IF(ABWR-1.) 211, 200, 210
12396      211 IF(ABWR-.1) 212, 200, 200
12464      212 TTEN = 10.
12500      IS = IS-1
12548      GO TO 214
12556      210 TTEN = .1
12592      IS = IS+1.
12640      214 WR(I,I) = WR(I,I)*TTEN
12808      WI(I,I) = WI(I,I)*TTEN
12976      GO TO 220
12984      200 CONTINUE
13020 C      EVALUTION OF DETERMINANT TAKING PRODUCT OF DIAGONAL ELEMENTS
13020      DO 7 I = 2, N
13032      J = I-1
13080      PRODR = (WR(J,J)*WR(I,I)-WI(J,J)*WI(I,I))
13428      PRODI = (WR(I,I)*WI(J,J)+WI(I,I)*WR(J,J))
13764      WR(I,I) = PRODR
13860      7 WI(I,I) = PRODI
13992      PRODR = PRODR*SIGN
14040      PRODI = PRODI*SIGN
14088      DET = PRODR*PRODR+PRODI*PRODI
14184      IF (DET) 111, 121, 111
14240      121 IF(MN) 4, 4, 111
```

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14296 111 PUNCH 115  
14320 115 FORMAT (/2X, 10HREAL PROD., 7X, 10HIMAG PROD. 7X, 11HDETERMINANT)  
14478 PUNCH 116, PRODR, PRODI, DET, MN  
14538 116 FORMAT (E14.8, 3X, E14.8, 3X, E14.8, 15)  
14602 ISD = IS+IS  
14650 PUNCH 117, IS  
14674 117 FORMAT (10X,44HMULTIPLY REAL AND IMAGINARY PRODUCTS BY 1.0E, 15)  
14818 PUNCH 118, ISD  
14842 118 FORMAT (10X, 28HMULTIPLY DETERMINANT BY 1.0E, 15)  
14954 IF(MN) 8, 9, 8  
15010 C DETERMINANT OF THE COEFFICIENTS IS SAVED FOR LATER COMPUTATIONS  
15010 9 BOT = DET  
15046 ISZ = IS  
15082 PRDIZ = PRODI  
15118 PRDRZ = PRODR  
15154 PHID = ATAN<sup>2</sup>(PRODI/PRODR)\*CONV  
15226 AMAG = SQRTF(DET)  
15262 PUNCH 109, PHID  
15286 109 FORMAT(/13HPHASE ANGLE =, E14.8, 7HDEGREES)  
15368 PUNCH 125, AMAG, IS  
15404 125 FORMAT(11HMAGNITUDE =, E14.8, 7H \* 1.0E, 15//)  
15486 C SET UP DETERMINANTS WITH COEFFICIENTS OF UNKNOWNNS REPLACED BY  
15486 C KNOWN TERMS  
15486 DO 10 MN = 1, NSOL  
15498 DO 11 I = 1, N  
15510 DO 11 J = 1, N

```
15522      WR(I,J) = AR(I,J)
15678      11 WI(I,J) = AI(I,J)
15906      DO 12 J = 1, N
15918      WR(J,MN) = FR(J)
16038      12 WI(J,MN) = FI(J)
16194      GO TO 50
16202 C    SOLUTION OF THE UNKNOWNNS
16202 C    POWER OF 10 READJUSTS TO CORRECT MAGNITUDE
16202      8 VAL = (10.**((IS-ISZ))/BOT
16298      XR(MN) = (PRODR*PRDRZ+PRODI*PRDIZ)*VAL
16430      XI(MN) = (PRDRZ*PRODI-PRODR*PRDIZ)*VAL
16574      10 CONTINUE
16610      PUNCH 106
16634      106 FORMAT (/46HSOLUTIONS OF THE SIMULTANEOUS LINEAR EQUATIONS/)
16766      PUNCH 3, N
16790 3FORMAT(5HORDERI4//5X4HREAL11X9HIMAGINARY7X11HPHASE ANGLE7X9HMAGNITUDE/)
17026      DO 13 I = 1, NSUL
17038      IF(XR(I)) 122, 123, 122
17118      123 PHID = 90.
17154      GO TO 124
17162      122 PHID = ATANF(XI(I)/XR(I))*CONV
17282      124 AMAG = SQRTF(XR(I)*XR(I)+XI(I)*XI(I))
17486      PUNCH 120, XR(I), XI(I), PHID, AMAG
17594      120 FORMAT(E14.8, 3X, E14.8, 3X, E14.8, 3X, E14.8)
17670      13 CONTINUE
17706      PRINT 900
17730      900 FORMAT(31HPAUSE, PUSH START FOR NEXT CASE)
```

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17816 PAUSE  
17828 GO TO 500  
17836 END

SYMBOL TABLE

39999 SIN  
39989 SINF  
39979 COS  
39969 COSF  
39959 ATAN  
39949 ATANF  
39939 EXP  
39929 EXPF  
39919 LOG  
39909 LOGF  
39899 SQRT  
39889 SQRTF  
39879 ABSF  
39869 ABSFF  
39859 AR 35869  
35859 AI 31869  
31859 FR 31669  
31659 FI 31469  
31459 WR 27469  
27459 WI 23469  
23459 XR 23269  
23259 XI 23069  
23059 CONV  
23049 18000000+03  
23039 31415926+01  
23029 000  
23019 \*0500  
23009 \*0101  
22999 \*0101  
22989 NSOL  
22979 N  
22969 \*0104  
22959 \*0104  
22949 \*0105  
22939 \*0105  
22929 \*0119  
22919 \*0119  
22909 \*0001  
22899 I  
22889 J  
22879 \*0100  
22869 \*0100  
22859 \*0116  
22849 \*0116

22839 \*0103  
22829 \*0103  
22819 \*0002  
22809 MN  
22799 0000  
22789 LIM  
22779 0001  
22769 \*0050  
22759 SIGN  
22749 10000000+01  
22739 \*0023  
22729 NUM  
22719 L  
22709 \*0018  
22699 DEN  
22689 001  
22679 \*0014  
22669 \*0015  
22659 \*0914  
22649 \*0024  
22639 \*0053  
22629 \*0004  
22619 \*0005  
22609 \*0110  
22599 \*0110  
22589 \*0016  
22579 WRT  
22569 WIT  
22559 K  
22549 X1  
22539 X2  
22529 002  
22519 IS  
22509 \*0200  
22499 \*0220  
22489 ABWR  
22479 \*0213  
22469 \*0211  
22459 \*0210  
22449 10000000+00  
22439 \*0212  
22429 TTEN  
22419 10000000+02  
22409 \*0214  
22399 \*0007  
22389 PRUDR  
22379 PRODI  
22369 DET  
22359 \*0111  
22349 \*0121  
22339 \*0115  
22329 \*0115  
22319 ISD

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22309 \*0117  
22299 \*0117  
22289 \*0118  
22279 \*0118  
22269 \*0008  
22259 \*0009  
22249 BOT  
22239 ISZ  
22229 PRDIZ  
22219 PRDRZ  
22209 PHID  
22199 AMAG  
22189 \*0109  
22179 \*0109  
22169 \*0125  
22159 \*0125  
22149 \*0010  
22139 \*0011  
22129 \*0012  
22119 VAL  
22109 \*0106  
22099 \*0106  
22089 \*0003  
22079 \*0003  
22069 \*0013  
22059 \*0122  
22049 \*0123  
22039 90000000+02  
22029 \*0124  
22019 \*0120  
22009 \*0120  
21999 \*0900  
21989 \*0900

APPLICATIONS OF NUMERICAL FILTERS IN  
THE POWER SPECTRAL ANALYSIS OF STATIONARY  
TIME SERIES

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We will focus our attention on the spectral analysis of finite length recordings of a physical process which is assumed to be random in nature. For deterministic functions such as periodic and aperiodic functions a harmonic analysis is usually carried out by Fourier series analysis and by Fourier integral analysis, respectively. The discrete line spectrum for a periodic function and the continuous spectrum for the aperiodic function may be determined analytically because these deterministic functions are "known for all values of time". Random series are a class of functions which are not deterministic and do not lend themselves to the same harmonic analysis techniques used for deterministic functions; that is, statistical methods must be used.

The Tukey technique, which is used here, is applicable to random time series which very closely approximate a

stationary random ergodic process. This computational procedure yields the variance spectrum of a time series.

Other names for the resultant computation are power density spectrum, second-degree spectrum, or quadratic spectrum; all of which refer to the distribution of variance as a function of frequency.

One begins with a recording of a physical process which is assumed to represent a sample of a random process.

The record must be free of "pure tone" or periodic components and transients. After sampling the record at equi-spaced intervals the linear trends and average should be removed.

Briefly, the Tukey method consists of computation of statistical estimates of the spectrum of a finite discrete time series by a numerical approximation of the Wiener-Khinchine equations. The procedure involves two

steps. First, one computes a set of mean lagged products of the time series. Another name for the set of mean lagged products is the autocorrelation function. The raw power spectral estimates are computed by application of a discrete finite Fourier cosine transform to the autocorrelation function. This transformation gives the desired frequency domain representation of the time series. Systematic statistical errors resulting from use of a finite amount of data appear in the raw power spectral estimates. The Tukey technique to obtain improved spectral estimates involves a smoothing or refining operation performed on the raw estimates.

Slide 1 shows the Tukey equations.

Slide 2 shows an example of a time series to which one might apply the Tukey analysis.

Slide 3 shows the power density spectrum of the time

series. Eighty percent confidence intervals are shown.

Your attention is directed to the fact that the power density graph has an upper bound at a point marked  $f_N$  and that no power estimates of higher frequency are plotted. This upper band set is known as the Nyquist frequency and is a function of the length of the sampling interval. A full discussion of sampling theory is beyond the scope of this presentation. However, a few brief remarks are in order.

When a continuous function is sampled at equi-spaced intervals, the question should be asked: "How well will the discrete set of sampled values represent the original function?" A continuous function of time is completely determined by its values at equally spaced intervals provided that the continuous function contains no frequencies higher than, say,  $W$  cycles per second, and

the ordinates are given at points spaced  $1/2 W$  seconds apart, the series extending for all time. This is a statement of the popularly referred to Shannon theorem.

Under consideration here is an analysis which is to be based on sampled values obtained from continuous records which are not infinite in extent and are not band limited. Analysis based on finite amounts of data is common to statistical work.

Of immediate concern is the selection of the sampling interval and the problem of aliasing. Consider two sine waves of equal amplitude, but different frequencies.

(See Slide 4)

Attention here is directed to a particular set of sine waves, differing in frequency, but having a common set of equally spaced sample values. Thus, given only the sampled values, a sine wave of a given frequency may

be confused with a sine wave of higher frequency.

Specifically, if a harmonic time function  $X(t)$  is sampled at equally spaced time intervals  $\Delta t$ , then a

frequency  $f_N = \frac{1}{2\Delta t}$

called the Nyquist or folding frequency, exists such that the functions with frequencies

$$f \pm nf_N, \text{ for } n = 0, 2, 4, \dots,$$

are not distinguishable.

Obviously, then, power contributed to a power spectrum at a given frequency  $f$  cannot be distinguished from powers contributed by frequencies  $f \pm n f_N$ . This translation of frequencies is known as aliasing. If the data actually contain power at frequencies greater than  $f_N$ , this power will be "folded back" into the principal band which extends from 0 to  $f_N$ . Power that is folded back results in a distortion of the true power spectrum in

the principal band.

To make the effect of aliasing negligible it is necessary to select a sampling interval "small enough" to place the Nyquist frequency beyond all significant power contributions.

Associated with each spectral estimate there is a confidence interval which depends on the number of degrees of freedom in the computation. If one assumes the distribution of the data to be Gaussian and that the distribution of the variability in the spectral estimates follows the so called "chi-square" distribution, then the number of degrees of freedom may be computed by the convenient formula:

$$k = \frac{2}{m} (N - \frac{m}{3})$$

where k = number of degrees  
N = number of sampled values  
m = number of the maximum log

The confidence intervals are then computed using the number of degrees of freedom. As the number of degrees of freedom is increased the confidence intervals decrease in size and the computed estimates are more reliable. The number of degrees of freedom is, generally speaking, directly proportional to the number of data points and inversely proportional to the maximum number of lags. Acquisition of more data may be impossible or economically unfeasible and reducing the number of lags reduces the number of spectral points in the frequency range from zero to the Nyquist frequency.

This brings us to the point of this paper.

In many physical processes the power density decreases very rapidly with increasing frequency. Often at the higher frequencies the power density of the process under investigation is of the same order of magnitude as the



noise background. One must sample the processes often enough to avoid aliasing which would cause the noise to "fold back" into the frequency range of interest. Then one must take many lags and compute many power density estimates in order to have a good look at the lower frequencies. The consequences of this are large confidence intervals and much computation.

In order to get around this problem one can operate on the original sampled data with a linear operator which is often called a numerical filter because of its mathematical resemblance to an electrical filter. Through use of filters one can change the frequency spectrum in a known and desirable way. In particular, a low-pass filter may be used to suppress the power near the Nyquist frequency and not significantly disturb the low frequency spectrum of a time series.

Slide 5 shows a power density spectrum computed before and after low-pass filtering.

Slide 6 shows a comparison between the mathematical model of an electrical filter which operates on a continuous electrical signal and a linear operator (a numerical filter) which operates on a set of equispaced sample values of a time series. Note that the time domain representation of the electrical filter is characterized by  $W$ , the impulse response or memory of the filter. The time domain representation of the linear operator is simply an array of numbers. In the frequency domain both the electrical and numerical filters have representations called the frequency response. It can be shown that the numerical filter is simply a numerical approximation to the mathematical model of the electrical filter.

Slide 7 shows a plot of the coefficients of a low-pass filter.

Slide 8 shows the frequency response of both a high-pass and a low-pass filter.

After the time series has been operated on by say, a low-pass filter, the new time series may be resampled using a larger sampling interval. That is, the set of sampled values may be decimated by taking every other value, every third value, etc. A new lower Nyquist frequency is associated with the power spectrum of the new time series since the new sampling interval is larger than the original one. The low-pass filter has suppressed the power at the higher frequencies and thus all but eliminated possible distortion caused by aliasing. Now the low frequency range may be investigated using fewer lags and thus keep the size of the confidence intervals small.

After the power spectrum has been computed the effect of the filter is removed using the frequency domain representation of the filter.

In various applications high-pass, band-pass as well as low-pass filters have been used. Such computations are used in geophysical applications such as analysis of temporal variations in the earth's magnetic field and in biomedical applications such as analysis of EEG recordings.

Slide 9 shows a macro-flow chart of a computer program, written in 1620 Fortran II, to accomplish the computations discussed in this presentation.

Listings of the program are available from the author. (User 5130).

# TUKEY EQUATIONS

## AUTOCORRELATION

$$C_r = \frac{1}{N-r} \sum_{q=0}^{N-r} X_q \cdot X_{q+r}, \quad r = 0, 1, 2, \dots, m < N$$

## RAW POWER DENSITY

$$V_r = \left[ C_0 + 2 \sum_{q=1}^{m-1} C_q \cos\left(\frac{qr\pi}{m}\right) + C_m \cdot \cos(r\pi) \right] \cdot \Delta t$$

## REFINED POWER DENSITY

$$P_0 = \frac{1}{2} (V_0 + V_1)$$

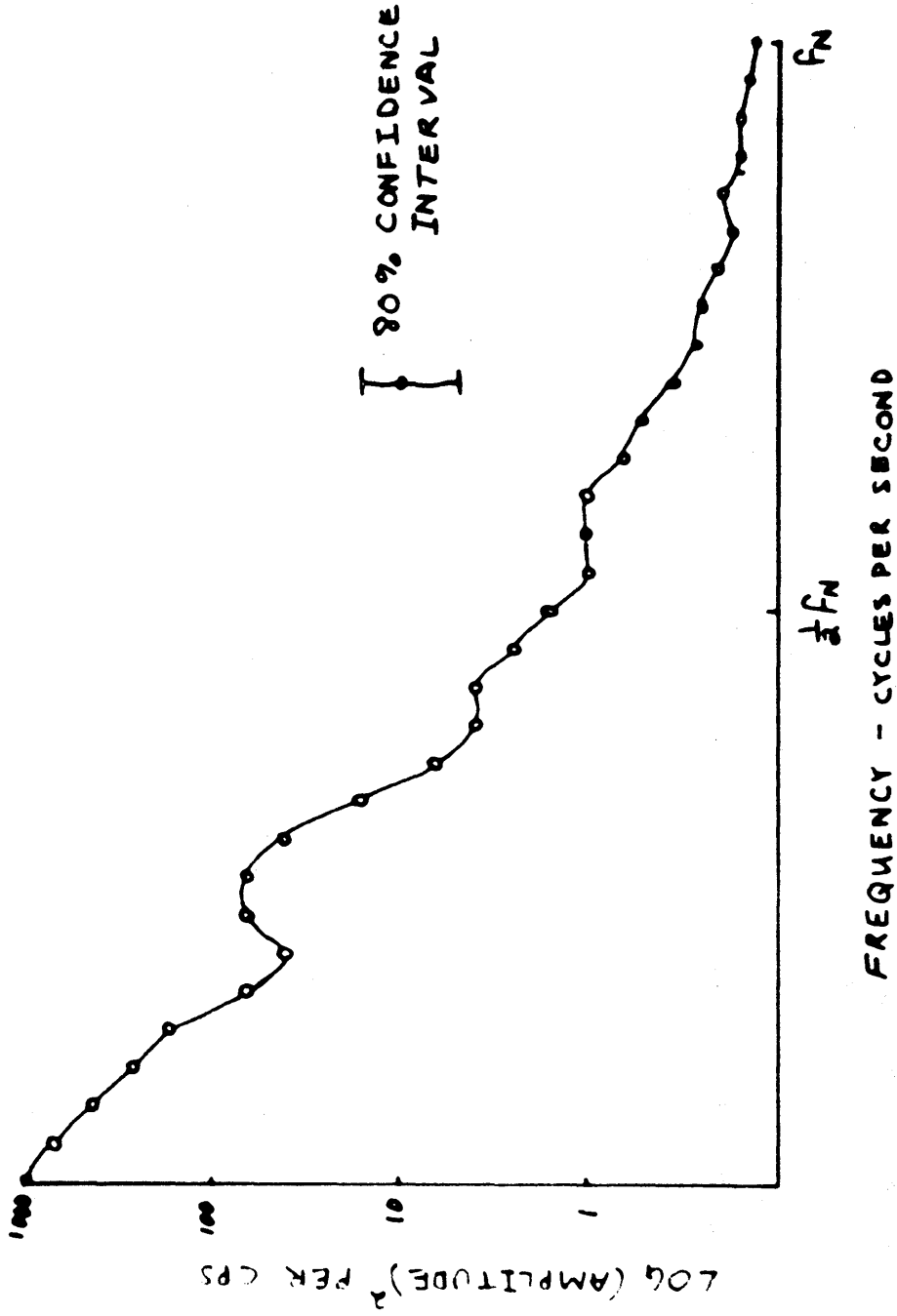
$$P_r = \frac{1}{4} V_{r-1} + \frac{1}{2} V_r + \frac{1}{4} V_{r+1}$$

$$P_m = \frac{1}{2} (V_{m-1} + V_m)$$

$$f_r = \frac{r}{2m\Delta t}$$



# POWER DENSITY SPECTRUM



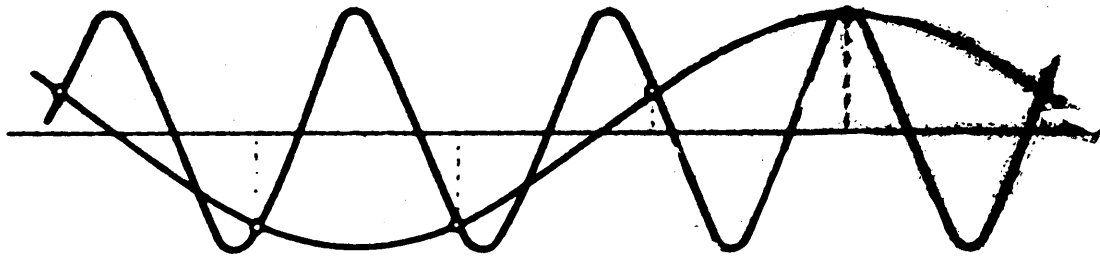
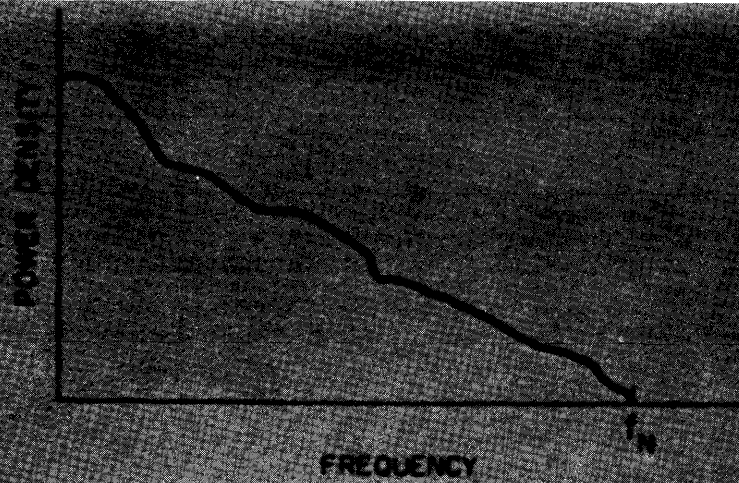


FIGURE 9  
SINE WAVES OF DIFFERENT FREQUENCIES WITH THE  
SAME SET OF EQUALLY SPACED SAMPLE VALUES

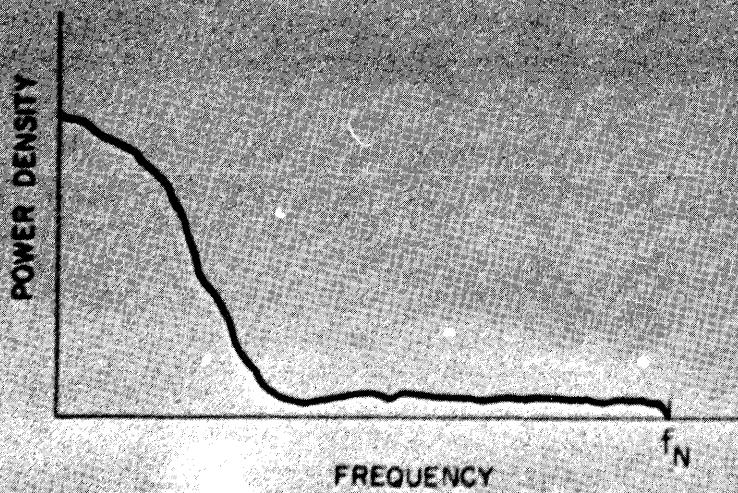
107

SLIDE 4





POWER SPECTRUM BEFORE LOW-PASS FILTERING



POWER SPECTRUM AFTER LOW-PASS FILTERING

FIGURE-10  
EXAMPLE OF A POWER SPECTRUM BEFORE AND AFTER  
LOW-PASS FILTERING

## ELECTRICAL FILTER



$$H(t) = \int_0^{\infty} G(t-\tau) W(\tau) d\tau$$

$$G(t, f) = e^{j2\pi f t}$$

$$H(t) = \int_0^{\infty} e^{j2\pi f(t-\tau)} W(\tau) d\tau$$

$$H(t) = e^{j2\pi f t} \int_0^{\infty} e^{-j2\pi f \tau} d\tau$$

$$H(t, f) = Y(f) G(t, f)$$

## NUMERICAL ANALOGUE

$$a_k = W(k\Delta t)$$

$$v(t) = \sum_{k=0}^k a_k u(t - k\Delta t)$$

$$u(t, f) = e^{j2\pi f t}$$

$$v(t) = \sum_{k=0}^k a_k e^{j2\pi f(t - k\Delta t)}$$

$$v(t) = e^{j2\pi f t} \sum_{k=0}^k a_k e^{-j2\pi f k \Delta t}$$

$$v(t, f) = u(t, f) \cdot Y(f)$$

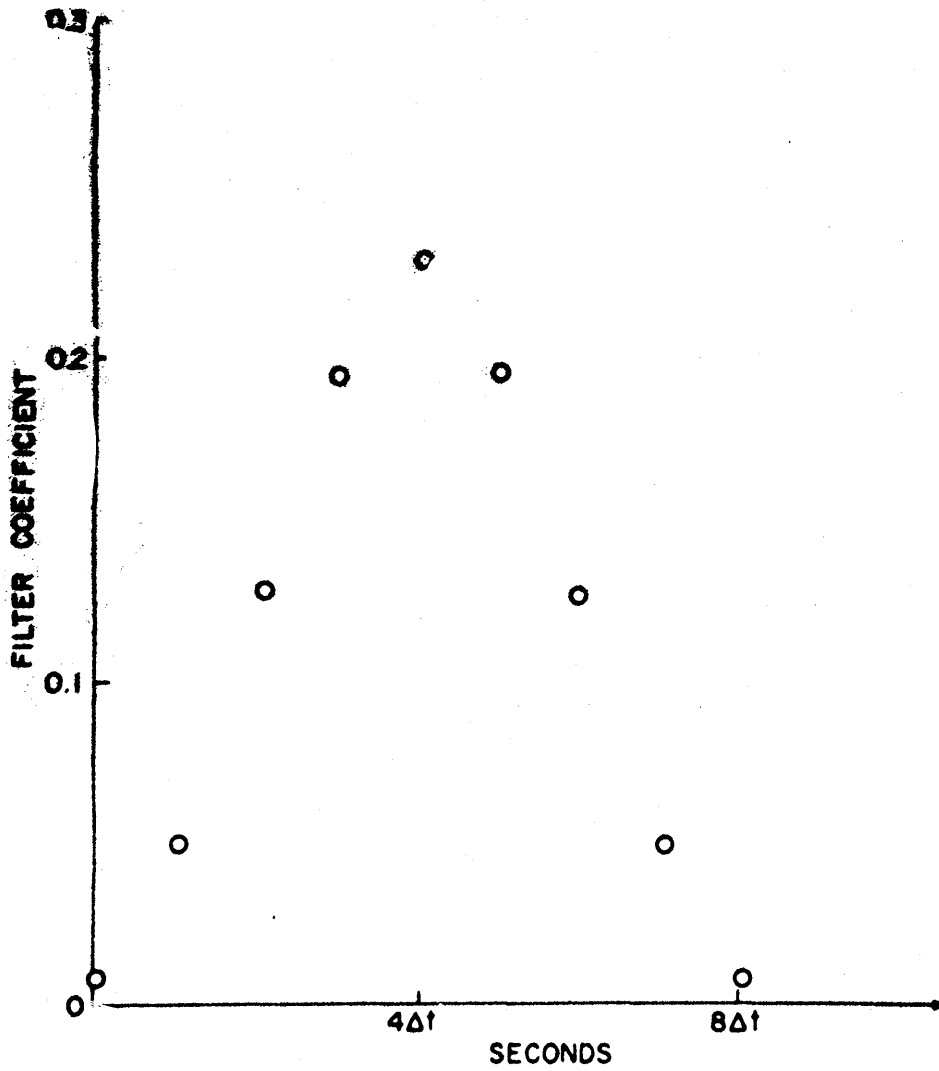


FIGURE 13  
 THE FILTER COEFFICIENTS OF THE LOW-PASS FILTER,  
 CASE B OF TABLE I

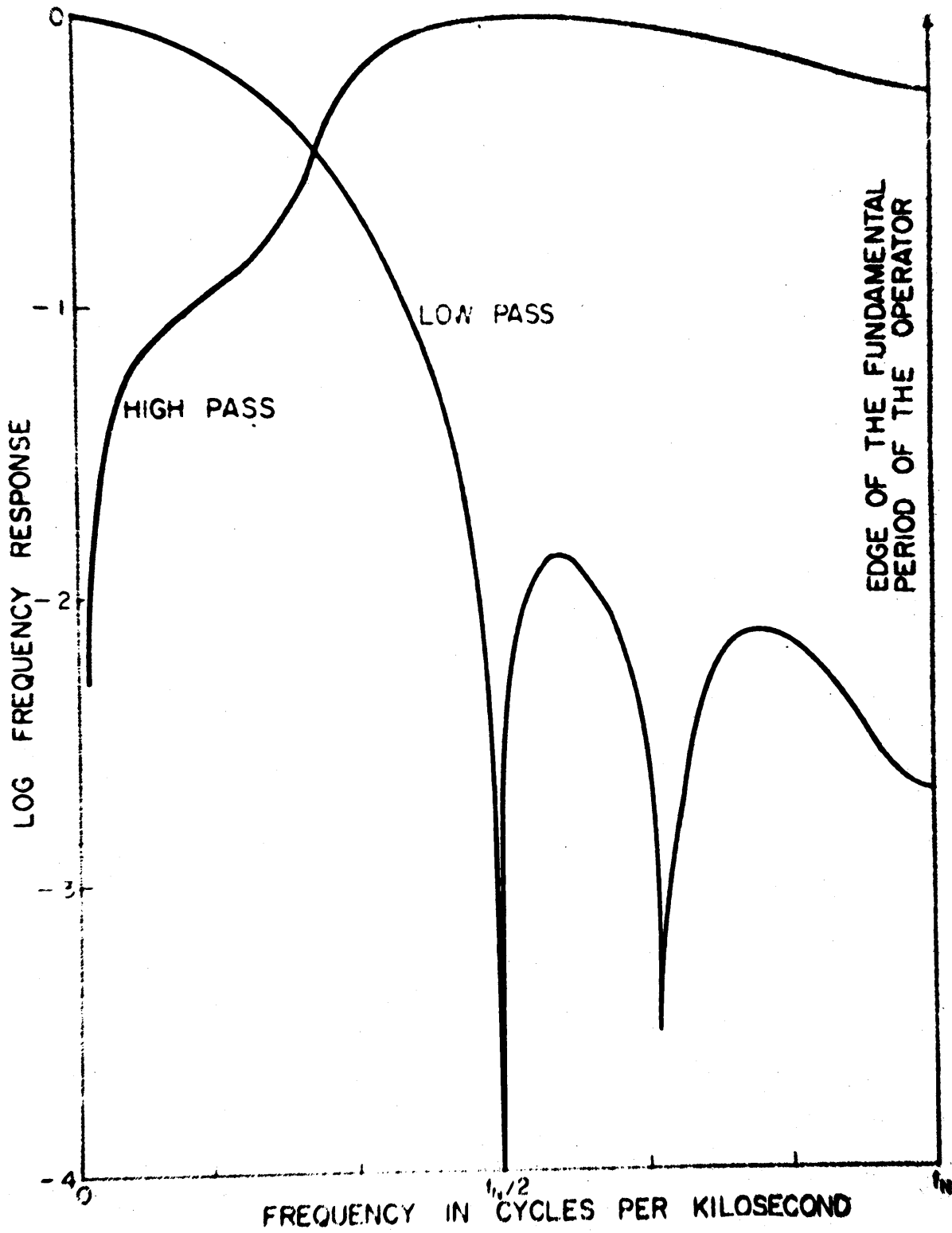


FIGURE 14  
 THE FREQUENCY RESPONSE OF LOW-PASS  
 AND HIGH-PASS FILTERS, TABLE I.

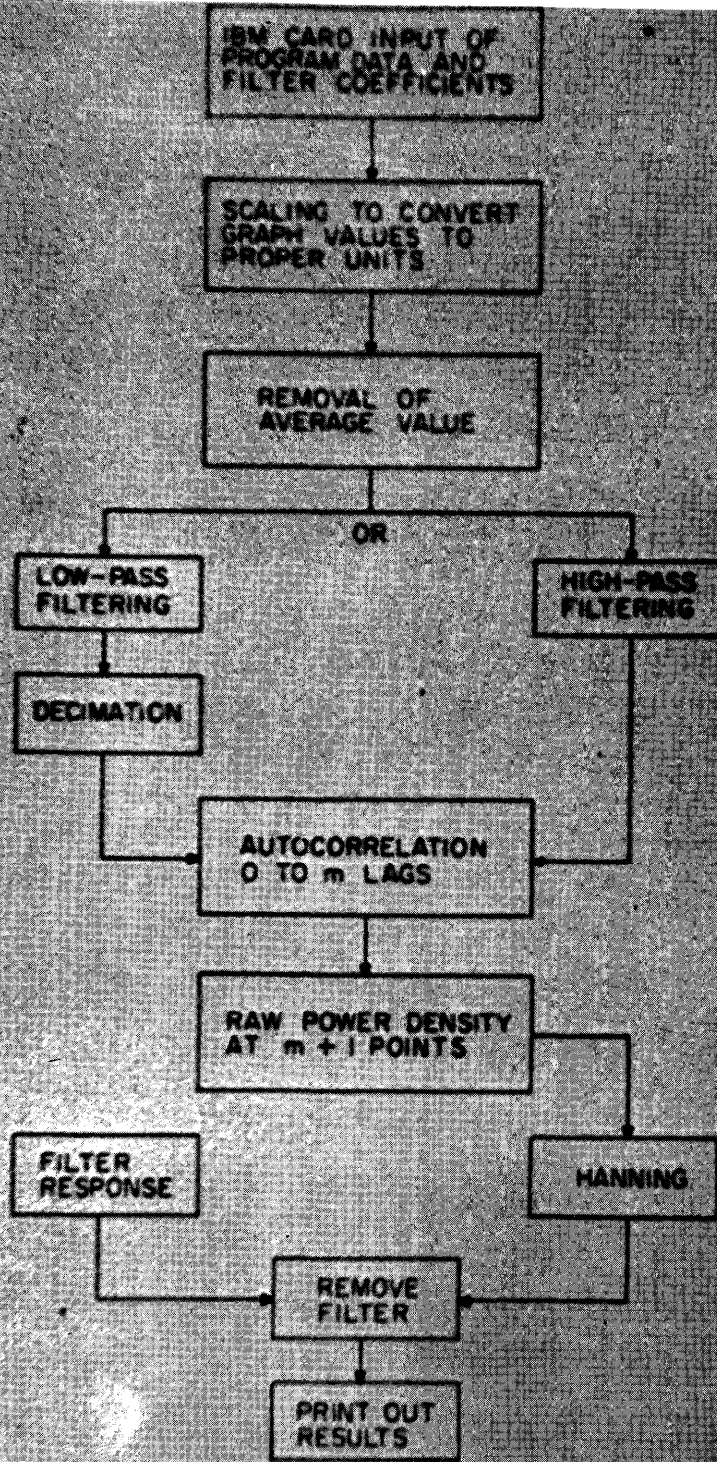


FIGURE 8  
 MACHINE COMPUTATION FLOW CHART FOR  
 A SINGLE TIME SERIES

SLIDE 9

IBM 1620 ASSISTS STUDENT COUNSELORS AT JUNIOR COLLEGE

Paul S. Chan  
IBM CORPORATION  
3610 - 14th Street  
Riverside, California

May 18, 1964

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  - (b) Scattergram of SCAT T vs. Chemistry 1A
  - (c) Scattergram of Mathematics Placement vs. Chemistry 1A
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## ABSTRACT

### IBM 1620 ASSISTS STUDENT COUNSELORS AT JUNIOR COLLEGE

The present study reports findings, based on the computed results from the IBM 1620, concerning the extent to which test scores on the college freshman testing program - such as the ACE, SCAT, Co-operative English Tests - are able to predict academic success or failure in specific junior college courses. Scattergrams have been created for those correlations of highest significance to assist counselors in estimating the incoming student's aptitude for college level study and in making a more accurate appraisal of the student's competence in a particular subject area.

Paul S. Chan  
May 15, 1964



## INTRODUCTION

Unlike private colleges, the state colleges, or the state university, California's public junior colleges are required by law to admit any resident of their districts who is a high school graduate or who is over 18 and able to profit from instruction.

Junior college administrators have interpreted this as meaning that they cannot deny admission to any applicant who has reached his 18th birthday, although virtually all now have retention policies which deny re-enrollment to students who fail to maintain a "satisfactory" grade point average. At one time, many administrators interpreted the legislative mandate to mean that they could not set any qualification for registration in any class. An apparent change in legislative sentiment has combined with the realities of post-war enrollment pressures to cause most junior colleges to search for some equitable means of screening from classes (particularly from transfer classes) those students who have little opportunity to succeed.

The freshman testing program has been an established practice at Riverside City College, a public junior college, for the past years. Although the counselors and admissions officers have been making extensive use of these tests to assist in laying out the academic path of many students, there have been no attempts until recently to make regular evaluations of the measuring instruments in use. Recently an IBM 1620 was installed at the college. One of the first projects to use the system was an attempt to determine the relationship between the test scores and the final grades in specific courses. It is anticipated that the results will improve placement of students in appropriate sections or courses, and selection of students for particular areas of concentration or pre-professional training.

## PURPOSE OF THE STUDY

The battery of tests - ACE, SCAT and others - were administered to the in-coming new students for the dual purpose of counseling and placement. Since this investigation was the initial application, the present study was to demonstrate the validity of the battery for these purposes.

Another aim of the study was to modify the battery to include only those tests best suited for the screening program. Excessive overlap of abilities measured by one test and those measured by another results in a waste of the student's time. Also, too great an array of scores for academic counselors might prove more confusing than helpful.

It was anticipated, too, that critical cut-off scores could be developed for each test, making it both practical and possible to advise the individual student, upon the basis of his score, just what his chances for success of failure in a specific course would be.

## DATA

This study involved over 800 students who were enrolled in Psychology 49, a freshman orientation course, and who had completed one or more of 25 courses which the college wished to examine.

There were fifteen predictors. These included:

- (1) three scores from the ACE (Quantitative, Linguistic, and Total)
- (2) the R.C.C. Arithmetic Competency Test of 40 items
- (3) three scores from the School and College Ability Tests (SCAT, Quantitative, Verbal, and Total)
- (4) six scores from the Cooperative English Tests, Form 1A-1960 EDITION (Vocabulary, Level of Comprehension, Speed of Comprehension, Total Reading, English Expression, and Total English)
- (5) overall high school grade point averages (to obtain this figure academic subjects and others such as typing, speech, journalism, and music courses were used. Physical education, military science and driver education were not used. Shop courses were used where it was the student's high school major.)
- (6) academic grade point averages (to obtain this figure only solids such as English, foreign languages, math at the algebra and higher level, history and sciences, but not including general science, were used.)

ACE and Arithmetic scores were easily obtained because they are a part of the placement battery of tests required of all new students. The SCAT and Cooperative English test scores were obtained by testing in the Psychology 49 classes and the two high school grade point averages were rather tediously obtained by employing an individual to compute the figures by hand.

The courses included chiefly transfer courses with a few not-transfer type courses and represented a cross-section of the major divisions within the college.

<u>DEPARTMENT</u>	<u>Course No.</u>	<u>Descriptive Title</u>
Anthropology	2	Cultural Anthropology
Art	1A	History and Appreciation of Art
Biology	1	General Biology
Business	1A	Principals of Accounting
Business	18A (hour)	Business Law
Business	50A (51A)	Elementary Accounting
Business	81A (50A)	Business Mathematics
Chemistry	1A	Chemistry
Chemistry	2	Introductory General Chemistry

<u>DEPARTMENT</u>	<u>Course No.</u>	<u>Descriptive Title</u>
Electronics	51	Electrical Fundamentals of Electronics
English	1A	English Composition
Geography	1	Introductory Physical Geography
History	3	American History
History	4A	History of European Civilization
History	6A	Political and Social History of the US
Math	3A	Analytic Geometry and Calculus
Music	20	History and Appreciation of Music
Nursing	1A	Introduction to Nursing
Philosophy	6A	Introductory Philosophy
Physical Science	1	Introduction to Physical Science
Physics	2A	General Physics
Political Science	3	American Political Institutions
Psychology	1A	General Psychology
Sociology	1	Introduction to Sociology
Spanish	1	Elementary Spanish

This battery of tests was originally selected to provide a basis for predicting over-all scholastic success and success in specific subject-matter areas: The ACE for general scholarship, with its Q and L sub-scores for areas of primarily quantitative and verbal content respectively; the Co-operative English Test for English and other areas which require considerable reading; the Mathematics Tests for placement in mathematics and allied physical science courses.

Final grades of the students in each of the chosen freshman courses were compared with their scores on each of the tests, The courses were chosen from four areas: Language, Humanities, Social Sciences, and Natural Sciences.

## METHOD OF ANALYSIS

Results of the test battery were separated into ten test variables:

- (1) Three scores, the Q, L, and T, were derived from the ACE;
- (2) Six scores from the Co-operative English Test;
- (3) Three scores, the V, Q, and T, from the SCAT; and
- (4) One score from Mathematics Placement Test.

All the test scores and course grades were recorded in punched cards.

An analysis program was written in Fortran.

Coefficients of correlation were computed by the 1620 between scores on each of these tests and final grades in each course. To substantiate the validity of the results, besides the correlation coefficient, regression line coefficients, standard error of estimate, and standard error of regression coefficient  $b$ , the significance of  $r$  and of  $b$  were analyzed. A summary of equations for these calculations can be found in Appendix D.

## CORRELATIONS BETWEEN TEST SCORES AND FINAL GRADES

	cases	MATH			CO-OP ENGLISH						SCAT			
		Q	L	T	1 (Vo)	2 (Le)	3 (Sp)	4 (Tr)	5 (Exp)	6 (Tot Eng)	V	Q	T	
<u>Language</u>														
Spanish	33	064	044	295	297	268	326	196	233	179	232	281	161	362"
<u>Humanities</u>														
Anthropology	22	131	287	397	290	089	122	018	040	228	182	478*	049	159
Art 1A	18	488*	009	491"	362	656"	287	100	546"	794*	737*	382	388	519"
Music 20	22	463"	233	363	363	286	458*	404	389	404	456*	396	330	018
Philosophy 6A	24	415	117	196	084	596*	457"	051	027	154	091	065	045	501"
average		399	159	362	275	407	441	144	251	495	367	330	203	299
<u>Social Science</u>														
Geography	18	201	169	488"	438	525"	612"	688*	675"	336	548"	616*	093	242
<u>Natural Science</u>														
Biology	40	125	436*	326"	433*	433*	334"	449*	379"	547*	536"	391"	375"	375"
Chemistry 1A	55	442*	185	165	189	160	033	048	106	146	116	208	329"	266"
Chemistry 2	30	599*	217	406"	386"	416"	316"	421"	473*	551"	598*	411"	405"	124
Electronics	13	101	400	348	359	439	421	405	425	501	492	392	588"	484
Mathematics 3A	23	578*	121	245	144	360	357	324	368	305	324	413	225	126
Nursing 1A	25	127	209	376	144	596*	457"	310	503"	294	452"	579*	558*	384"
Physical Science	21	534"	342	306	450"	133	582"	449	377	528"	538"	065	491	328
average		278	218	241	218	326	290	260	295	320	342	275	385	232

\* Indicates .01 level of significance

" Indicates .05 level of significance

## RESULTS

The results of this study are reported in Appendix A, a table presenting the correlations between test results and course grades. Within each curricular area, the average correlation with each test is also given. All the correlations coefficients in Appendix A at the .01 level of significance are marked with an asterisk and at the .05 level with double primes.

Grades in some courses appear to correlate relatively well with scores on all the tests, while those in other courses showed low correlations with most of the test scores. For example, biology has 12 out of 13 sub-scores with correlation at either .05 or .01 level of significance and chemistry has ten out of 13, whereas Spanish and electronics have only one out of 13 at .05 level of significance. Some explanations may be offered for this phenomenon. One is that the differential magnitude of the correlations depends partly on the magnitude of the reliabilities of the grades in those courses. Sectionings of a course will certainly be a factor to affect the magnitude. Another factor is that grades in some courses are based on objective-type examinations, while in others on a more subjective basis.

The relatively high predictive power of the mathematics placement test in the Natural Science Division is more or less expected. However, an almost equivalent result was found in the Q part of the SCAT Test. This is an indication that it may be possible to obtain the same predictive information from either of the tests, so duplication of student's effort can be avoided. It is quite unexpected that Spanish correlates with only the total score of the SCAT Test in the entire battery. Also, electronics correlates only at .05 level of significance, with Q part of the SCAT Test. It is possible that this phenomenon is due to the fact that SCAT Tests involves not only the psychological functions commonly measured by tests of verbal ability, but also a particular type of reasoning ability important in academic success which is not assessed by any other tests employed in the present battery.

The two parts, speed of comprehension and total reading, of the Co-operative English Test show high correlation with geography. This can be explained because of the fact that the Social Studies courses normally require more speed in reading and in comprehension. The significant correlation at .01 level between philosophy and the vocabulary part of the Co-operative English Test certainly implies the requirements to succeed in the course.

In general, the six parts of the Co-operative English Test correlate relatively better than any of the three parts in the ACE Tests, with all the selected courses. This is illustrated by the r-values of .794 with Art, .458 with Music, etc.

The tendency was noted also for correlations to be relatively high or low with reference to separate courses rather than to the different tests. It was hypothesized that this phenomenon might be the result of difference among the courses in inter-section standardization reliability of grading, or use of objective examinations.



## SUMMARY

In this study of the value of a battery of aptitude and achievement tests for the prediction of junior college freshman grades, test scores were correlated with final grades in a variety of freshman courses. The individual correlations appeared small, but the relative predictive power was demonstrated clearly.

The following major conclusions concerning the predictive significance of the present battery appear to be warranted:

- (a) Overlapping of tests in the battery used is evidenced, suggesting that such an extensive array of examinations is somewhat superfluous and repetitive. Both over-all and individual course predictions could be made with even greater accuracy with a more abbreviated battery.
- (b) From the scattergram, it was found that it is feasible to determine the cut-off score in screening and to obtain more insight in the statistical probability of achievement of a student in a particular course.
- (c) Because of the small number of cases in this particular study, a caution against placing too much weight on individual test scores in guidance, selection or placement is in order.

SCATTERGRAM OF SCAT T-SCORE vs. CHEMISTRY 1A

	<u>F</u>	<u>D</u>	<u>C</u>	<u>B</u>	<u>A</u>
99-95			3		1
94-90		2	3		
89-85	1	2	2	2	
84-80	2	3	5	1	
79-75	2	3	1		1
74-70	1	3	4	1	
69-65	4	1			
64-60	3				
59-55			1		
54-50	1				
49-44		1			1
<hr/>					
TOTAL	14	14	19	4	3

SCATTERGRAM OF MATHEMATICS PLACEMENT  
vs.

CHEMISTRY 1A

	<u>F</u>	<u>D</u>	<u>C</u>	<u>B</u>	<u>A</u>
44-40				1	
39-35	3	3	11	3	2
34-30	9	7	7	1	
29-25	4	4	1		1
24-20	1				
<hr/>					
TOTAL	17	14	19	5	3

APPENDIX (D)

SUMMARY OF EQUATIONS

(1) Variances

$$S_x^2 = \frac{n\sum x^2 - (\sum x)^2}{n(n-1)}$$

$$S_y^2 = \frac{n\sum y^2 - (\sum y)^2}{n(n-1)}$$

(2) Regression Line

$$b = \frac{n\sum xy - \sum x \sum y}{n\sum x^2 - (\sum x)^2}$$

$$a = \frac{1}{n} (\sum y - b\sum x)$$

(3) Correlation Coefficient

$$r = \frac{\sqrt{[(n\sum xy) - \sum x \sum y]^2}}{\sqrt{[n\sum x^2 - (\sum x)^2][n\sum y^2 - (\sum y)^2]}}$$

(4) Standard Error of Estimate

$$S_{y/x} = \sqrt{\frac{n-1}{n-2} (S_y^2 - b^2 S_x^2)}$$

(5) Standard Error of Regression Coefficient b

$$S_b = \frac{S_{y/x}}{S_x \sqrt{n-1}}$$

(6) Significance of r

Compare  $|r|$  with the critical value in statistical table for 2 variables and n-2 degrees of freedom.

(7) Significance of b

Compare  $t = \frac{|b|}{S_b}$  with the critical value in statistical table for n-2 degrees of freedom.

**1620 COMPUTER UTILIZATION  
IN A WIND TUNNEL DATA  
ACQUISITION SYSTEM**

by

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1620 Users Group  
Brown Palace Hotel  
Denver, Colorado

18 June 1964

**NORTHROP CORPORATION**

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**ABSTRACT**

This paper describes how the IBM 1620 computer was teamed with a high-speed digital data acquisition system and two tape units to perform on-line processing of wind tunnel test data. The total installation is located in the Northrop Norair wind tunnel complex comprised of three tunnels: subsonic, transonic-supersonic, and hypersonic. The processing installation provides a central data acquisition and reduction function for all three tunnels, even simultaneously when necessary.

The high-speed data acquisition section scans, measures, and digitizes test data, introduces identification information, and records the data on magnetic tape for instantaneous reading by the 1620, in a read-after-write manner. The 1620 then reduces the data into tabulations meaningful to the aerodynamics research engineers, enabling them to make early evaluation of test run results and to proceed with model changes if called for.

During off-line operations, the computer is available for other applications, and has full control of the tape units.

## INTRODUCTION

Today, more than ever, competition in the aerospace industry is very keen and time is one of the most important elements to be utilized. For this reason, a company that makes use of wind tunnels must also have a satisfactory test data acquisition system and a means of automatically reducing the collected data as soon as it becomes available.

In the following paragraphs you will learn how we at Northrop have improved our techniques in this area. Our wind tunnels will be described as well as our data acquisition system to which a 1620 computer is coupled. Also of interest will be the changes we designed into the 1620 computer to make it suitable for our applications and the programs we have written to fulfill our objectives.

To the general public, wind tunnels are environmental chambers used to test model planes, but to the aerodynamicist, wind tunnels are probably the most superior devices used in aeronautical and aerospace research and development. Because of modern wind tunnels, today's test pilots are no longer the nerveless stunt men of the past, but professional engineers. Wind tunnels offer both fast and accurate data as well as the ability to simulate the different types of atmospheric conditions of any time of day or year. However, they are by no means new tools. Years before the Wright Brothers famed flight at Kitty Hawk, wind tunnels, crude as they were, gave valuable aerodynamics data which proved the feasibility of powered flight. The original wind tunnel employed by Orville and Wilbur Wright is on exhibit at the Air Force Museum, Wright-Patterson Air Force Base in Dayton, Ohio.



## WIND TUNNEL TESTING FACILITIES

Dominated by its 100,000-cubic-foot vacuum sphere is the supersonic-hypersonic wind tunnel facility at Hawthorne, California. This space age test facility provides test velocities from Mach 0.5 to Mach 14 with temperatures to 3000 degrees and simulated altitudes to 200,000 feet. To my knowledge, no privately-owned wind tunnel in the United States can produce the combined heat, pressure, velocity and run time that are obtainable with the one at Hawthorne. This relatively new, dual-circuit facility provides a greatly expanded capability for aerodynamics testing on advanced aircraft, missiles and space systems. It consists of two separate wind tunnel circuits: transonic-supersonic (Mach 0.5 through Mach 5) and hypersonic (Mach 6 through Mach 14). Design models can be tested for periods of at least 30 seconds in the supersonic circuit and up to one minute in the hypersonic circuit. The hypersonic tunnel can accommodate up to six 30-second runs each eight-hour shift. More test runs of proportionately shorter duration are possible.

Test sections, in which the models are mounted for aerodynamic study, measure two feet square in the supersonic circuit and 30 inches in diameter in the hypersonic circuit. A special "free jet" section in the hypersonic circuit allows removal of a model from the air flow while air flow is being established, thus protecting the model from excessive heat loads. The pressing of a button promptly injects the model into the flow stream. In a transonic or supersonic run, air passes from storage through a settling chamber (to smooth the airflow and remove any turbulence), is expanded through a nozzle (to establish Mach number), flows through the test section and then is forced through a "second throat" to reduce its velocity and to recompress it to atmospheric pressure before it exhausts through a muffler.

In a hypersonic run, air must be expanded so much (to achieve the higher velocities) that its temperature could actually be reduced to a point where the air would turn to liquid. To prevent liquefaction, an electrically fired heater containing a 16-ton bed of 3/8-inch alumina pebbles heats the air to temperatures as high as 3000 degrees Fahrenheit before it reaches the hypersonic nozzle. When the air is cooled by expansion, its temperature is therefore still high enough to keep it from liquefying.

From the test section, the hypersonic air passes through a "second throat" as in the supersonic circuit, to reduce velocity and then through a cooler to remove heat. It is then discharged into a large 100,000 cubic-foot vacuum sphere. The vacuum sphere is essential to hypersonic operations in order to achieve the high velocities desired in the test section. With storage pressure fixed at 3,200 pounds per square inch, the required pressure ratio obviously cannot be met by discharging the "used" air to atmospheric pressure (14.7 pounds per square inch). A low-pressure atmosphere is necessary and this is the function of the vacuum sphere.

About 100 feet from the supersonic-hypersonic facility and in another building is the 7' x 10' subsonic wind tunnel which went into operation in the year 1956 and was used in the very successful development of the Northrop T-38 Talon supersonic trainer, F-5 fighter, and Laminar Flow Control (LFC) airplane. During those tests, the output of test data was punched onto cards, carried to a remotely-located IBM 704 computer installation, processed and returned in a relatively long turn-around time (normally about three days; on emergency basis about four hours).

## DATA ACQUISITION SYSTEM

Today in the same building that houses the subsonic tunnel, is the data acquisition system, which we are very proud to possess. It was designed and built to our specifications by the Astrodata Corp. It serves all three of our wind tunnels. The data from any two of the three remotely-located tunnels can be transmitted to this center simultaneously.

The analog data, supplied by transducers at each of the tunnels is digitized by an analog-to-digital converter (ADC) in the central data system. The digital data from the ADC is then sent to the formatting generator where it is joined by other digital data from the model-position encoders, the time-of-day clock and also the switch settings from both the transmitting site and the central data system. The switch settings provide fixed information such as the barometric pressure, the test number, the run number and the date. The formatting generator then assembles and prepares the data for recording on magnetic tape. The records produced by the formatting generator are of variable length and automatically padded to contain an integral multiple of six characters, so that the resulting magnetic tape recordings can be used with both the 7090 and 1620 IBM computers. The ability to read the system-generated tapes by the 7090 computer proved very valuable during system checkout, because the 1620 computer was not adapted to handle magnetic tapes until later.

There are two types of records produced by the data acquisition system. The first of these is the title run record which identifies the test run by a test number, a run number, four parameters, the barometric pressure, the day and the time of day, and the model position by roll, yaw and pitch. The activation of the title push button switch will initiate output of a title run record consisting mainly of the above information provided through manually-set, thumbwheel switches. The second of the two types of records produced by this system is the data record. A data record is generated when the data circuit is closed (manually or automatically). The data record consists of an identification header, the time of day, the model position, and data from all site input channels programmed for the specific test.

## COMPUTER UTILIZATION

The 1620 computer employed is a Mod I with 40,000 core storage positions. It is equipped with most of the special, built-in features (indirect addressing, hardware divide, and floating point arithmetic). This computer is attached to the data acquisition system by an umbilical cord; it has been programmed to read and reduce the data as it is being recorded on any one of the two magnetic tape units. The reading is accomplished in a read-after-write manner, termed "eavesdropping." The information is introduced into the computer by the read gap, which is positioned a distance of .300-inch behind the write gap of a two-gap read-write head, almost immediately after the information is written onto the tape by the system. The normal function of the read gap, which is to provide parity checking during the recording process, was extended to make this possible. The two magnetic tape units used are Datamec D2020. These units are IBM compatible, using either 200 bpi or 556 bpi tape formats at 30 ips tape speed. The Central Data System (CDS) records at the 556 bpi density.

Eavesdropping allows the computer to sample the data as it is being recorded without interfering with the recording process itself. During the eavesdropping or on-line mode, as it is sometimes called, all the tape units are under the control of the CDS. Upon receipt of a signal from the 1620, the CDS causes the first character and associated parity bit to be transmitted to the 1620. Each character and associated parity bit continues to be transmitted until the longitudinal redundancy check character (LRCC) is encountered. The computer cannot initiate tape movement by attempting to read a tape while in this mode; therefore, a read tape instruction hangs up the computer until the CDS moves the tape to record new information. Besides the eavesdropping mode, the computer is also able to operate in an off-line mode. During the off-line mode, a selected tape unit (any of the two) may be read or written by the 1620 as if it were its own. These two modes of operation are manually selected.

Reduced punched card data is generally generated and plotted off-line during tests. An IBM 407 printer is also available in this center and is used to print much of the punched card output.

## COMPUTER HARDWARE MODIFICATIONS

The 1620 computer performing the data reduction is unique. Three new instructions had to be designed and the computer modified to permit their use for this special application. The design and implementation of these instructions into the computer required several months. In addition, other instructions were adapted to permit the reading, writing and other handling of magnetic tapes.

The three new instructions pertain specifically to the use of magnetic tape.

BST, backspace magnetic tape (36XXXXXX01300),  
REW, rewind magnetic tape (36XXXXXX02300) and  
WEF, write end of file (36XXXXXX01200)

Two instructions that refer to paper tape normally, RNPT, read numerically paper tape and WNPT, write numerically paper tape, were modified to read magnetic tape (RMT) and write magnetic tape (WMT), in the numerical mode.

RMT, read magnetic tape (36YYYYYY00300) and  
WMT, write magnetic tape (38YYYYYY00200)

In order to allow for tape redundancy and end of file testing, the functions of the following sense switch testing codes were extended.

BC1, branch console switch 1 on (46YYYYYY00100) and  
BC2, branch console switch 2 on (46YYYYYY00200)

When a BC1 instruction is executed, a branch takes place if either sense switch 1 is on or if a tape redundancy occurs. Likewise, the BC2 instruction also serves two purposes: a branch will occur if either sense switch 2 is on or an end of file mark is sensed. These two sense switches must be in their off position during magnetic tape operations. The redundancy and end of file indicators are not reset by any of these two instructions; they are reset only when the selected tape is put into motion again.

## SOFTWARE

Although the magnetic tapes normally may be read and written with FORTRAN coded programs by utilization of the paper tape statements, the tapes produced by this system can only be read by SPS or machine type programs. This is due to the various field widths contained within the records written by the system. The problem of reading tapes was quickly resolved by the writing of an SPS subprogram that could be called by and loaded with FORTRAN coded programs.

The SPS subprogram was designed to operate in two modes. The first of these modes, as directed by the arguments of the FORTRAN program, causes a compacted system record to be read from the tape. Each of the fields of the record is then extracted and expanded to a six-character field width. Flags are placed over the leftmost positions of each of the fields and the fields are then transmitted to their prescribed COMMON locations as integers. The second of the two modes requires the subprogram to search through the tape (not used during the on-line operation) for a particular title run record that agrees with the run and test numbers indicated in the arguments of the calling program. When the appropriate record is found, the information from the record is processed in the same manner as it was in the first mode. If the record is not found an indicator is placed in a communication field, reserved in COMMON for this purpose.

The two tape reading modes of the subprogram, described above, are very useful and make the 1620 an even more important asset to the overall system. The first of these two modes provides the user with integer data that is FORTRAN-compatible. The second mode, in addition to performing the same task as the first mode, assists in retrieving previously recorded data.

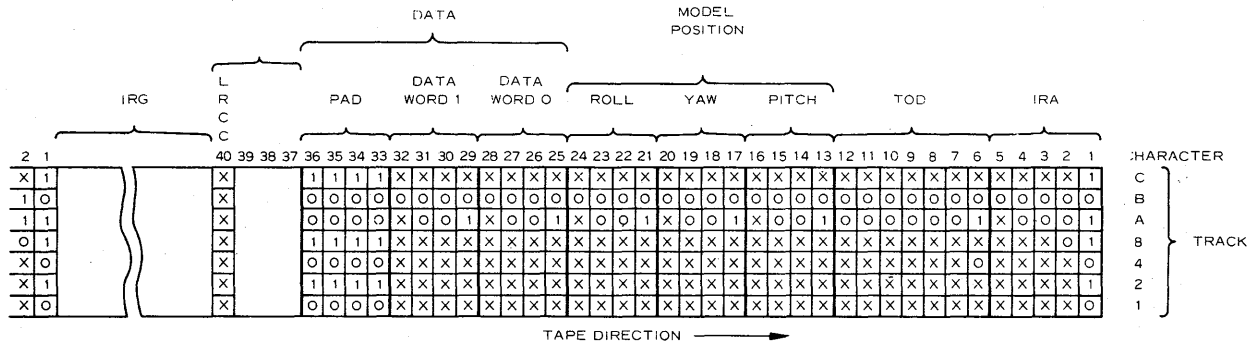
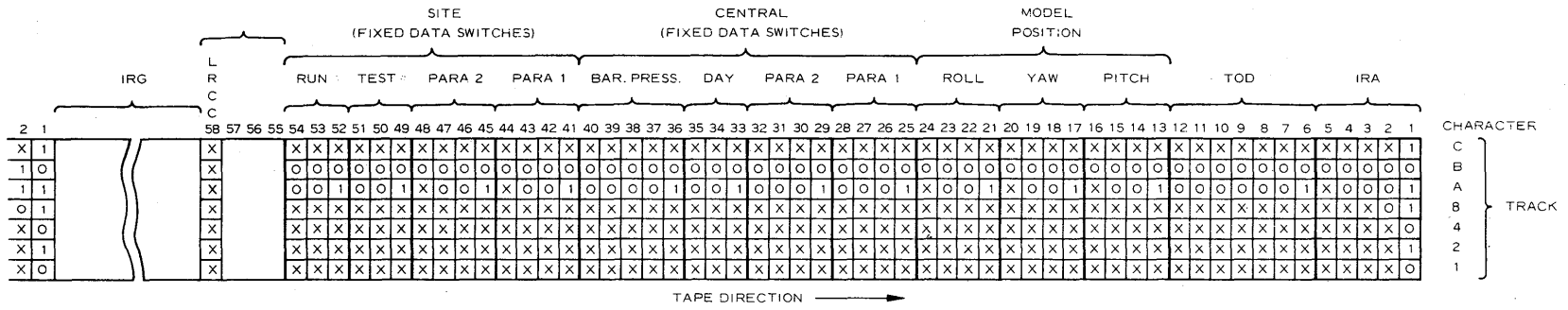
## RECORD FORMAT

The formatting generator produces two types of records and these are the title run record and the data record. The purpose of the title run record is to provide identification for the data records that follow it. The title run record consists of 54 characters of the following information:

Characters 1 - 5	Information Retrieval Aid (IRA)
Characters 6 - 12	Time of Day (TOD)
Characters 13 - 24	Model Position (Pitch, Yaw and Roll)
Characters 25 - 32	Parameters from the CDS to be used for computations or further identification
Characters 33 - 35	Day (001 thru 366)
Characters 36 - 40	Barometric Pressure
Characters 41 - 48	Parameters from the transmitting site to be used for computations or identification
Characters 49 - 54	Test and Run numbers

The IRA indicates the record type and identifies the test site. In the data record, it also gives the number of channels (data words) that were recorded.

The data records produced by the formatting generator are of a variable length. The length varies with the number of data words that are recorded. Characters 1 - 24 of the data record are of the same format as those of the title run record. Characters 25 and above represent data words. Each data word consists of four characters. As many as 100 data words can be recorded in one record. The record is automatically padded to contain an integral multiple of six characters.



X, IN THE ABOVE ILLUSTRATIONS, MEANS THAT EITHER A 1 OR A 0 MAY BE GENERATED.

### CONCLUDING REMARKS

This paper has described how an automatic data acquisition and processing system was developed to perform a vital function in the modern wind tunnel complex at Northrop Norair. The major benefits of this computerized system can be stated as follows:

1. It reduces wind tunnel data immediately when it is most needed.
2. It permits quicker and more effective adjustments to be made to the model within the test chamber.
3. It shortens the time spent in carrying out a series of tests.

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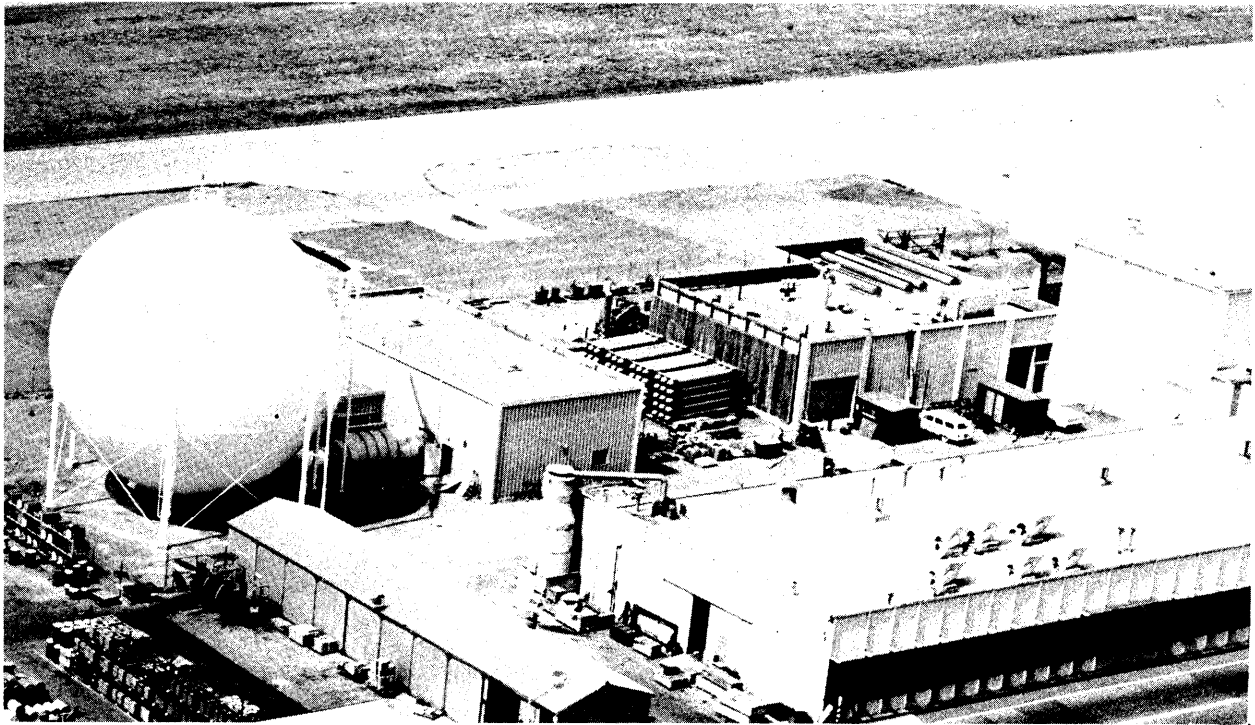
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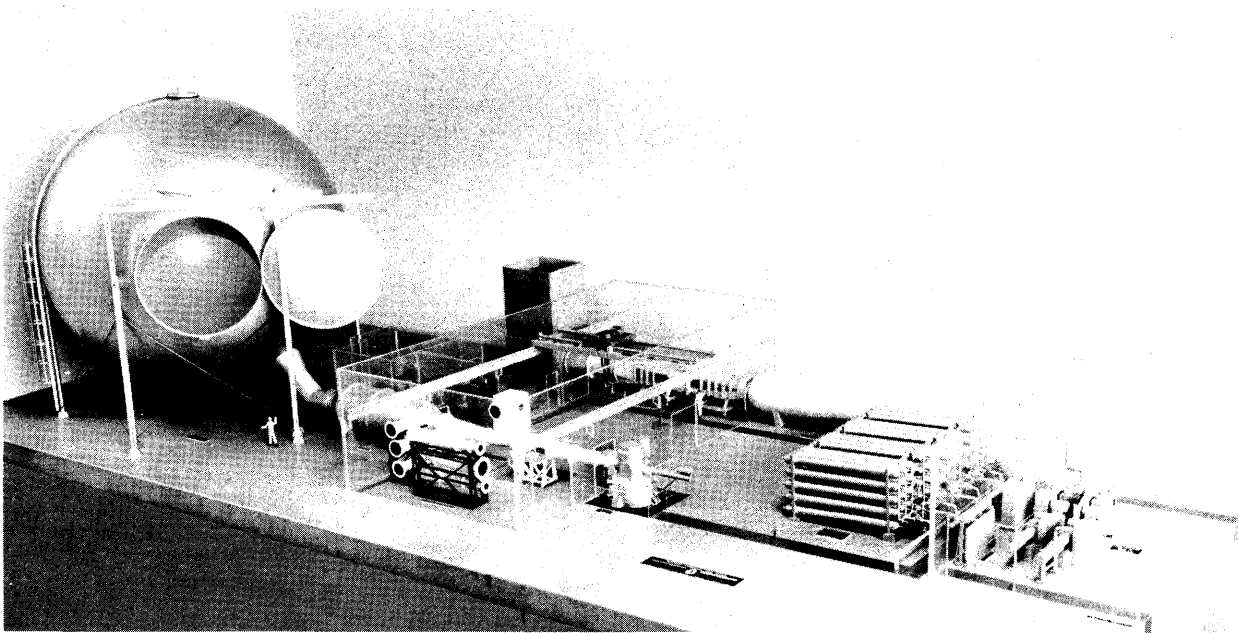
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1043

**OVERALL VIEW OF WIND TUNNEL COMPLEX, SHOWING SUPERSONIC — HYPERSONIC FACILITY AT LEFT AND SUBSONIC AT EXTREME RIGHT**



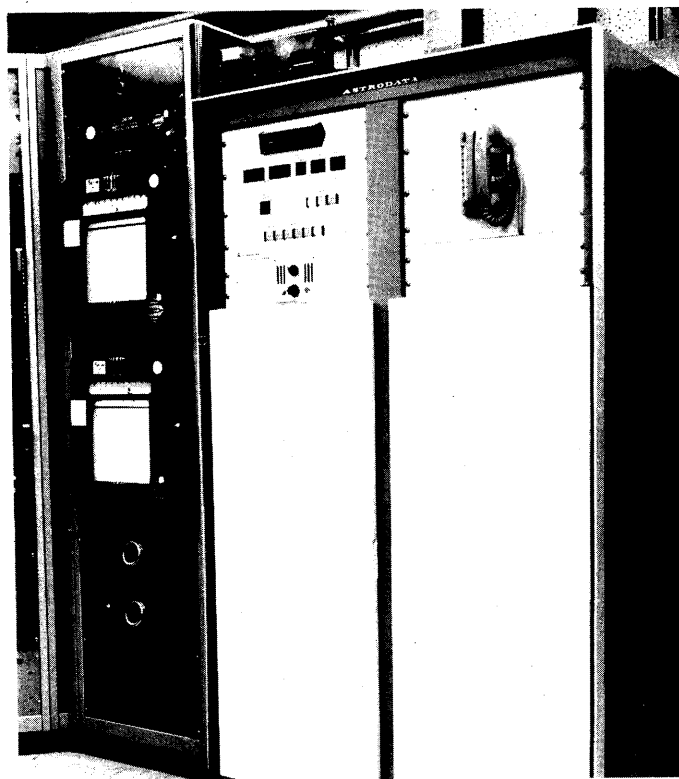
12113

**MODEL OF SUPERSONIC — HYPERSONIC CIRCUIT**



12115

CENTRAL DATA ACQUISITION SYSTEM



SIGNAL CONDITIONING CABINET AT SUPERSONIC SITE

1620 IPL-V

A NON- NUMERIC PROBLEM SOLVING TOOL

by

Wendell Terry Beyer

An essay  
submitted to the Department of Mathematics  
of the University of Oregon  
in partial fulfillment  
of the requirements for the degree of  
Master of Arts

April 1964

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## Preface

This paper is composed of three sections. Section I introduces the need for computer languages similar to IPL-V, section II outlines the IPL-V language, and section III describes the IPL-V implementation for the IBM 1620 computer. A detailed description of the IPL-V instructions and a sample problem are contained in the appendix. A list of selected references is given at the end.

## I

Stored program digital computers were initially developed as devices for performing complex arithmetic calculations at high speeds. At first, the task of programming these machines was burdensome because all programming was done in machine language. However, programming languages were soon developed as an aid to the programmer, beginning with low level assembly languages for specific machines and eventually evolving into high level, machine independent languages such as FORTRAN, ALGOL, and COBOL. Due to the arithmetic origins of the computer, these languages were designed to assist the programmer in the coding of arithmetic or numeric problems.

For a long time, however, it had been known that the digital computer, with its ability to analyze data and take differential action, was not inherently limited in scope to numeric problems. Indeed the problem of translating source statements from a high level language like FORTRAN into machine code is itself a problem basically non-numeric in nature. Other problems for which computer solutions were sought include chess, bridge, analytic differentiation and integration, language translation, pattern recognition, study of learning and self-organizing systems, information retrieval,

theorem proving, and most recently theory developing.

As interest in these and similar problems grew, certain questions arose. Is the present form of digital computer, designed with numeric computations in mind, necessarily the best for non-numeric problems? If not, what better designs might there be? Is it in fact possible to believe that one design will be capable of handling the majority of non-numeric problems? Is it possible to develop a high level language which will do for non-numeric computation what FORTRAN does for numeric computation?

Today these questions remain largely unanswered. No one has succeeded in developing a high level language designed for non-numeric computing although work is being carried on in this area. Some computer designs have been developed which seem to yield a better method of attack on non-numeric problems than that afforded by numeric computers.

To more fully appreciate the problems confronting the designer of a non-numeric computer, it is necessary to examine some of the common characteristics of the non-numeric problems listed above. These problems cover a wide variety of topics and one might suspect that there is little in common among them; however, four characteristics do appear in most of the problems.

First, each problem is non-numeric in part or in

whole. The great computational power of modern numeric computers is not needed.

Second, in most of the problems there is a need for a unit of data more complex than a simple number or array of numbers. For example in analytic differentiation, some method of representing algebraic formulae is needed. In language translation or theorem proving some method of representing syntax or theoretical relationships must be provided.

Third, in many of the problems the assignment of specific areas of memory to contain certain types of information is difficult or impossible since the form, structure, and amount of information is not known at the time a program is set in action. For example, in many cases it is not known what form a self-organizing system will take, or what concepts, and hence information, a theory developing program will yield.

Fourth, it is often desirable to have certain portions of a program call on themselves as subroutines. This is called recursion and is useful in differentiation or game playing where a routine may call on itself to look ahead a move.

A successful non-numeric computer, if it is to have general applicability, must be designed to meet these four needs. Similarly, any language aimed at non-numeric work must fill these needs.



IPL-V is an abbreviation for Information Processing Language V, a highly successful and widely used language designed for non-numeric computing. The IPL languages were developed at the RAND Corporation by Newell, Shaw and Simon, beginning in 1954 with IPL-I, a language for playing chess. Of the IPL languages, IPL-V is the only one which has seen widespread use. The language is well-documented and a manual for programmers is available.<sup>1</sup> In the next section IPL-V is outlined and the manner in which it meets the four problems posed above is discussed.

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<sup>1</sup>See reference [3].

## II

The IPL-V language may be regarded as an assembly language for a non-numeric computer, the IPL computer, or as a medium level language which is machine independent and is executed on numeric computers by an interpreter program. It is interesting to note that an IPL computer has never been built, and all work done in IPL-V is accomplished by means of interpreters. Nevertheless it is useful to describe the IPL-V language in terms of the IPL computer.

It is the function of the IPL computer to manipulate symbols, that is, to accept as data, members of a certain set of symbols, to store these symbols in memory, move the symbols from one location to another, compare the symbols, make decisions based on these comparisons, organize the symbols in memory in a meaningful manner and produce as output a sequence of symbols. For this reason IPL-V is often referred to as a symbol-manipulation language.

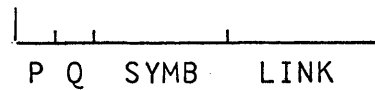
The memory of the IPL computer is divided into cells, and it is the addresses of these cells which form the symbol manipulated by the computer. That is, an IPL symbol is the address of a cell in the IPL memory. The meaning assigned to these symbols is arbitrary. Thus regardless of the contents of cell 14613,

the address 14613 may represent New York City in a military problem, Act II of Hamlet in a literature analysis, or the principle of mathematical induction in a theorem proving problem.

Since it is inconvenient for a programmer to deal directly with memory addresses, the IPL-V language allows a more convenient external representation of symbols. The thirty-six characters A B C ... Z \$ = . + - \* / ) ( and , are called regional characters. At the beginning of his program a programmer may assign to each regional character a continuous block of cells in memory. The block of cells assigned to say A is called the A region and the individual cells in this region are referred to by the symbols A0 (or simply A) for the first cell, A1 for the second cell, etc. Any symbol naming a cell in one of the thirty-six regions is called a regional symbol. The assembler translates regional symbols into the corresponding addresses. In addition the IPL computer has the ability to transform the address of any regional cell into the correct regional symbol during output operations. The address of any cell not assigned to a region is a non-regional symbol, and may be represented by the programmer in a variety of ways.

Each cell in the IPL memory contains two digits called the P and Q digits and two addresses called the SYMB for symbol and the LINK. A typical cell in memory

is represented by the following diagram:



The individual portions of a cell are not addressable.

Cells may be used for one of three purposes: to contain an instruction for the IPL computer, to contain data, or to contain information necessary to the functioning of the IPL computer.

There are a fixed number of cells of the third type and three regions are automatically set aside to contain them. The H, W, and J regions always contain the same cells in memory. The cells of the H region function as registers and indicators in the IPL computer. The W region contains some cells usable by the programmer as temporary storage and other cells used in exercising a certain degree of control over the operation of the computer. Each cell in the J region represents and contains the first instruction of a built in subroutine, of which there are 188 in a complete system.

With the exception of the H, W, and J cells, any cell in memory may be used to contain data or an instruction, and during the course of a program, may contain both.

A cell containing data may be of two types. A data term is a cell containing special alphanumeric or

numeric information, while the P and Q digits indicate the type of information.<sup>2</sup> A standard data cell is a cell used to store an IPL symbol. The symbol is stored in the SYMB and the P and Q digits indicate the type of symbol.<sup>3</sup> The LINK of a standard data cell also contains a symbol, the use of which will be described below. A data cell containing the symbol "+" might look as follows:

0	0	14613	00000
P	Q	SYMB	LINK

where 14613 is the address of the first cell in the "+" region. Unless attention is to be called to the P and Q digits, this will be represented by

+	0
---	---

<sup>2</sup>All data terms have a Q digit of 1 which serves to distinguish them from standard data cells which have a Q digit of 0, 2, or 4. The P digit of a data term indicates the type of information stored in the data term as follows:

P=0	Decimal integer
P=1	Floating point number
P=2	Alphanumeric
P=3	Octal number

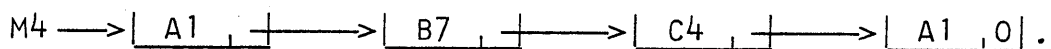
<sup>3</sup>Standard data cells usually have a P digit of 0 although they may be specially marked by a P digit of 1. The Q digit indicates the type of symbol contained in SYMB as follows:

Q=0	SYMB is regional
Q=2	SYMB is local
Q=4	SYMB is internal

The data terms play a rather minor role in the computer, usually serving as storage locations for numeric information; while the role of the standard data cell is central to the operation of the computer.

In dealing with symbols of arbitrary meaning, the IPL computer answers the first need of a non-numeric computer, that of dealing with non-numeric information. These symbols do double duty, serving sometimes as the addresses of cells in memory and at other times representing the concept assigned by the programmer. However, the IPL symbol, being an address, is basically no more complex than a number.

The need for a complex unit of data is fulfilled by the list, a basic unit of data in the IPL computer. A list is a sequence of data cells which are joined together by having the link of each cell contain the address or name of the following cell. A list of the symbols A1, B7, C4, and A1 in that order would be represented by the following diagram,



where the arrows indicate the cell referenced by the LINK of a cell. Note the use of the symbol 0 in the link of the last cell. This symbol is called the termination symbol and indicates that the list terminates at that point. The name of the first cell in the above list is

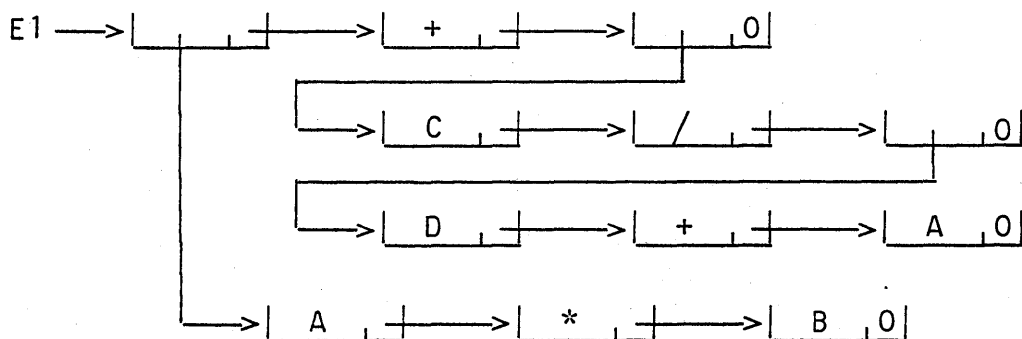
M4 and the list is also referred to by that symbol.

Given cell M4, any symbol on the list may be reached by passing from link to link.

Far more complex structures may be created by using the SYMB of some cells on a list to contain the names of other lists. The Q digit of a cell on a list may be used to indicate whether the SYMB contains an abstract symbol or the name of a sublist which is to be considered part of the structure.<sup>4</sup> The number of structures possible is limited only by the programmer's imagination, but for simplicity only lists will be considered below.

Because of the list, IPL-V is called a list-processing language, as are other languages which use the same concept. The language contains subroutines for list manipulations such as copying, printing, searching, or erasing lists. An example is a subroutine which will

<sup>4</sup>For example, the algebraic expression  $A*B+C/(D+A)$  may be represented by the structure E1 below which expresses the structure of the expression in a manner not possible in a linear list representation.







number of cells in memory is of concern to the IPL-V programmer.

It might seem that locating an unused cell in memory would be difficult, but this problem is handled in an elegant and efficient manner. After assembly, all unused cells are linked together to form a list named H2 and called the available space list. During processing when a cell is needed, one is removed from H2 for use; and when a cell is no longer needed by the programmer, it is returned to H2.

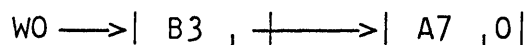
The list organization also allows cells to be used as though they were capable of storing more than one symbol. Suppose for the moment we have a symbol stored in cell W0, say A7, and we need to temporarily store a second symbol, say B3, also in W0.

W0 → | A7 , 0 |

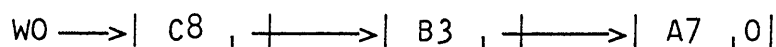
We execute an IPL instruction causing the computer to push down cell W0. That is, an unused cell is removed from H2, inserted behind cell W0, and a copy of the symbol in W0 is placed in the new cell, creating the following list:

W0 → | A7 , 0 | → | A7 , 0 |

Now that a copy of A7 has been made, B3 may be placed in W0.

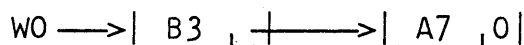


We may go even further and store C8 in W0 before removing B3, by pushing down W0 again, then storing C8.

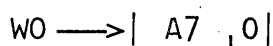


The list created in this manner is called a push down list but is no different from any other list.

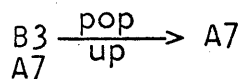
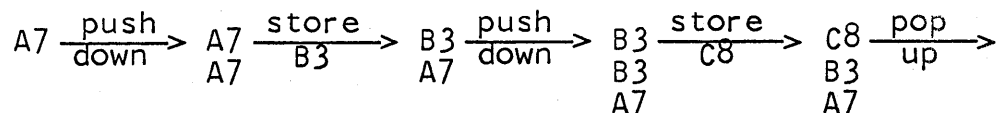
When the symbol C4 is no longer needed in cell W0, a pop up instruction is executed. This operation copies the second symbol on the list into the first cell and removes the second cell from the list, returning it to H2.



One more pop up, and W0 is returned to its original state.



The preceding sequence of events may be summarized by writing the push down list vertically.



The push down and pop up instructions enable a subroutine and main routine to use the same storage cells. A set of working cells, W0 through W9, are provided for temporary storage. When a subroutine needs temporary storage, some of these cells are pushed down, then used as storage. Any information stored by the main routine in these cells is preserved by the push down operation. Before terminating, the subroutine pops up these cells, returning them to their original state.

The ability of the IPL computer to allow recursion, the fourth need of a non-numeric computer, is also based on the push down operation. The cell H1, called the current instruction address cell, contains at any given time the address of the instruction currently being executed by the IPL computer. When an instruction is completed, the address of the next instruction is obtained and placed in H1. Like any other cell in the memory, H1 may be pushed down. When one routine calls on another as a subroutine, H1 is pushed down by the computer, saving the address of the instruction in the main routine where processing is suspended. The address of the first instruction in the subroutine is placed in H1 and that instruction is executed. Processing now continues along the subroutine and the computer is said to have descended a level. Before terminating, the subroutine may call on itself or another subroutine.

Again H1 is pushed down, saving the point at which processing was suspended in the subroutine, and processing continues at a lower level. When a subroutine terminates, H1 is popped up and the routine one level up resumes action. A combination of the manner in which H1 is used and the ability of the working cells to keep the contents of routines on different levels from becoming mixed, allows a subroutine to call on itself.

The instructions in the IPL computer are kept in lists. The P, Q, and SYMB of a cell make up the instruction and the LINK indicates the next instruction. The IPL computer follows instructions from cell to cell down a list rather than executing instructions sequentially in memory. This allows routines to be manipulated with the list processing subroutines. It is conceivable that a main routine could construct a subroutine using list processing subroutines, execute that subroutine, then erase it, that is, return all of its cells to the available space list.

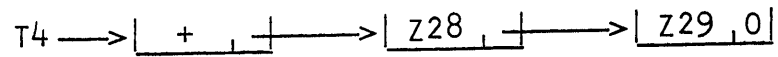
In communicating information to a subroutine, a special cell H0, the communication cell, is used. The symbols required as inputs by the subroutine are placed in H0 using the push down operation. The subroutine accepts these inputs, removing them from H0, and before terminating, places all output symbols in H0 where they are recovered by the main routine.

In addition to producing output symbols, some subroutines produce a yes or no answer. For this purpose a cell called the test cell, H5, is provided. The test cell may be in one of two states, "+" or "-", and an instruction is provided to allow conditional branches or transfers within the program on the basis of the state of the test cell.

There are only eight basic instructions in IPL-V, most of the processing being done by the numerous subroutines. Two instructions are used for placing symbols in H0, one instruction for calling on subroutines, two instructions for removing symbols from H0, one instruction each for popping up cells or pushing down cells, and one instruction for conditional branching on the status of the test cell. The P digit determines the type of instruction and SYMB contains a symbol, the name of a cell, or the name of a subroutine, depending on the context. The Q digit is used in connection with SYMB for three levels of addressing. For more complete information concerning instructions and a sample routine, see the appendix.

The external form of IPL-V is quite simple. Lists, instructions or data, are written vertically on the coding sheet. Each line represents a cell and space is provided to indicate the name of the cell and the P, Q, SYMB, and LINK of the cell. If a link is left blank, the cell is

assumed to link to the cell on the following line of the coding sheet and the name of the following cell may also be left blank if its memory location is unimportant. Thus to create the list



we write on the coding sheet

<u>NAME</u>	<u>PQ</u>	<u>SYMB</u>	<u>LINK</u>
T4		+	
		Z28	
		Z29	0

### III

The University of Oregon IPL-V system for the IBM 1620 computer, developed and written by the author and John D. MacDonald, was designed with two objectives in mind. It was intended first as an educational device to acquaint students with list processing and symbol manipulation problems, and second as a system for checking out IPL-V programs before running them on larger computers. In view of the educational aim, operating speed was sometimes sacrificed for operating ease and additional safeguards. Because of the speed and size of the 1620, the system was never intended as a production tool.

The 1620 system is based on the specifications of IPL-V set forth in the manual<sup>5</sup> and is fully compatible with those specifications, though not all options are available on the 1620 system. Operating on any 1620 equipped with card I/O, indirect addressing, automatic divide, and special instructions, the system provides approximately 640 IPL cells at run time with a 20K memory. An additional 1,660 cells are available with each additional 20K of memory. The system operates at approximately 80 IPL instructions per second and is equipped

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<sup>5</sup>See reference [3].

with all tracing and monitoring features specified by the manual. These features include operator or program controlled trace with output on any unit, automatic trapping on error conditions, and flexibility in trap recovery.

The system consists of three decks, the assembler, subroutines, and the interpreter, which are loaded in that order with the source deck placed between the assembler and subroutines. The assembler loads into the lower portion of memory and assembles the source deck directly into the upper portion, producing an assembly listing on option. Next the subroutine deck is read by the assembler and those subroutines called for are loaded into memory. After the last card of the subroutine deck has been read, the interpreter loads into the lower portion of memory, occupying the space previously occupied by the assembler; the computer halts; and execution begins when START is pressed.

The internal form of an IPL cell is a twelve digit field with an odd address. From low to high address the cell contains the P, Q, SYMB (five digits), and LINK (five digits).

Provisions are made for writing additional subroutines in SPS and including them in the source deck. It is also possible to reserve blocks of space in the 1620 memory for use by other systems. Methods



of setting up linkage between systems are described in the documentation.

The documentation is in the form of an appendix to the manual<sup>6</sup> with cross references. A master copy of the documentation is maintained on cards for easy editing and reproduction.

During the summer and fall of 1963, a preliminary version of the system was written. This version was distributed to approximately twenty participating users for field testing and was used in a one term seminar in IPL-V programming at the University of Oregon. Students in this seminar used the system for problems such as analyzing poetical structure, construction of Farey sequences of numbers, calculation of all closed paths in a planar graph, and construction of a machine for playing Hex. The system has also been used for map coloring and analytic differentiation.

The preliminary version does not contain block handling, auxillary storage, read/write, floating point, save for restart, or post mortem dump routines. During the summer of 1964, a final version will be written, which will include all features except auxillary storage processes. The final version will be submitted to the 1620 Users Group's General Program Library for distribution.

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<sup>6</sup>See reference [3].

## Appendix

### The IPL-V Instruction

The Q digit of an instruction operates on the SYMB to produce a transformed symbol S as follows:

- Q=0 S is SYMB.
- Q=1 S is the symbol contained in the cell whose name is SYMB.
- Q=2 S is the symbol in the cell whose name is contained in the cell named SYMB.

For example, if we have the following cells in memory,

```
A1 → | T4 ,0 |
T4 → | J8 ,0 |
```

and the SYMB of the instruction contains A1, the Q digit produces the following transformations:

PQ	SYMB	S
0	A1	A1
1	A1	T4
2	A1	J8

The transformed symbol S is stored in a register; the SYMB portion of the original instruction is never altered in memory.

After the transformed symbol has been obtained the P digit determines the action as follows:

- P=0 call on the subroutine whose first instruction is in cell S.
- P=1 push down H0 and place a copy of the symbol S in H0.

- P=2 copy the symbol in H0 into cell S, then pop up H0.
- P=3 pop up cell S.
- P=4 push down cell S.
- P=5 same as P=1 except H0 is not pushed down first.
- P=6 same as P=2 except H0 is not popped up afterward.
- P=7 if H5 is -, transfer to cell S for the next instruction. if H5 is +, continue.

### Sample Problem

As an example of how the instructions are used, we will write a short subroutine below. It will be necessary to understand the operation of two of the J routines.

J2 accepts two inputs in H0. Each input is a symbol. J2 compares the symbols and sets H5 "+" if they are equal and "-", if not. J2 leaves no symbols as output in H0 and the two input symbols are no longer in H0 after J2 terminates.

J60 accepts one input which is the name of a cell on a list. If that cell is the last cell on the list, J60 sets H5 "-" and leaves the input as output. If the cell is not the last cell on the list, J60 places the name of the following cell in H0 and sets H5 "+".

We now code the routine E4. E4 is a routine which evaluates a function of X at a given point. More clearly, E4 accepts a symbol representing a given point,

say A, and a second symbol assumed to be the name of a list representing a function of X. For example, the list F1 below:

```

F1      X
        *
        L
        O
        G
        (
        B
        /
        X
        )      0      "xlog (b/x)"

```

E4 then evaluates the function at A by replacing every occurrence of the symbol X on the list by the symbol A to yield the list:

```

F1      A
        *
        L
        O
        G
        (
        B
        /
        A
        )      0      "alog(b/a)"

```

E4 should leave no output in H0. In addition since E4 does not set H5 as part of its output, the status of H5 should be the same after execution of E4 as before. But E4 must call on J2, which does reset H5. For this reason, it will be necessary to push down H5 at the beginning of E4 to save its status, then to pop it up at the end to restore its status. Two storage cells, W0 and W1 will also be needed. It is assumed that the routine which called on E4 input the name of the function

list first, then the symbol representing the point. A little study and liberal use of a black board as a simulator will make the operation of E4 clear. The symbols 9-1, 9-2, and 9-3 are called local symbols and are used for internal branching within the routine.

<u>Name</u>	<u>PQ</u>	<u>SYMB</u>	<u>LINK</u>	<u>Comments</u>
E4	40	H5		Preserve H5
	40	W0		Preserve W0
	40	W1		Preserve W1
	20	W0		Move "point" to W0
9-1	12	H0		Input symbol in list cell
	10	X		Input X
	00	J2		Compare symbols
	70	9-2		Go to 9-2 if not equal
	60	W1		Copy list cell address in W1
	11	W0		Input point symbol
9-2	21	W1		Move point symbol to list cell
	00	J60		Find next list cell
9-3	70	9-3	9-1	If no list cell, clean up
	30	W0		Restore W0
	30	W1		Restore W1
	30	H5		Restore H5
	30	H0	0	Pop up H0, terminate

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PETROLEUM EXPLORATION AND PRODUCTION APPLICATION  
FOR THE IBM 1620 AND PLOTTER

By

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Delivered at:

IBM 1620 Users Group  
Western Region Meeting  
Denver, Colorado  
June 17-19, 1964



It was refreshing to hear Dr. Edward N. Brandt, of the University of Oklahoma Medical School Biostatistical Laboratory, say in his keynote address that the problems dealing with computers in the field of medicine are such that they are basically related to and parallel the problems which are encountered in the oil industry. Dr. Brandt also related that the use of computers in medicine has required that the users better define their problems, which gives them a better understanding of the overall situation. The same can be said about the use of computers in the oil industry.

In the next 15-20 minutes, I plan to tell you a little about the Oil Information Center which is an integral part of the University of Oklahoma Research Institute. I will discuss the Oil Information Center:

1. Why and how it was established
2. The goals and objectives
3. How it is connected with the University computer usage generally and the IBM 1620 specifically
4. What we are presently doing, and
5. Where we are going

#### I. GENESIS OF THE OIL INFORMATION CENTER

Two independent oil men in Oklahoma, Mr. Ward Merrick, Ardmore, and Mr. Howard McCasland, Mack Oil Company, Duncan, were concerned about three

apparently unrelated situations and problem areas in Oklahoma. These three problem areas were:

1. No attempt had ever been made to gather groups of oil field related information on a library basis.
2. The Oklahoma Corporation Commission needed an assist in some of their data processing problems and engineering calculations.
3. The computers at the University of Oklahoma were not being utilized as much as could be reasonably expected by local industries, particularly the oil industry.

The concern of these two independent oil operators led them to the concept of the Oil Information Center and as a direct result they furnished the impetus by supplying financial assistance through the medium of their personal foundations. A two-year budget was set up for the initial phase of this Center.

One obvious objective of the Oil Information Center was that sooner or later it must become self-supporting from earned income. It was felt by all concerned that these problem areas just mentioned would be the strong nucleus upon which the objective of self-support would be reached.

After a series of conferences, oil industry executives and University people agreed that the logical central location for libraries of oil information would be on the campus at the University of Oklahoma. The categories of information which seemed desirable to collect were electric logs, scout tickets, drillstem tests, sample logs,

and Oklahoma Corporation Commission completion forms. The University of Oklahoma has been famous for years in the quality and quantity of graduates pointed toward the oil industry. The University has probably turned out as many petroleum geologists, petroleum geophysicists and petroleum engineers as any university in the United States.

The Oil and Gas Conservation Department of the Oklahoma Corporation Commission needed assistance with some data processing problems. They wished to work directly with a group who could help them in their work, on whose integrity they could rely and in whom they could have confidence. The Oil Information Center devised a plan to prepare computer programs to assist with some of these problems, and Commission representatives gladly accepted this plan.

## II. OPERATIONS OF OIL INFORMATION CENTER

### A. Introduction to University Relations

The actual operation of the Oil Information Center is concerned with various areas of effort. A major area is connected with university activities. These are:

1. The graduate program of the University of Oklahoma
2. The Oklahoma Geological Survey, The University of Oklahoma Schools of Geology and Petroleum Engineering

3. Conducting seminars on oil related topics
4. Attracting people in the oil industry to the campus

#### Geology Graduate Student

In checking the records I found no evidence indicating that any geology graduate student had used the computers or plotter to assist them in their master's thesis work. I sought out someone who might be interested in using the computer and found a Humble Oil geologist, on leave from his company to do master's work, and who was willing to work with me. Since the geologist was not a programmer, arrangements were made for his programs to be written for him and through the cooperation of the Computer Lab his key punching was accomplished. This graduate student's thesis was on the geology of an oil field in North Central Texas. His study of the electric logs on each well furnished him with formation tops, well elevations, etc. for his study of 25 different formations. With this information punched into cards he was ready to use the 1620 and plotter to prepare his isopach and subsea calculations and his many maps. The computer program as written was general enough that calculations could be made for isopach thicknesses, subsea formation tops, and sandshale-limestone ratios. This is an example of what can be done in working with graduate students and we hope to encourage others along these lines.

Oklahoma Geological Survey and University of Oklahoma  
Schools of Geology and Petroleum Engineering

The Oil Information Center has attempted to work closely with the Oklahoma Geological Survey and the University of Oklahoma Schools of Geology and Petroleum Engineering. The libraries of oil field information being gathered by the Oil Information Center are a valuable complement to the Core and Sample Libraries now existing at the University of Oklahoma. The Geological Survey uses the electric logs, sample logs, drillstem tests, etc. in their state-wide geologic investigations. The Schools of Geology and Petroleum Engineering can use the same information as teaching aids.

Conducted Symposiums

An important activity in the university phase of our operation is the conduction of symposiums. The Oil Information Center, in conjunction with our libraries of information and computer services, has conducted two symposiums on the campus. One was related to our Drillstem Test Library to which we were able to get good industry speakers from all over the Southwest.

The second symposium was directly connected with the Mid-Continent Well Data System in Oklahoma City. In addition to the speakers at this meeting, the Oil Information Center in cooperation with the University of Oklahoma Computer Lab demonstrated an information retrieval program. I shall discuss this demonstra-

tion in more detail in a few minutes. These symposiums have been extremely helpful in our relationship with oil industry people, particularly on the operating level. The sharing of new ideas and approaches is always helpful.

#### Bring People to the Campus

Directly through the efforts of the Oil Information Center a large number of people have been directed to or through the University of Oklahoma campus. Our seminar on drillstem testing attracted 148 people for two days of meetings. The Mid-Continent Well Data System Symposium was for one day and was attended by 65 people.

Major oil company and consulting geologists from Tulsa, Ardmore, Norman, Ada, and Oklahoma City have been to the Oil Information Center libraries for various reasons. Major oil company representatives have also been to our computer installations using our computer and plotter services. Others have investigated the services which we have to offer in order to determine how this information could be beneficially used by their company.

#### Industry Effort

To the best of my knowledge, this is the first industry wide effort of information gathering undertaken by the University of Oklahoma. Acceptance of the oil libraries could well lead to the establishment of the gathering of information in other fields of endeavor.

## Oil Industry in Oklahoma

With the advent of oil industry data retrieval pilot studies in West Texas, the Oil Information Center found it advisable to conduct their own pilot project on the digitizing of scout tickets and a retrieval program to recover this information. The Autwine field in Kay County, Oklahoma, was chosen for this study for several reasons. The field has more than one producing zone; it produces both oil and water; both major oil companies and independent oil operators have wells in the field. Scout tickets were received on 122 wells which included some surrounding dry holes, and the information was keypunched to our predetermined format.

A computer program was written for our 1410 to retrieve certain information from these cards. The program was written to gather certain usable groups of information:

1. List the wells which cored the Red Fork formation,
2. List the wells and the detailed results of all drillstem tests in the Red Fork formation,
3. List the casing programs in each well,
4. List the formation tops from some wells,
5. List each well that penetrated the Mississippi formation, and
6. List the details of the acid and fracture treatments on each producing Red Fork well.

These are some of the categories of information chosen to be retrieved for this demonstration. This information is typical of that which is used by the exploration geologist and the petroleum engineer in some of their everyday problems.

## Oklahoma Corporation Commission

### Preparation of Oklahoma Guymon-Hugoton Gas Allowable Schedule:

Due to the large amount of paper work which they process, and their general work load, Gas Conservation Department personnel often were two or three months late in the preparation and distribution of the Guymon-Hugoton Gas Allowable Schedule. By the time the operators of the well and the purchasers of the gas received the schedules they were practically of no value.

The Oil Information Center worked as liaison between Corporation Commission engineers and the Computer Lab programmer so that a computer program could be written to calculate the monthly gas allowable for each well in the field. When Corporation Commission personnel prepared this gas allowable schedule on a desk calculator, they required approximately 70-75 manhours per month. After an estimated five hours of keypunching and keyverifying per month, the IBM 1410 makes these calculations to prepare this gas allowable schedule in 0.4 hours per month.

### Calculate one-point back pressure test:

An Oklahoma Corporation Commission statewide rule makes it mandatory for all allocated gas wells to annually report a one-point back pressure test. This information is used in assigning per well gas allowables for the following year.



An estimated 1,800 - 2,000 of these tests are filed with the Commission each year and the Gas Engineer is required to check each of the calculations. The Gas Engineer informed me that with no interruptions he could check five or six of these calculations per hour. This meant that two or two and one-half man-months per year was spent in checking these previously calculated tests. An O.U. Computer Lab programmer wrote a program for our 1410 to make these calculations. The 1410 processes these tests in 4.25 hours, which is a significant dollar saving estimated at 3-1/2:1. This Gas Engineer is now freed to do more productive and original work for the Commission, which represents the true saving.

#### B. Introduction to Commercial Applications

Our other major effort is the industrial commercial activities. We have worked directly with:

1. Major oil companies
2. Independent oil operators
3. Oil-field service companies
4. Petroleum consultants

In mid-1963 IBM released a group of programs from their 1620 library, which are called the Petroleum Package. These programs were written by experienced petroleum engineers, geophysicists, and geologists for a rather wide range of

commonly encountered exploration and engineering problems. The engineering programs deal with primary oil recovery, secondary recovery, economic evaluations, casing design, gas production rates, flash calculations, etc. The exploration programs deal mainly with geophysics, but are also related to map contouring, electric log analysis, dipmeter calculations, map preparation, etc.

In the past ten years petroleum oriented companies have become more dollar conscious and overall economics have played an ever increasing part in top management decisions. Computers are being used more and more to funnel detailed geophysical, geological and petroleum engineering information to these top management people for their perusal in making their decisions.

In the recent past it was not feasible to make many groups of calculations in the fields of geophysics, geology and petroleum engineering. These calculations were known applications and approaches to their problems but were too detailed and too time consuming for the engineer or geologist to justify spending the time from his other daily duties. With the advent of computers, it became more realistic to consider making some of these calculations. Also, in the past, the necessary data to make these calculations were not gathered knowing that they would never be used. Such is not the case now, and it should be pointed out that the gathering of these data in many cases makes for a more efficient operation on all levels.

In several application areas the use of digital computers is becoming more valuable as magnetic tape recording devices are used in the field. Some of these instances are:

1. Electric logs (and their companion logs)
2. Dipmeter surveys
3. Geophysical field surveys

Many of the large oil field service companies are installing magnetic tape recording devices in their field trucks. This will lead to a more detailed study of data now being received but not efficiently used.

However, most of the commercial work which we have done in our 1620 Lab is related to geophysical problems. The reason is rather obvious when the users were questioned. In many instances geophysicists were not making certain known approaches to their problems because of the number of manhours required to prepare the data, make the calculations and plot certain information. The use of computers and digital plotters now makes it more practical to better utilize data gathered in the field by geophysical crews.

As some of you know, a reflection seismograph crew costs an oil company between \$15,000 to \$60,000 per month depending on the overall services rendered and the field equipment involved. As in most any other service operation, reflection seismograph field crews can and do have certain problems.

If the field data are being processed on computers as work progresses, the errors can quite easily be rectified. However, if there is a large time lag between the error and its discovery, it may not be so easy to make the necessary adjustments.

A geophysical group of a major oil company in Oklahoma City has been our largest user of commercial time on our 1620 and plotter. This District office is responsible for the geophysical work in all of Oklahoma, all of Kansas, the Texas Panhandle, North Central Texas and the northern 2/3 of Arkansas. In addition to the reflection seismograph field crews gathering new data, they are continually reviewing old seismic records previously shot by themselves or by other companies.

One geophysicist pointed out the following, relative to the information gathered from 300 shot-points. The time required to hand calculate and hand plot this data from 300 shot-points would be an estimated two man-months. To use computers, this same amount of work would require an experienced geophysicist one week, another week to key-punch, one to one and one-half hours on the 1620 for calculation, and five and one-half to six hours on the 1620 and on-line plotter. This represents a vast saving of time as well as money.

One geophysicist pointed out that the use of our 1620 computer on their reflection seismograph field data makes it possible for them to better utilize the

information which can be gathered from seismic records. He said that they can now prepare ten to twelve useful sub-surface maps where previously they were fortunate if they were able to get five to six maps from a set of seismic records.

Dan Merriam of the Kansas Geological Survey and John Harbaugh of Stanford University through their joint effort developed a computer program to assist in the location of mineral deposits. (1) Based on certain known geological and/or geophysical information and certain mathematical computations trend surfaces are fitted so that the sum of the squared deviations is the least possible value.

The trend surface analysis may be used to:

1. Predict projected depths to geological units within an area,
2. Delineate unconformities or changes in structural patterns, and
3. Extend better "geologic guesses" into adjacent unknown areas of no control.

Close agreement exists between local structural features and trend-surface residuals. The residual maps were found to stress or emphasize trend relationships not otherwise clearly observed from original data and to emphasize the local component of the structural pattern by essentially removing the regional component or regional dip. Inasmuch as in many regions the oil and gas producing areas are systematically associated with structural features, there is the possibility that a study of the residuals will indicate previously overlooked areas favorable for additional oil exploration.

The Oil Information Center plans to take advantage of the existence of this program but we plan to rewrite the program to use the IBM 1620 and plotter rather than using the printer to prepare the map.

John P. Dowds, a successful petroleum consultant in Oklahoma City, has worked on the laws of probabilities and the application of statistical methods to help analyze the problem of obtaining commercial oil or gas production. Dowds, in a recent paper, stressed that "exploration geologists and geophysicists need to become statistically minded and to think of locating oil and gas fields as a problem in applied possibilities." (2)

Dowds uses entropy for his mathematical model to learn of favorable trends and patterns in searching for logical locations for drilling new oil or gas exploration wells.

Dowds determined a long time ago that his calculations were too difficult and the number of these calculations required were too many to be done by hand. An Oil Information Center programmer recently wrote programs to Dowds' formulae for his entropy calculations. These are now being run on our 1620 and plotter. The final output to be studied for purposes of exploration is a series of contour maps. Dowds is representing a large independent oil operation in Oklahoma City in their search for sizable oil or gas reserves.

James M. Forgotson, Jr., research geologist with Pan American Petroleum Corporation in Tulsa, said in a recent Oil and Gas Journal article that the use of electronic computers to evaluate electric logs is very practical. He said, "The speed with which these computations can be performed makes the analysis of many zones or formations in thousands of wells practical." Forgotson went on to point out that "without the aid of the computer, approximately eight manhours are required to calculate shaliness, saturation ratio, and favorability criterion for four zones in one well." He also made an interesting comparison stating that "with the use of computers approximately one and one-half man-months would be required to process four zones in 1,000 wells while without the use of computers fifty-four man-months would be consumed." (3)

### III. SUMMARY

The Oil Information Center is serving a useful purpose to the University of Oklahoma, to the Oklahoma Corporation Commission, and to the oil industry in general in Oklahoma.

With the 1410, 1620 and the plotter now in the University of Oklahoma Computer Lab, we are able to offer computer services to:

1. Major oil companies
2. Independent oil operators
3. Consultant geologists and petroleum engineers
4. Oil field service companies

Using the 1620 Petroleum Package of programs has proven successful up to a point even though the large majority of commercial time which we are able to sell has been to companies who have written their own programs.

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A CONTROL SYSTEM APPROACH  
TO  
AUTOMATIC JET ENGINE TESTING

1620 User's Group  
Western Region  
June 17, 18, 19 - 1964

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## A. History of Jet Engine Testing:

After the first jet propelled airplane was captured from Germany by the United States in World War II, development of the jet type aircraft has proceeded in rapid fire fashion.

The first truly great use of the jet airplane came about as a result of the Korean War. In a few short years since the early 1950's, the development and production of the jet engine has proceeded at an amazing rate.

With the production of the first jet also came problems in the maintenance and overhaul of these complex, high thrust engines. At the beginning, especially during the Korean War, maintenance and repair was carried out in the remote airstrip locations and centralized repair facilities using the out-moded piston engine repair and test facilities. The piston engines had not required the highly substantial and instrumented test facilities that the newer high thrust jet engines were requiring; so, many of the first tests were performed in a crude makeshift manner.

In the initial stages, many of the repair personnel became engine test personnel. Because the jet engine development had proceeded in a hurried fashion, adequate testing procedures were lacking; so many of the first test cell personnel found themselves preparing their own through pooling, interchanging and accumulating their experiences. Many of the basic principles of these early testing technical procedures are still in use today. Also, the great majority of today's test cells are modified piston engine of low thrust jet engine test cells and their instrumentation leaves a lot to be desired. Much of the instrumentation was installed on a "guess and try" basis.

Since the early 1950's the production rate and number of jet engines in the air has risen considerably. With these increases also came increases in the number of engines to be overhauled and repaired. The test facilities in many instances have been updated with new instruments. The engine manufacturers have also been allowed time to adequately prepare better testing procedures. Even with all of these improvements there still remains two pressing problems. They are: (1) the large number of engines awaiting the testing facility and (2) the advent of the higher thrust (turbo fan, J75, etc.) jet engine has again outdated the test facilities.

## B. Present Test Problems:

Because of the rapid expansion of the test facilities to accommodate the increased workload of jet engines and the complexities of the higher thrust engines, many problems arose in acquiring an adequate balance between the production and quality control functions.

These problems are presented in the following sections. They are grouped into areas in order to present a detailed view of each. It should, however, be noted that the problems actually overlap into other areas and even overlap each other. Many times a

particular problem arises because of testing techniques, instrumentation, and the facilities being used.

## 1. Present Testing Methods

In order to fully understand the problems associated with the present testing techniques, the following is submitted as a general discussion of the overall testing procedure.

The typical jet engine test cell has two or three men assigned to it. During the initial installation phases two men perform all necessary physical connections. This will include steps (a) and (b) of the test procedure. During the running/testing of the engine, one of these men will control the throttle and instrumentation necessary to run the engine and make recordings while the other man makes the balance of the necessary recordings at the appropriate times and places. A third man acts as an inspector. His job is to observe the readings being made and perform a reasonableness check on certain limits to see if recording errors have been made. He also takes observed readings and corrects them to a standard day (sea level or other) condition for comparison with the technical order specifications.

On the final analysis he either accepts or rejects the engine based upon its performance within the limits and specifications of the manufacturer's technical order. As the engine is routed to test from final assembly it is complete as required by technical order to the final piece of safety wire.

### a. Dressdown

Upon receiving the engine at the test area, numerous steps are necessary in order to prepare the engine for testing. The first step is checking for possible external damage which might have occurred during transportation. The engine normally is assigned to a particular test cell prior to dressing for test. Special plugs, fittings, and some harness have to be removed in order to install test equipment. Special test harness is installed in order to obtain individual thermocouple readings for temperature spread checks. Various pressure taps are installed throughout the engine in order to obtain internal air, oil and fuel pressures. Engines are so designed that internal pressures must meet certain limits. If engine internal ratios are below values outlined by engine manufacturer, it becomes necessary to change some specific internal clearance in order to obtain required ratios.

The next step is to install a workhorse tailcone or afterburner. Altogether, there are approximately ten test fittings and adapters that must be installed in addition to temperature harness. Engine mounting adapters and bellmouth adapter rings are installed. Finally the engine oil tank is filled to capacity. This about completes the initial dressing. If the engine is designed for an after-burner, then an AB is attached. There are additional functions of preparations to be performed after the engine enters the test cell.

b. Preliminary Check

After installation of the engine in the test stand, it is necessary to perform some inspections at particular times. This will include such inspections as freedom of compressor rotation and making sure no foreign objects are present in the compressor inlet. It is necessary to accomplish this type inspection prior to installing the bellmouth and inlet screen. If an inspection is performed after installation of the bellmouth, it is quite easy to overlook some small item which might result in compressor damage.

c. Preliminary Shakedown

After engine is properly secured in the test stand with all pressure and temperature connections, attached, a complete shakedown is accomplished by a quality inspector. This shakedown is necessary to pick up anything which may have been overlooked during the engine installation.

d. Functional Component Check

The next step is a functional component check out. This consists of selecting the main fuel control emergency system, afterburner system, and anti-icing valves for functional operation. These checks are necessary prior to starting the engine in order to replace such items that may be faulty.

e. Dry Run

Prior to starting the engine a dry run is performed in order to flush preservative oil from the fuel components, pressure fuel and oil system. Leaks are sometimes found during this check. Afterwards, the dry run oil system is replenished and the pressurizing valve sense line reconnected.

f. Running Prior to Acceptance

The engine is then ready for a start. After the engine has started and reached idle R. P. M. a complete shakedown is made to check for air, oil and fuel leaks. If no abnormal conditions are found, power is advanced toward top power and preliminary checks are made on oil pressure, E. G. T. and vibration.

g. Performance Runs

After the preliminary run has been completed, the engine is ready for a performance test run. This test run consists of numerous functions in order to test the basic engine and its attached components as a complete assembly.

Other checks that follow during the actual performance test are acceleration checks, simulated afterburner runs, emergency system runs, oil consumption check and performance calculations.

The test run begins with an initial power advance after start to approximately nine thousand RPM. This is necessary in order to obtain specific data for test run calculations and warm up the engine oil. The engine oil must be heated to actual operating temperature in order to obtain valid consumption during test run. Oil temperature must be noted at the time oil level is checked on a sight gauge and again at completion of the test run. Oil temperature at the time of the final check must be within  $\pm 2$  degrees F. of the initial temperature. Oil consumption is actually determined by visually observing a sight gauge. This sight gauge is calibrated to the engine oil tank and actually seeks the oil level within the engine tank. The oil level sight gauge is marked with ten increments to the inch and each increment represents a specific amount of oil.

The data collected during the initial warm up period is used to determine the exact power position required for various test runs. Four power runs ranging between seventy-five per cent and take off are required in order to help determine the quality of the engine.

Test run power positions are determined by charts representing given thrust positions. All data from such charts represent standard day conditions biased for temperature variation. Actual thrust requirements are subtracted from points corresponding to various power positions by using compressor inlet temperatures. Once having obtained required corrected thrust output, this data must be converted to actual time conditions. This correction is a function of present time condition variations from a standard day and test cell correction factors.

Each individual run has a time duration of five to twenty minutes depending upon the position of power. Recordings of internal pressures from compressor inlet to turbine discharge are made. Temperatures of air inlet, oil, fuel and turbine discharge air are logged. Other recordings such as fuel flow, thrust, turbine discharge pressure, RPM and vibration are necessary.

All data logged directly related with the functional operation of the engine must be corrected to a standard day condition. This data is also corrected for compressor inlet temperature, barometric pressure and test cell correction factors. There are approximately 175 calculations performed during the test run. Thirty-five points are plotted on special graphs in order to determine if any maximum limit has been exceeded. Also plot points are necessary in order to determine minimum RPM required to obtain guaranteed rated thrust. Other correction factors which are necessary pertain to the emergency fuel flow and cooling air ratio. Other checks of the emergency system consist of acceleration procedures and engine starts. Such steps are necessary in order to determine if the emergency system has the ability to operate properly and take over engine operation in the event the main system fail.

Cooling air ratio is a necessary factor in order to determine if a sufficient

amount of air is being furnished to the hot section parts. If the air ratio is below values outlined by the engine manufacturer, damage could occur to some parts.

#### h. Simulated AB Runs

After completing the necessary performance checks, the engine afterburner system is simulated. The complete afterburner system is subjected to all functions of operation without actual firing. The method used is simply to rout afterburner regulator fuel back to the pump inlet. The ignitor valve will fire, nozzle control will function and afterburner regulator will meter fuel. This system actually is quite practical insofar as all fuel is returned to the inlet supply.

## 2. Testing Techniques

Jet Engine Testing has many problems associated with the techniques encountered using the present manual methods. Some of these problems can be directly associated with human capabilities and reactions during the test cycle. These represent man's inability to cope with the complex situations and the split-second decisions at a speed and with the accuracy required for maintaining a high quality test procedure.

Other problems can be attributed to inaccuracies in the existing mechanical and electrical means of transmitting test data to the test cell personnel from its primary source on the engine. These problems are created because a primary signal in the form of an electrical pulse, voltage or current, pressures, and temperatures must be converted to a mechanical means of display for use by the test operator.

#### a. Standard Tests

A standard test is defined as one in which the test procedure for each type, model, and series engine is conducted in the same manner each time it is conducted, e.g. all data are gathered the same, analyzed the same, and all decisions are made under the same rules without variance. This does not mean that the magnitude of each number in the recorded data will be the same each time, but the manner and intervals at which the recordings are made remain constant.

If an engine is tested and found to be acceptable under one set of ambient conditions, it should also be acceptable when tested under another set of changed ambient conditions. The procedure for testing after overhaul contains the necessary charts and calculations to correct all recordings to a standard day condition; thus, all data should be acceptable under the standard test limits each time it is taken, if it is acceptable at any one of the times.

Even though the testing instructions gives a description of the major procedures to be followed in testing a jet engine, it would become an insurmountable task to specify to the test cell personnel all the exact steps to be taken during the test.

Located at the test facility are many different operators and inspectors, (quality personnel). Because each man is capable of thinking and making individual decisions, he will conduct a jet engine test in a different manner. Because the technical order allows the variations in the manner in which a major test step may be conducted, each operator will not perform each step the same. This situation as well as inconsistencies in the test cell instrumentation will create many different techniques in testing and a possible multiple variations on the acceptance or rejection of an engine under varying conditions.

Not all of the problems associated with the Jet Engine Test can be completely removed by achieving the standard test alone. However, in the process of achieving this standard test many of the "ills" of the present method of testing would have to be eliminated.

The achievement of a standard test can only be realized after correction of the problems in the forthcoming sections.

#### b. Correlation of Test Cells

In most test facilities there are two or more test cells. In order to obtain a standard set of test data on a engine test in one or more of these cells, it is necessary to inter-correlate the cells.

Either a "gold plated" or standard engine that has been tested in the manufacturer's cells is tested in the production cells. This process is commonly known as calibrating a cell. It involves running the standard engine in the production cell, comparing the data gathered with the instrument recordings made in the manufacturer's cells. This will produce a correlation or correction factor to be used with each cell.

The correlation of one cell may require from five to eight hours to complete-longer if trouble is encountered. Trouble is common. Difficulties arise from changing cell ambient conditions (air temperature, humidity, etc.) inaccuracies of data from readout mechanisms, changing of test cell personnel, etc.

The accuracy of the data acquired in a final test phase will be directly dependent upon the degree of accuracy obtained in calculations of each cells correlation factor. Not only are the inaccuracies involved a problem, but there are extra man-hours, fuel costs, and engine wear characteristics incurred.

Rather than correlating every ninety days as is now required,  $\bar{X}$  (average) and R (range or deviation) charts of all instrument reading deviations from those readings



of the production correlator engine would produce cell correction factors. This would allow a constant updating of the cell correction or correlating factor as well as indicating trending abnormalities that may be developing. Using the present techniques of testing jet engines; it is impossible to gather sufficient data, calculate the  $\bar{X}$  and R's of the recordings, and do the correlating.

The data recordings must be gathered and analyzed over a sufficient period of time to detect trending conditions. This usually involves such things as EPR's, EGT's,  $N_1$  and  $N_2$  speeds and their average and range deviations from the standard engine recordings.

Even if it were possible to gather the data, the magnitude of the calculation and analysis is enormous and would require many manhours.

### c. Penalty Runs

It may become desirable after either a major test or a test segment completion, to conduct a penalty run. The penalty run would involve running a small segment of the test, several test segments, or the complete test.

After the completed test and the performance calculations have been made, the engine results could indicate an off specification; thus, requiring the need for a recheck of the calculations and test recordings.

Many times when a borderline situation exists, the inspector will call for the same recheck. Because of the inconsistency in testing methods, calculations, and decisions, the inspector may feel it necessary to repeat a portion of the test in order to gather additional data for analysis, or verification of calculations and recordings. Even when a penalty run is made, conditions may exist (the need for simultaneous readings) that cause the data accuracy to be insufficient, e.g., it is impossible to obtain simultaneous recordings under the manual methods. Because the operator and inspector know of the inconsistencies that exist, several extra minutes or hours along with many extra gallons of fuel may be consumed in conducting the penalty run in order to obtain sufficiently accurate data for a correct test.

## 3. Instrumentation

There are many and varied problems in the instrumentation areas. The sensing elements on most instruments are reliable and accurate. However, the actual readout mechanism is very difficult to keep within the calibration limits. Because most readout mechanisms present problems of nonlinearity in changing from one setting to another, time and manpower must be spent on a periodic (usually monthly or bimonthly) basis to insure accurate calibration. Many times an instrument can become erratic in its reading and the test cell personnel not become aware of it until a new calibration is made. In the mean time many "good" engines have been rejected and many possible "rejects" are flying or in storage.

#### 4. Human Error

Throughout the test procedure recordings are being made on a second timing intervals. Many of these readings should be made simultaneously, but because of the human inability to observe and record on a "split" second basis many of the readings will change by large increments before they can all be recorded. This is especially true during acceleration and deacceleration of the engine.

Because all the instruments are not located at a 90° angle with the eyes of the man making the recording and because many of his recordings are made at a fast rate, it has been found that many recordings have been made with large errors (sometimes a completely gross transition error is made). A 5 lb. pressure error or 5% temperature error is enough in some readings to reject a "good" engine or accept a "bad" engine.

#### 5. Rerun Statistics

If after the sequence of test events, calculations and plotting of data the engine does not perform according to the technical specifications it is not always rejected and sent to the rework area immediately. After a series of checks on the calculations made by himself, gross range errors on readings, or minor detectable instrument error, the inspector will apply his knowledge in conjunction with the trouble shooting points listed in the TO to diagnosing the area of trouble in the engine. These diagnostics will then be sent back with the engine to the rework area (overhaul line).

If he feels that some element of doubt is present in a reading or calculation, portions of the test or the complete test may be performed again. This re-running may consist of re-trimming the engine, re-running the performance runs, or giving an AB function check. Many times the ability to diagnose the problem area relies solely upon the experience and background of the inspector in charge of the test. The majority of the inspection personnel have not gained this type of experience. Because of this inexperience, many of the engines may be re-run or rejected needlessly. If the engine must be re-run several times in order to find the source of trouble, large quantities of time and fuel are consumed.

If we consider the price of a complete overhaul of an engine ranging from \$12,000 to \$15,000, the needless reject of a good engine or the improper diagnostic of an engine for overhaul becomes an expensive waste.

Because of the advance in design of the jet type engine year by year, it becomes a large task to keep the test personnel updated on the new techniques accompanying the advance design engines. During the period of time when the modernization of the cell and training of personnel are being done, many costly errors are made.

If we consider a facility that tests 2,000 engines per year, the annual fuel bill will be approximately \$550,000 per year. It has been estimated that 40 per cent of this fuel bill can be attributed to running reworked engines or performing a portion of a test over again (because of improper readings or calculation errors).

The preceding sections describe in general the testing procedures and some of its existing problems. Do not be "misled" by the seemingly simple test procedures described. There are many things not covered in as minute a detail as possible; also not mentioned are the many splitsecond decisions that must be made during the course of the test and at times when possible malfunctions occur.

## 6. Capacity

In cases of national emergency, or increased workload responsibility, the need for increased test capacity in the high thrust cells could develop into a major production "bottleneck."

Pressure could be relieved in these situations by creating extra shifts of men to handle testing and facilities maintenance; however, the increased utilization of the cells under the present test time and procedures would increase many fold the manhours of maintenance as well as cause the quality of the engines released to the field to be inferior because of this increased pressure.

In either case, the cost of an increased workload under these conditions can become enormous.

## 7. Safety

There are several events that could take place to endanger the lives of personnel working in the test cell while an engine is running. No "concrete" solution will be found to completely remove all these danger areas. The technical order regulations specify where and at what time personnel may be in the cells while the engine is running. Because of unusual circumstances, the rules are many times "bent" to fit the situation. In many of these cases, danger may be at its peak.

Examination of the possible dangers of these situations reveals that there is a possibility of fuel leaks and thus flash fire while trimming. The bleed valve may also dump excess air overboard while decelerating. The force of this air can be enough to knock a man off his feet. There are also dangers from any engine part or accessories not being securely fastened and thus breaking away.

## II. SOLUTION TO THE PROBLEM

### A. Introduction

Some of the problems existing in testing a turbo jet engine have been discussed in the first section of this paper. Not all of the intangible problems were brought out, but inference was made to them.

The forthcoming discussion is submitted as a possible approach to the solution of many of the problems that act as a plague to the efficient and correct testing of a jet engine.

There are many alternatives to the degree of automation that can be applied by the use of a computer in a jet test cell. The primary problem rests on two factors: (1) What degree of control should the system have and (2) Whether the system should be a primary "slave" to the operator or the operator a "slave" to the computer.

The one chosen discusses a completely closed loop operation (In this instance, the running of the test including start-up shut-down via an IBM 1710 Control System, related hardware and any special features). The advantages and disadvantages of operating in a manual and open-loop mode as compared to the chosen approach are discussed.

A great majority of the following information has been derived by working with prospective customers in the jet engine test area; however, due to reasons which will not be discussed, customers' names will not be mentioned.\*

### B. Previous Work

#### 1. Data Logging

One of the first attempts at applying an on-line device for the logging and reduction of engine test data was tried by the U.S. Naval Airforce. A special device for these purposes was built by Gilmore Industries (3) to perform such a function.

The primary design of this device was for gathering piston type engine data. Many of these were later modified to receive data from test cells geared for jet engines.

The data logger was usually located in a prototype cell where certain special test runs could be made.

The data logger was primarily an analog type sensing device. Its primary readouts were instrument faces, graphical x-y plotting and type writer data that had

\* Contact author for further information.

been converted by an analog-to-digital converter to a scaled digital form. The number of channels or sensing and readout elements depended upon the elaborateness of the model ordered. The acceptance of this system was "poor" especially in commercial installations (where a few are found war surplus) where the price/performance ratio was much too great.

This piece of equipment contained the same hinderances as the analog computer does. No logical ability coupled with an "exponential" increase in price for flexibility, plus inadequate readout accuracy. Enough of this type of gear to log data in one cell often times cost as much as the digital computer components to control multi-cells.

## 2. Research and Development Jet Engine Testing

One of the first companies to apply a computer to the role of gathering and reducing test data was Pratt and Whitney. The computer is an IBM 1410 with a special interface (Analog-to-Digital Converter) to take data gathered during tests conducted for research purposes. The system acts as a data monitor. It logs and reduces data only during the time the engine is in the performance run phases. Special instrumentation has been added to detect malfunction of components at high temperatures and fast speeds. After one test has been conducted, the instrument leads are then automatically connected to an engine awaiting test in another cell.

Because the purpose of this system was to do only a data logging and data reduction job, no further effort has been made to perform a close-loop function.

Cases of research and development do not readily adapt themselves to a close-loop operation. There are many times when extraordinary or special tests need to be conducted which would not be compatible with the programs that had been written for test.

There are also under development in the NASA Space Program the adaptation of fast general purpose computer to missile checkout. This program like all other programs in jet engine control is in its infancy.

## 3. Industrial Testing Systems - Discrete Process

Industry has entered an era in which the processing of production and product performance information must be incorporated as a part of the manufacturing operation. As the profit squeeze continues along with the need for increased production, the cost of manufacturing the product must be reduced to maintain or improve the profit position. Much has been done and is being done to reduce the cost of making the product through advances in technology and by automation. However, the costly operation of quality assurance which continues to receive more demanding tasks is not keeping abreast of its production counterpart. To parallel the giant step made in manufacturing through automation, the quality assurance program in industry has

made and continues to make drastic advances through in-process test and inspection systems. Some of the industrial testing applications using industrial process control systems are:

1. Space Vehicles - Analog to digital converter used in logging, reducing and analyzing data on space vehicles in environmental chambers.
2. Potentiometers - Final testing of potentiometers.
3. Automobiles - On-line quality control to determine defects in assembly as they happen.
4. Aerospace Nose Cones - FM Tape playback of data telemetered from missiles.
5. Nuclear Research - On-line recording of information from a spark chamber.
6. Atomic Powered Naval Ships - On-line measurements and computation of shielding experiments.

In general Industrial Testing with control systems controls plant test procedure, analyzes product test data and contributes to production test equipment the capabilities of:

1. Testing dynamically at production speeds.
2. Correlating the test data for each product.
3. Determining the classification of each produced unit based upon specification.
4. Sorting product unit after final test.
5. Storing test data for future analysis.
6. Initiating reports during production runs.
7. Checking and calibrating of test equipment during production runs.
8. Scheduling produced product.
9. Determining critical trends as they develop.<sup>1</sup>

The preceding paragraphs have shown the development of automatic control systems in the continuous process industries and manufacturing operations involving discrete processes. In each case one of the main objectives is increasing the quality of the end product. It can also be clearly seen that automatic testing is not an idea with unproven results but the missing link between production and quality.

1. IBM Application Brief, No. K20-1725

## C. Automatic Jet Engine Test Control System

### 1. System Design Requirements

In the preceding sections of this report the various phases of the actual jet engine test were discussed in moderate detail. These are functions performed by the operator, recorder, and inspector.

The following describes the functions the control system will perform in regard to the various test phases. The functions are necessary to deliver a high quality engine with minimum cost.

#### a. Control of the Independent Variables to Set Up and Sequence Tests

The jet engine control system will select various test phases for an individual type, model and serial number engine-use information gathered from the engine in the test cell, such as pressures, temperatures, flows, etc., and determine appropriate test sequences and procedures taken from the Technical Order to send control signals to the engine in the cell.

By designing the test phases as a series of logical steps, the system will use each test phase as a sub-program and execute the over-all series of sub-programs under control of a master monitor routine.

#### b. Data Acquisition and Control

Each instrument pick-up will be connected to a transducer which will be connected to a multiplexer and terminal unit which will be connected to an analog-to-digital converter. The analog-to-digital converter will provide a digital voltage to the control system main frame. The main frame will scan all instrument leads for each pressure, temperature, flow, etc and convert these by the use of equations into meaningful engineering values. These values will then be used to control the system.

The system will also convert a digital value to an analog voltage for control of the throttle, trimmer and other relay switches in order to control the speed, thrust, fuel flow, and other controllable variables.

#### c. Calculation of Performance Parameters

After gathering all data (instrument readings), one of the test phases will correct all data to a standard day (usually sea level) condition in order that all parameters may be compared against the T. O. limits for trimming and reject status.

d. Operator Guide for Engine Adjustment

Such things as warning messages, trim guides, test status, etc, will be logged for the operator. Any transducer reading will be available upon operator demand.

Any time the engine must be stopped or shut down by the control system, a message will be logged on the typewriter giving the reason, a complete diagnostic, and recommendations for repair or rework.

e. Automatic Instrument Calibration

This can be done by either or both of the following:

1. Comparison of a known standard signal with the transducer output from this signal.
2. Comparison of the transducer output to other related signals. This will, in essence, tell if the signal is abnormal (too high, too low, or fluctuating). From this an automatic calibration can be done. This will insure against catastrophic results from a faulty transducer.

f. Check Calibration of Installed Engine Transducers

Many pick-ups are installed on the engine during dress-down, thermocouples, tachometers, etc. It is possible for one of these to be faulty (disconnection or off specification in the thermocouple not detected during test). By using the calibrate feature, control system abnormalities may be detected before the actual test.

g. Conduct Penalty Runs

After the major test phases are completed and the acceptability of the engine is ascertained, it may be necessary to re-conduct portions of the major test or call upon special penalty run procedures to be executed. This need arises when certain T. O. limits have been exceeded or an engine has been accepted on a marginal condition. This will insure correctness of data and calculations as well as insuring that an out-of-limit condition was not a transient. The ability to automatically select and execute these routines under control of the automatic system will improve both speed and accuracy of the over-all test.

h. Engine Diagnostics

During the running of a sub-portion of the test or after completion of the major test, conditions may arise that will indicate off specifications in the engine



or one of its components. By gathering data at high speeds, using past historical data on engine rejects, failure incidents and rework data, and building a series of logical steps or a mathematical model of certain sections of the engine, it will be possible to determine the exact cause of the abnormality and make recommendations for repair.

There will be a learning process by the system. As more and better data is gathered, the logical model will improve.

The ultimate aim of the system is to furnish complete re-work information to the engine penalty line. In many cases, this will save time and prevent unnecessary rework of an engine.

i. Logging and/or Punch-Out of Test Data and Engine Data

After a major test has been completed, all instrument readings, calculations, and diagnostics remarks will be stored on the disk storage unit. The operator in the cell control room will execute a request to the central control system room via the manual entry control. The control system then will print a completed log or run sheet giving the three items above for each test phase.

The log may be used by the operator to select penalty or re-runs if it appears that a component or recording is marginal to the limit.

Several carbon copies may be produced so that copies may be sent to all authorized personnel.

An engine data plate card will be printed to accompany the engine and a military run data card punched for Quality Analysis.

j. Store Test Programs and Parameters for All Type, Model and Serial Number Engines

The control system will use a mass random access unit for storage of the test parameters and limits for all engine models, types and serial number that will be tested. This type storage insures immediate access to all types of engine programs for complete asynchronous testing and control for the test facility.

Mass storage will allow the system to be open ended for expansion to future cells. By the use of this mass storage, a better and more complete engine diagnostic can be performed (as pointed out in the previous section). The system will be designed to allow the updating of all engine technical orders on a daily basis.

As a secondary function, statistical data will be stored for analysis. By storing summary data, critical trends can be detected early. All causes for rejects or defects can be stored by type, model, and serial number. Summary data will be quickly available upon management request.

k. Detect Emergency and Unsafe Conditions and Take Appropriate Action

The fast instrument scanning speed of the control system permits dangerous trends to be detected in many of the instrument readings and appropriate corrective action to be initiated to prevent occurrence of out-of-limit conditions. In out-of-limit situations, the system will quickly bring the test and engine to a halt to prevent serious damage.

One of the most important things to consider when designing the actual control system is achieving a high degree of reliability. Two types of failures can occur. The failures and corrective actions are:

Type I - Transient Failures

These are internal system transmission errors and occur on a transient basis. In this case, the system will record the failure and try twice more to perform the operation. The recording will be used by maintenance engineers for regular preventive maintenance (once per week). A transient type error will usually be eliminated in three attempts.

Type II - Complete Component Failure

In this case, the system will try to by-pass the bad component switching to manual control or bring the engine to a safe stop. A by-pass procedure will be incorporated for emergency action.

l. Quality Analysis

- 1) The system will use store data to perform reliability calculations for engine and individual components.
- 2) The quality analysis will produce data assurance for a better test engine.

m. Production

- 1) Scheduling - Using advanced techniques such as linear programming, a master plan will be prepared for scheduling the cells.
- 2) Planning - Better methods of machine and manpower utilization can be prepared.

2. Control System and Interface Description

To approach the problem of determining the necessary hardware, one must keep three factors in mind. They are (1) design functions as determined in the meeting of section 1 (2) instrumentation-present and future, (3) and layout of the basic test cell.

If we notice the basic test cell layout as shown in illustration 1, it shows the location of the control room as being between two test cells. If there are more than two cells (there are usually several more) then it is logical there will be two or more control rooms. Because a typical control system will control more than two cells, it will be necessary to locate the computer in either a remote location or in the rear of one of the test cells. When this is done, there arises necessity for remote communication devices.

Attention should be drawn to the design function to operate in conjunction with this communications device. Whether the operator or the central computer system is the "slave". It will be necessary to place a device for the operator (inspector, etc.) to select the particular test function he wishes to perform. It will also be necessary for him to get return information from the instrument readings, pertinent calculations and emergency or troubleshooting messages.

Many times it will be necessary for the test cell foreman to have information concerning phases of test of engines in each cell in order to coordinate the overall movement of engines in test. He will also need access to stored statistical information pertaining to reject, re-run and other engine test functions. Many times upper level management will inquire of the cell foreman on these statistics. Things that could be available on an inquiry basis would be:

1. Number of rejects/month on a certain model number.
2. Major cause of rejects.
3. What was done for correction.
4. Ranges and standard deviations from set standard operating limits.
5. Etc.

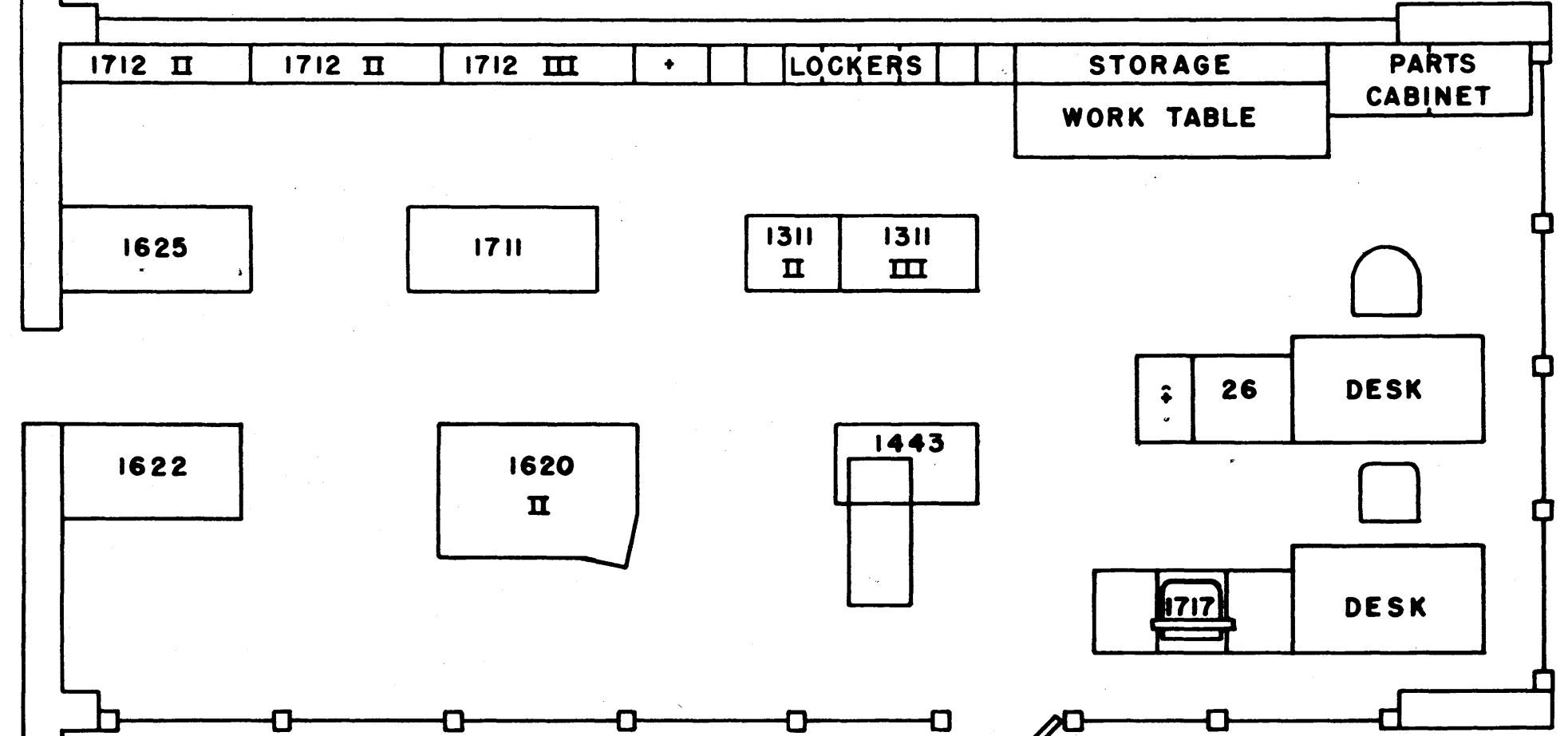
Using the system described in illustration 2 and treating the requirement as 8 test cells, each component and its function will be discussed.

a. Central Computer (1620 Model II)

Because of the speed needed to accomplish the sampling of the necessary instrument leads, making all necessary calculations, actually sending output signals, for control and receiving feedback input signals for correction, the 1620 Model II with a 60,000 position memory was chosen. The 1620 as the heart of the 1710 Control System contains the necessary machine instructions and programming systems (Executive System-e.g., monitor) to operate in conjunction with an asynchronous test system design.

b. Auxiliary Storage (2-1311 Disc Files)

Even though the majority of the skeleton test functions are the same for all engines, there still remains different test parameters for each engine model. Each of these parameters must be stored for immediate access. Because no central computer memory would be large enough to contain all test program phases, these must



1712 II

1712 II

1712 III

+

LOCKERS

STORAGE

WORK TABLE

PARTS  
CABINET

1625

1711

1311  
II

1311  
III

1622

1620  
II

1443

26

DESK

1717

DESK

+ 1712 IV, V, VI, VIII  
: CARD FILE

CONTROL ROOM

SCALE: 1/4" = 1'

207

also be stored for immediate access as they are called by the skeleton control program. Also contained in auxiliary storage would be necessary diagnostic routines available upon request, as well as emergency limit and correction routines.

A second disc file would be used as in intermediate store area for input/output information, if all input/output devices are busy and they would be available to store quality and production control data gained as a by-product of each test. This data would be available upon inquiry from management.

c. Interface Equipment (1711 and 1712's)

In order to attach all necessary points for 8 cell (see Appendix B) and convert the analog (electrical) signal into a digital form in a sufficient time period, the analog to digital converter (1711) has the ability to convert 200 points/second. In order to handle all necessary analog input points, analog output points, contact indicating and operating relays for an eight (8) test cell facility, it is necessary to have three (3) multiplexing and terminal units (1712's) to the system.

d. Test Cell Input/Output Gear (1713, 1715, 1717)

Located adjacent to each instrument control panel will be an IBM 1713, 1715 and 1717.

The operator will have the option with the IBM 1713 manual entry device-through a set of coded instructions-to dial in either the command for a complete test or portions of a test. The command will be dialed through the use of twelve (12) rotary knobs with zero (0) to nine (9) selection ability.

An enter key will be hit, the information will go via the SIOC channel and interrupt the computer, the computer will read the rotary knobs and start the processing.

e. Interface

All instruments that furnish an electric signal of a standard form will be sent via shielded cabling to the 1712 multiplexing unit, all non standard (pulsed, etc) and pressure type signals will be transformed via transducers (in the test cell control room) to an electrical form and sent to the central system complex.

All pickup signals from the engine are easily adjusted to the standard 1710 signals; however, more specialized servos must be bought or designed to control the throttle and trimming mechanism. There are several types of stepping motors or feedback systems on the market today that can handle these tasks.

All existing cell instruments will remain intact as manual back-up for the system. Through a specially designed panel, the operator will be allowed the option of switching to either automatic or manual system at any time.

### III. ECONOMIC JUSTIFICATIONS

The justifications for considering the "Control Systems Approach to Jet Engine Testing" can be broken into tangible, intangible and possible savings categories. The justifications can vary depending upon the application. Some of each are listed as follows:

#### A. Tangible

##### 1. Increased Engine Throughput:

This can be accomplished by

- a. Simplifying the testing procedure.
- b. Decreasing delay in such things as trimming and shakedown.
- c. Operator Guide Print-Out for prompt emergency and testing actions.

The best time estimate for engine throughput with no major hindrances is 5 hours 55 minutes. As previously mentioned, the average throughput is approximately eight (8) hours for an engine with time running up to twelve (12) hours if there are several re-runs or persistent trouble exists.

The control system would increase the capability of the cells to take on added workload without added facilities. This need would arise in wartime emergency for federal customers and with added contractual obligations for both commercial and federal.

##### 2. Reduced Manpower Requirement/Engine

This would free inspection and operating personnel for a greater engine throughput. One operator would be substantial for testing procedures, where the present system utilizes an inspector and two (2) shops or production personnel.

##### 3. Avoiding Re-Run of Engines

- a. By eliminating bad instrument calibration--erroneous transducer signal.
- b. Bad instrument reading--can be eliminated. The signal will originate completely at the transducer and eliminate the nonlinearity of the instrument read-out mechanism. Operator error in reading will also be eliminated, e.g., simultaneous reading of instruments.

#### 4. Decreased Fuel Costs

This can be saved with automatic trim procedures and avoidances of excess penalty runs.

#### 5. Decreased Calibration Costs

By automatically calibrating the transducers the computer will give a correction or tare factor for the back-up reading devices. The time between calibrations will decrease. The maintenance costs will correspondingly decrease.

### B. Intangibles

#### 1. Better Engine Quality

- a. Better checked out engines through more certain detection of off-specification units.
- b. Simultaneous recording of the instruments, thus insuring proper data for checking limit parameters.
- c. Consistent methods of testing, thus insuring proper acceptance or rejection of an engine.

#### 2. Decreased Re-Work Costs Through Better Diagnostics

As mentioned previously, the system with its on-line mass storage can furnish pinpoint diagnostics to eliminate complete overhaul for minor abnormalities or defects.

#### 3. Data Assurance

This assurance can be derived through getting simultaneous readings or reliability in instrument calibration and will result in better customer (the pilot or branch of the armed services) satisfaction.

#### 4. Increased Safety for Personnel, Equipment and Property

#### 5. Increased Readiness Program on First Line Aircraft

#### 6. Reduced Paperwork Handling

Complete and accurate unit performance logging. Here we have better customer satisfaction through hard copy records. Diagnostics are automatically printed to be sent back for re-work.

#### IV. SUMMARY AND CONCLUSIONS

This paper has attempted to discuss one particular approach to the application of a digital computer to the closed loop control of a jet engine test cell.

As has been pointed out, there are many approaches to consider in designing a system for a particular application. In summary, the things that must be considered are repeated and listed as follows:

1. Degree of Control Desired-Open or Closed Loop
2. How Much the Operator is "Slave" to the Computer or Vice Versa
3. How Many Test Cells Must Be Controlled Simultaneously
4. The Functions that the Customer Wishes the System to Perform

The justifications for a control system can be varied, depending upon what the customer wishes to accomplish; however, the "state of the art" of automatic control in jet engine testing is in its infancy and there are many justifications in all cases.

It should be pointed out that the general approaches and ideas used are applicable to many other industries and are not limited to jet engine testing.



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**GOODYEAR**  
**GOODYEAR AEROSPACE**  
CORPORATION  
ARIZONA DIVISION  
LITCHFIELD PARK, ARIZONA

GENERALIZED FILTER NETWORK

A/C STEADY STATE ANALYSIS PROGRAM

by

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AAP-18911

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## GENERALIZED FILTER NETWORK

### A/C STEADY STATE ANALYSIS PROGRAM

This program has been written to make possible comprehensive surveys of theoretical filter designs. It opens up a more sophisticated range of filters to theoretical consideration and evaluation. The program input is general enough that almost any filter network consisting of cascaded inverted-L or symmetrical lattice sections may be handled easily.

The minimum machine requirements are a 1620 with 40 K core storage, auto divide, and indirect addressing. The source language is Fortran II. There are 6 subprograms plus the mainline program.

Filter design has been speeded in recent years with the advent of tables of normalized low-pass filter element values\*. Even if these tables are used, this program allows the designer to compute the effects of component tolerances, finite Q's, and mismatched terminations. These introductory remarks have centered around filter design, but it will be apparent that the program is useful for analyzing any RLC network, e.g., amplitude or phase equalizers.

The filter designer needs to know how a proposed design will perform over a particular range of frequencies before making recommendations to those who will implement the design. Manual calculations of the desired performance parameters over a range of frequencies can be quite tedious and are highly subject to human error. The problem is complicated because theoretical calculations do not always reflect the actual performance, particularly at higher frequencies due to stray reactances. The net result frequently is that a minimum number of proposed designs are evaluated at a few pertinent frequencies from which data plots of parameters versus frequency are made. Curve definition is rarely good, making it virtually impossible to compare proposed filter designs on important fine points.

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\*For example, see 1964 Microwave Engineers' Handbook, pp. 91-95, Horizon House-Microwave, Inc., Dedham, Mass.

This program will evaluate theoretical performance of filters over as broad a range of frequencies as desired giving precise values of the performance criteria. The net result is excellent definition of a greater number of performance curves in much less time. The designer is freed of computational drudgery and may consider more sophisticated designs without concern for the difficulties of manual evaluation.

The parameters computed by the program are insertion loss, impedance, input voltage phase angle (phase shift), and time delay (envelope or group delay). Insertion loss is defined as that loss resulting from inserting a given four terminal filter between the input and output terminals of a network where the output is non-reactive and the input can be a complex impedance driven by a voltage signal source. The phase angle of the input voltage is referenced to the voltage across the output terminals whose magnitude and phase are arbitrary. The impedance calculated is at the input of the filter looking toward the output and is in rectangular form.

The program is based upon a report titled "Electronic Digital Computer Analysis of Cascaded Networks" by R. H. Tuznik and D. H. Wood of Westinghouse Electric Corporation. The report presents a method of defining in general terms a filter network and then, through recursion techniques, computing the impedance and voltage at the filter input for any desired frequency.

The program can handle networks having either basic ladder or symmetrical lattice sections. These sections consist of  $Z_A$  and  $Z_B$  basic section impedances as in Figure 1.

A maximum of four different sections in cascade can be handled. This is a limitation imposed by core storage. Note that the program will process a group of from one to four cascaded sections any desired number of times. Thus it is possible to analyze a filter consisting of four repeated groups of four cascaded sections by setting NCYCL, a program input variable, to four.

A section impedance  $Z_A$  or  $Z_B$  as in Figure 1 can be as simple as a single element or as complex as in Figure 2.

A basic section impedance may consist of from 1 to 3 orders each of which may be as complex as above or as simple as a single element. The maximum number of parallel branches in an order as well as the maximum number of orders again are limited by available core storage.

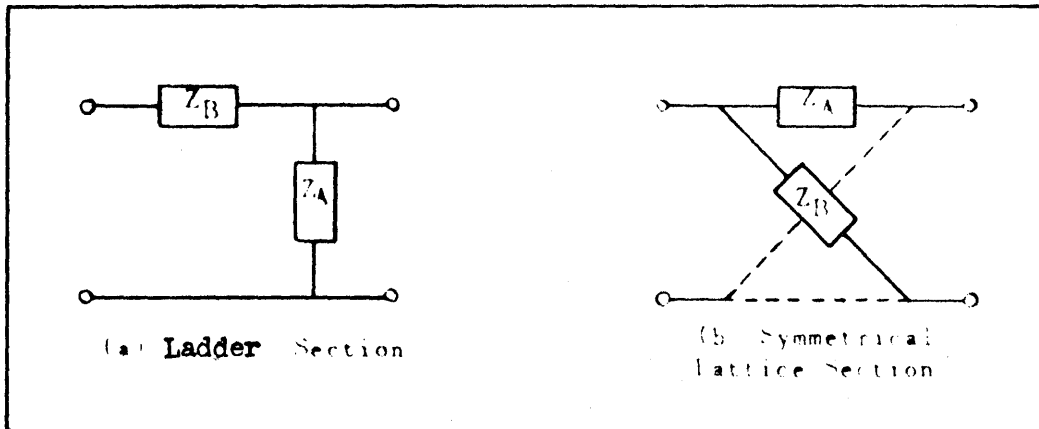


FIGURE 1

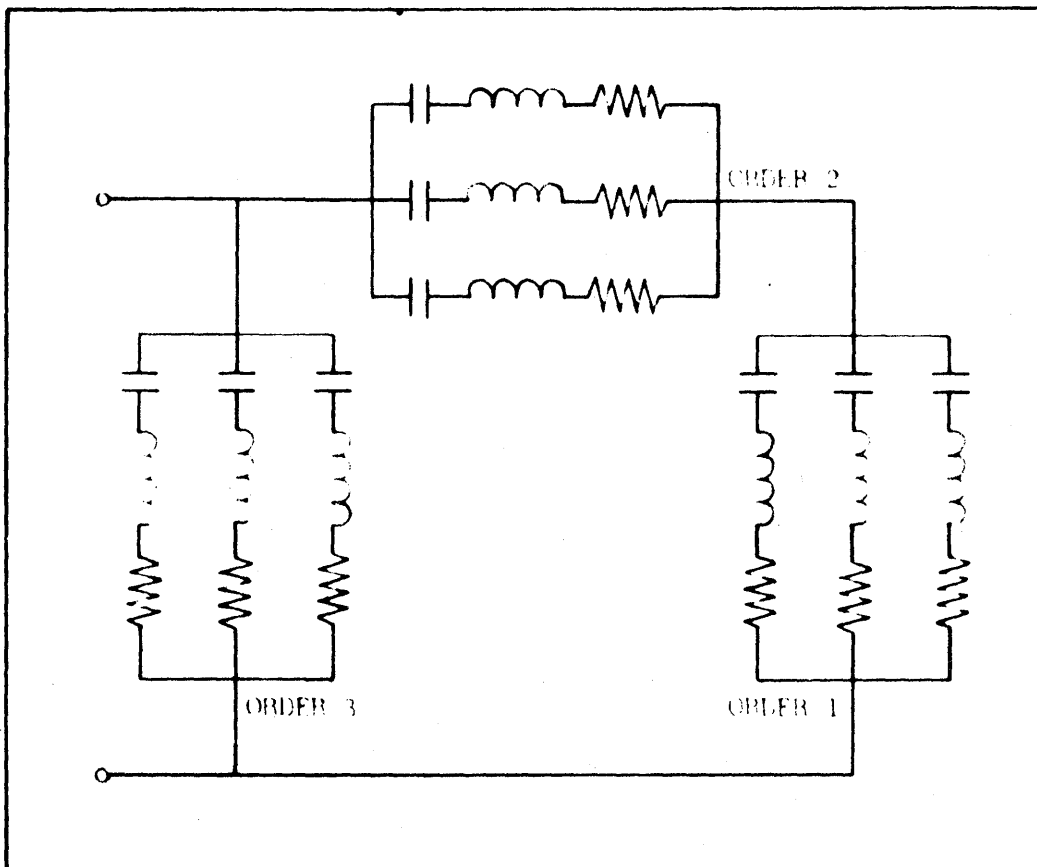


FIGURE 2 - BASIC SECTION IMPEDANCE

The program positions the filter to be analyzed between an input and output termination as in Figure 3.

The input termination may be as complex as shown or a single element. The output termination  $R_o$  must be a pure resistance.  $V_o$  is an input variable and usually is set equal to 1 volt at zero phase. The program computes  $V_i$ , the phase angle of which is used to determine time delay.

The designer may desire to analyze a proposed filter for a frequency sensitive network. The frequency sensitive network must be representable in the form shown in Figure 3 as the input termination. Another way is to include the frequency sensitive network as part of the filter with a single resistance as the input termination. The first way has the advantage of allowing the NCYCLE variable to be other than one.

The program begins by reading cards with the following information:

1. Frequency range including maximum, minimum, and increment.
2. Output termination resistance and voltage.
3. Number and type of cascaded sections in filter.
4. Number of times filter sections are to be cycled.
5. Filter section element values.
6. Input termination element values.

The program immediately punches out cards with these values. Then computation begins at the output termination and progresses section-by-section toward the input termination. Complex impedance and voltage necessary to produce the output termination conditions are calculated at the input of each filter section and so on recursively to the input termination.

The recursion equations programmed are for inverted-L or ladder sections and symmetrical lattice sections. These are in separate subprograms named LADDER and LATTICE respectively. These will be outlined in Appendix A. The program could be expanded to handle other types of basic sections such as the tee, bridged-tee, pi, etc., by writing additional recursion subprograms. Certain of these types of sections can be rearranged into ladder sections for evaluation with the existing program.

Examination of Figure 1 suggested a programming short cut. Both types of basic sections consist of two-terminal basic section impedances as in Figure 2. A single combination of two subprograms is used to compute these impedances.

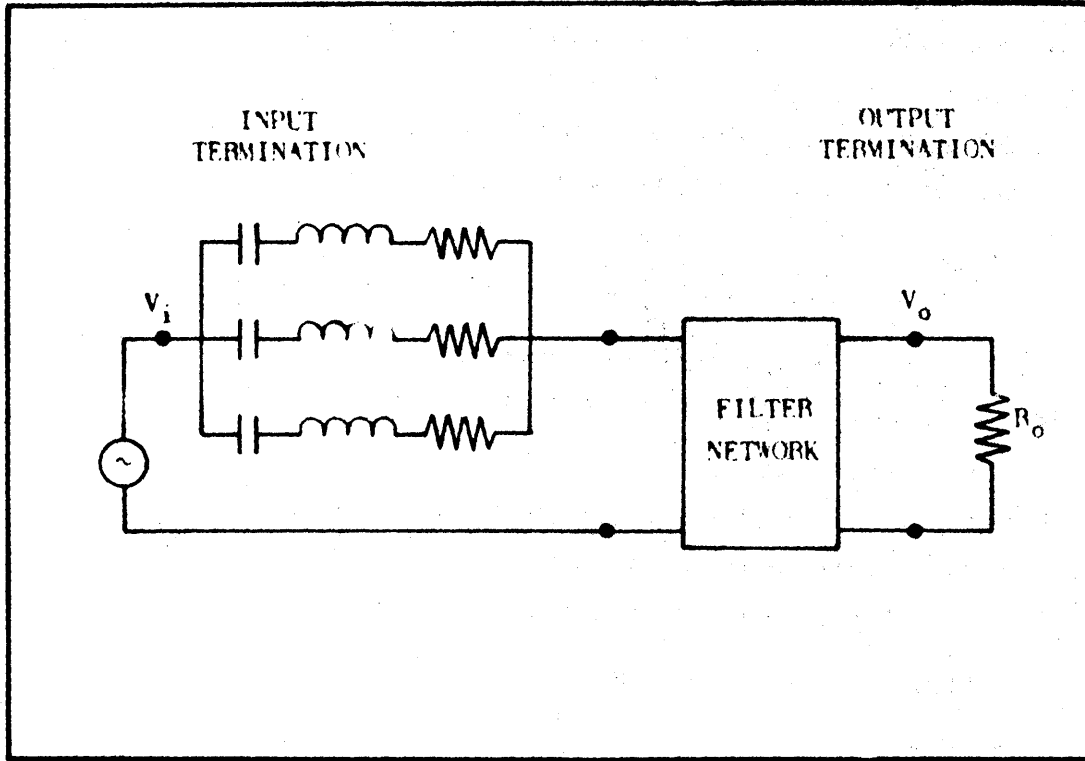


FIGURE 3 - FILTER NETWORK WITH TERMINATIONS



The first, DRPTZ (Driving Point Z) computes the immittances of an order. The second, ORDER, combines the immittances into the basic section impedance. When the first order of a particular basic section impedance  $Z_A$  has been processed, the program begins with the first order of the second basic section impedance  $Z_B$ . When all orders of a particular section have been processed the program shifts control to the appropriate recursion subprogram LADDER or LATTICE. All basic sections in a particular filter must be either inverted-L or symmetrical lattice type.

#### DETAILED DESCRIPTION OF PROGRAM INPUT DATA

Program input consists of four categories of cards which must be input in the order following for each network to be analyzed.

1. Title Card.
2. Frequency and Output Termination Card.
3. Basic Section Data Cards
4. Input Termination Cards

Note that jobs may be stacked.

#### Title Card

Any 80 alphanumeric characters for identification of output.

#### Frequency and Output Termination Card (F11.5 format)

##### Cols.

- |         |   |
|---------|---|
| 1 - 11  | FMIN - Minimum frequency in megacycles.                     |
| 12 - 22 | FDEL - Frequency increment in megacycles.                   |
| 23 - 33 | FMAX - Maximum frequency in megacycles.                     |
| 34 - 44 | ZOR - Output termination resistance in ohms.                |
| 45 - 55 | VOR - Output termination voltage (real part) in volts.      |
| 56 - 66 | VOI - Output termination voltage (imaginary part) in volts. |

#### Basic Section Data

This data describes the basic filter sections. Data should be input section by section beginning at the section nearest the Output Termination.

## Inverted-L (Ladder) Sections

### Preliminary Card (I5 format)

<u>Col.</u>	
5	NSECT - Number of filter sections (Max. of four).
10	LETTR - Must be two.
15	NCYCL - Number of times basic sections are used.

### Order Card (I5 format)

<u>Col.</u>	
5	IORD (K) - Maximum number of orders in $Z_A$ or $Z_B$ impedance of $K^{\text{th}}$ section. Maximum of three.

### Resonator Card (I5 format)

One card for each pair of  $Z_A$  and  $Z_B$  orders. This card must immediately precede the  $Z_A$  and  $Z_B$  parameter cards.

<u>Col.</u>	
5	MP - Number of resonators (parallel branches) in $Z_A$ .
10	MQ - Number of resonators (parallel branches) in $Z_B$ .

### $Z_A$ Parameter Cards (El4.8 format)

One card per MP resonator specified on preceding resonator card.

<u>Cols.</u>	
1 - 14	C - Capacitance in Farads.
15 - 28	AL - Inductance in Henries.
29 - 42	R - Resistance in Ohms.

### $Z_B$ Parameter Cards (El4.8 format)

One card per MQ resonator specified on preceding resonator card.

<u>Cols.</u>	
1 - 14	C - Capacitance in Farads.
15 - 28	AL - Inductance in Henries.
29 - 42	R - Resistance in Ohms.

## Symmetrical Lattice Sections

This data is identical to Inverted-L data except that LETTR on the preliminary card in Col. 10 must be 1 and resistive pi pad parameters must be input. The card for the pi pad between the output termination and the first basic section must follow the preliminary card. Cards for

pi pads following basic sections must follow the Order Card.

Pi Pad Card (E14.8 format)

Cols.

1 - 14 RA - Resistance of pad shunt arms in ohms.

15 - 28 RB - Resistance of pad series arm in ohms.

SAMPLE PROBLEMS

Three examples will be given, one of which will be quite simple and the other two slightly advanced.

Sample Problem 1

This problem will illustrate the input for a simple inverted-L (LADDER) filter network.

Given the two-section ladder filter in Figure 4 with element values as below, calculate filter parameters from 1.0 mc to 10 mc with a step of 1.0 mc.

Use output voltage of 1.0 volts at zero angle as reference.

Output Termination

$$R_o = 50 \text{ ohms.}$$

First Ladder Section

$$R_1 = 1000 \text{ ohms}$$

$$C_1 = 1.0 \mu\text{f}$$

$$L_1 = 2.0 \mu\text{h}$$

Second Ladder Section

$$R_2 = 100,000 \text{ ohms}$$

$$C_2 = .01 \mu\text{f}$$

$$L_2 = 1.0 \mu\text{h}$$

Input Termination Section

$$R_T = 50. \text{ ohms}$$

Preparation of Input Data - Problem 1

Title Card

Sample Problem 1 - Network Analysis Program - Date

### Frequency and Output Termination Card

<u>Cols.</u>	<u>Value</u>	
1 - 11	1.0	Minimum Frequency (mc) FMIN.
12 - 22	1.0	Frequency Increment (mc) FDEL.
23 - 33	10.0	Maximum Frequency (mc) FMAX.
34 - 44	50.0	Output Termination Resistance (ohms) ZOR
45 - 55	1.0	Output Termination Voltage Real Part (volts) VOR
56 - 66	0.0	Output Termination Voltage Imag. Part (volts) VOI

### Basic Ladder Section Cards

#### Preliminary Card

<u>Col.</u>	<u>Value</u>	
5	2	Number of filter sections
10	2	Ladder section signal
15	1	Number of times basic sections are used

#### Order Card (First Section)

<u>Col.</u>	<u>Value</u>	
5	1	Maximum Number of Orders in First Section

#### Resonator Card (First Order)

<u>Col.</u>	<u>Value</u>	
5	1	The only order of the $Z_A$ impedance of the first section has a single resonator.
10	1	The only order of the $Z_B$ impedance of the first section has a single resonator.

#### $Z_A$ Parameters Card

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.10000000E-05	Capacitance in Farads - $C_1$ .
15 - 28	+.00000000E-99	Inductance in Henries
29 - 42	+.10000000E+04	Resistance in Ohms - $R_1$ .

NOTE: There is only one  $Z_A$  parameter card for the first section because there is only one resonator.

Z<sub>B</sub> Parameter Card

Cols.      Value

1 - 14    +.00000000E-99

Capacitance in Farads

15 - 28    +.20000000E-05

Inductance in Henries - L<sub>1</sub>.

Order Card (Second Section)

Col      Value

5            1

Maximum number of orders in second section.

Resonator Card (First Order)

Col .      Value

5            2

There are 2 resonators in the only order of the Z<sub>A</sub> impedance of the second section.

10          1

There is 1 resonator in the only order of the Z<sub>B</sub> impedance of the second section.

Z<sub>A</sub> Parameter Cards

Cols.      Value

1 - 14    +.10000000E-07

Capacitance in Farads - C<sub>2</sub>.

Cols.      Value

1 - 14    +.00000000E-99

Capacitance in Farads

15 - 28    +.00000000E-99

Inductance in Henries

29 - 42    +.10000000E+06

Resistance in Ohms - R<sub>2</sub>.

Z<sub>B</sub> Parameter Card

Cols.      Value

1 - 14    +.00000000E-99

Capacitance in Farads

15 - 28    +.10000000E-05

Inductance in Henries - L<sub>2</sub>.

Order Card (Input Termination)

Col.      Value

5            1

There is a single order in the input termination.

Resonator Card (First Order)

Col.            Value

5                    1

There is one resonator in the input termination and it is treated as a  $Z_A$  type impedance.

10                   0

No  $Z_B$  resonators.

$Z_A$  Parameter Card

Cols.            Value

1 - 14    +.00000000E-99

Capacitance in Farads

15 - 28    +.00000000E-99

Inductance in Henries

29 - 42    +.50000000E+02

Resistance in Ohms -  $R_T$

This completes the input data cards for sample problem 1. A listing of these cards is found on Page 22. A listing of the output is found on Page 23.

Sample Problem 2

This problem will illustrate the input for a two section symmetrical lattice filter network.

Given the network in Figure 5 with two symmetrical lattice sections separated by resistive pi pads:

Output Termination

$R_o$     -    100 ohms

First Pi Pad

$RA_o$     -    10000 ohms

$RB_o$     -    5 ohms

First Section

$C_1$     -    30  $\mu$ f

$C_2$     -    .02  $\mu$ f

$L_3$     -    .1  $\mu$ h

Second Pi Pad

$RA_1$     -    10000 ohms

$RB_1$     -    5 ohms

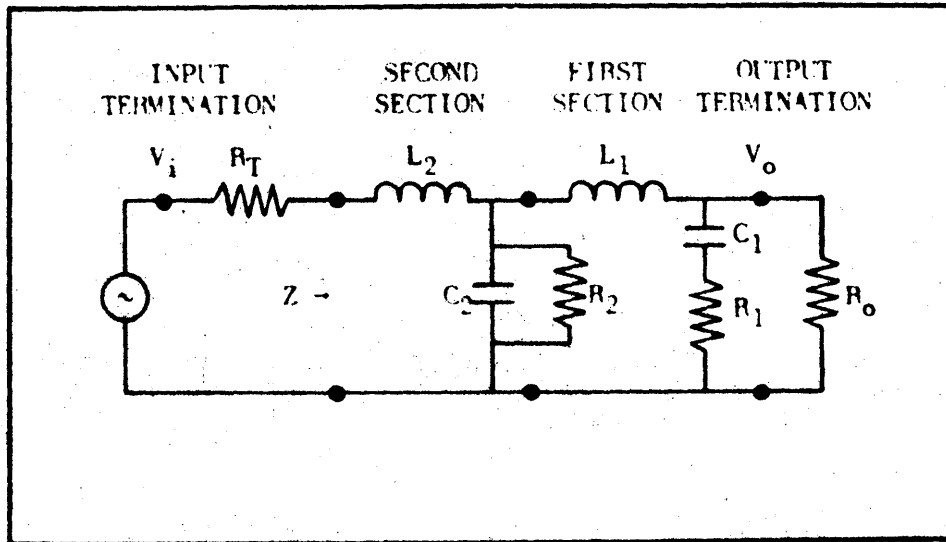


FIGURE 4 - SAMPLE PROBLEM 1

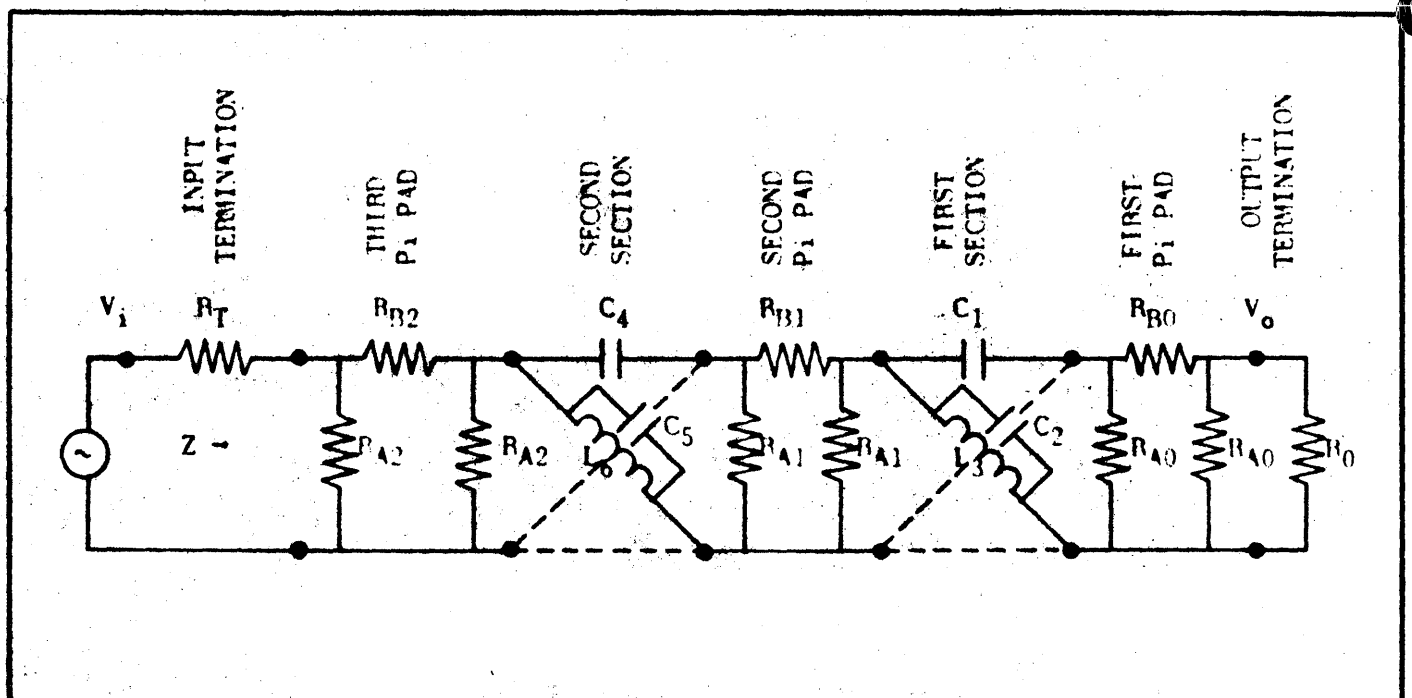


FIGURE 5 - SAMPLE PROBLEM 2

Second Section

$C_4$  - 300  $\mu\text{f}$   
 $C_5$  - .4  $\mu\text{f}$   
 $L_6$  - .15  $\mu\text{h}$

Third Pi Pad

$RA_2$  - 10000 ohms  
 $RB_2$  - 5 ohms

Input Termination

$R_T$  - 100 ohms

The system in Figure 5 is to be evaluated over a range of frequencies from 50 to 100 megacycles with 2 megacycle steps. The program always computes the same type of information, namely impedance, phase angle of  $V_1$ , insertion loss, and time delay.

Preparation of Input Data - Problem 2

Title Card

Sample Problem 2, Network Analysis Program - Date.

Frequency and Output Termination Card

<u>Cols.</u>	<u>Value</u>	
1 - 11	50.	Min. Freq. in Mc.
12 - 22	2.	Delta Freq. in Mc.
23 - 33	100.	Max. Freq. in Mc.
34 - 44	100.	$R_o$ in Ohms.
45 - 55	1.0	$V_o$ in Volts (real part).
56 - 66	0.0	$V_o$ in Volts (imag. part).

Basic Lattice Section Cards

Preliminary Card

<u>Cols.</u>	<u>Value</u>	
5	2	Number of filter sections.
10	1	Lattice section signal.
15	1	Number of times basic sections are used.



First Pi Pad Card

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.10000000 E+05	RA <sub>0</sub> in ohms
15 - 28	+.50000000E +01	RB <sub>0</sub> in ohms

Order Card (First Section)

<u>Cols.</u>	<u>Value</u>	
5	1	Maximum Number of Orders in Section

Second Pi Pad Card

Same as First Pi Pad Card

Resonator Card (First Order)

<u>Col</u>	<u>Value</u>	
5	1	A single resonator in Z <sub>A</sub> .
10	2	Two resonators in Z <sub>B</sub> .

Z<sub>A</sub> Parameter Card

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.30000000 E-10	Capacitance in Farads - C <sub>1</sub> .

Z<sub>B</sub> Parameter Cards

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.20000000 E-07	Capacitance in Farads - C <sub>2</sub> .

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.00000000 E-99	Capacitance in Farads.
15 - 28	+.10000000 E-06	Inductance in Henries - L <sub>3</sub> .

Order Card (Second Section)

<u>Cols.</u>	<u>Value</u>	
5	1	Max. Number of Orders in Section.

Third Pi Pad Card

Same as First Pi Pad Card.

Resonator Card (First Order)

<u>Cols.</u>	<u>Value</u>	
5	1	A single resonator in $Z_A$ .
10	2	Two resonators in $Z_B$ .

$Z_A$  Parameter Card

<u>Col.</u>	<u>Value</u>	
1 - 14	+.30000000 E-09	Capacitance in Farads - $C_4$ .

$Z_B$  Parameter Cards

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.40000000 E-06	Capacitance in Farads - $C_5$ .

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.00000000 E-99	Capacitance in Farads.
15 - 28	+.15000000 E-06	Inductance in Henries - $L_6$ .

Order Card (Input Termination)

<u>Col.</u>	<u>Value</u>	
5	1	Maximum of One Order allowed in Input Termination.

Resonator Card (First Order)

<u>Col.</u>	<u>Value</u>	
5	1	There is one $Z_A$ Resonator in the Input Termination.
10	0	The Input Termination Cannot Have A $Z_B$ Resonator.

$Z_A$  Parameter Card

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.00000000 E-99	Capacitance in Farads.
15 - 28	+.00000000 E-99	Inductance in Henries.
29 - 42	+.10000000 E+03	Resistance in Ohms $R_T$ .

This completes the input for sample problem 2. A listing of these cards is on page 24. The resulting output is on page 25.

### Sample Problem 3

This problem will illustrate three features not displayed by the previous problems, namely the use of NCYCL for multiple usage of a group of cascaded sections, the use of a section having 3 orders, and a complex input termination.

This filter has a repeated group of two sections, that is, the First and Third Sections are identical, also the Second and Fourth Sections. The block labeled Z3 represents the two-terminal three-order impedance shown in Figure 7. The frequency range of interest is 5 to 100 Mc in steps of 5 Mc.

Element values are as follow:

#### Output Termination

$R_0$  - 100 Ohms.

#### First Section

$C_1$  - 1  $\mu$ f  
 $R_2$  - 1000 Ohms.

#### Second Section

$C_3$  - 1  $\mu$ f  
 $R_4$  - 1 Megohm  
 $R_5$  - 1000 Ohms  
 $L_6$  - .2  $\mu$ h  
 $C_7$  - .3  $\mu$ f  
 $C_8$  - 1.  $\mu$ f  
 $R_g$  - 100 Ohms  
 $L_{10}$  - 1  $\mu$ h  
 $R_{11}$  - 50 Ohms  
 $C_{12}$  - .1  $\mu$ f  
 $L_{13}$  - .2  $\mu$ h  
 $R_{14}$  - 100 Ohms  
 $L_{15}$  - .01  $\mu$ h  
 $R_{16}$  - 100 Ohms  
 $C_{17}$  - .005  $\mu$ f  
 $L_{18}$  - .5  $\mu$ h  
 $R_{19}$  - 100 Ohms  
 $R_{20}$  - .1 Megohms  
 $C_{21}$  - 1  $\mu$ f

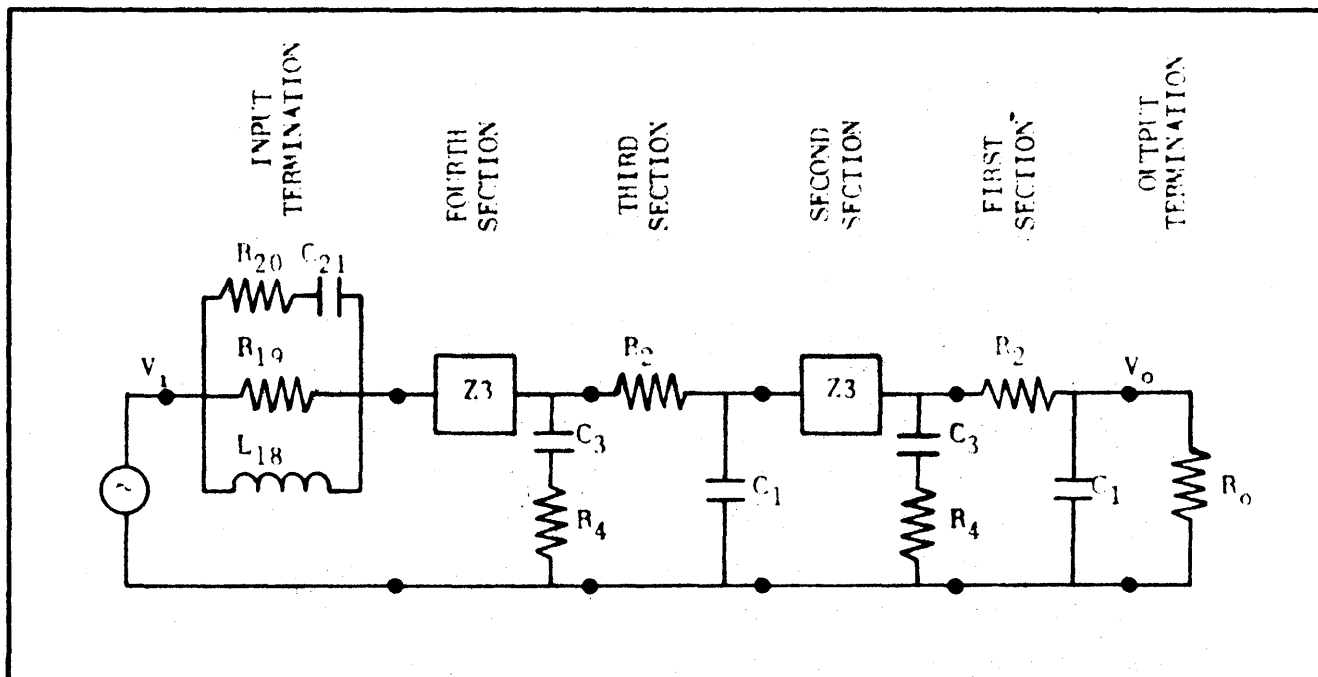


FIGURE 6 - SAMPLE PROBLEM 3

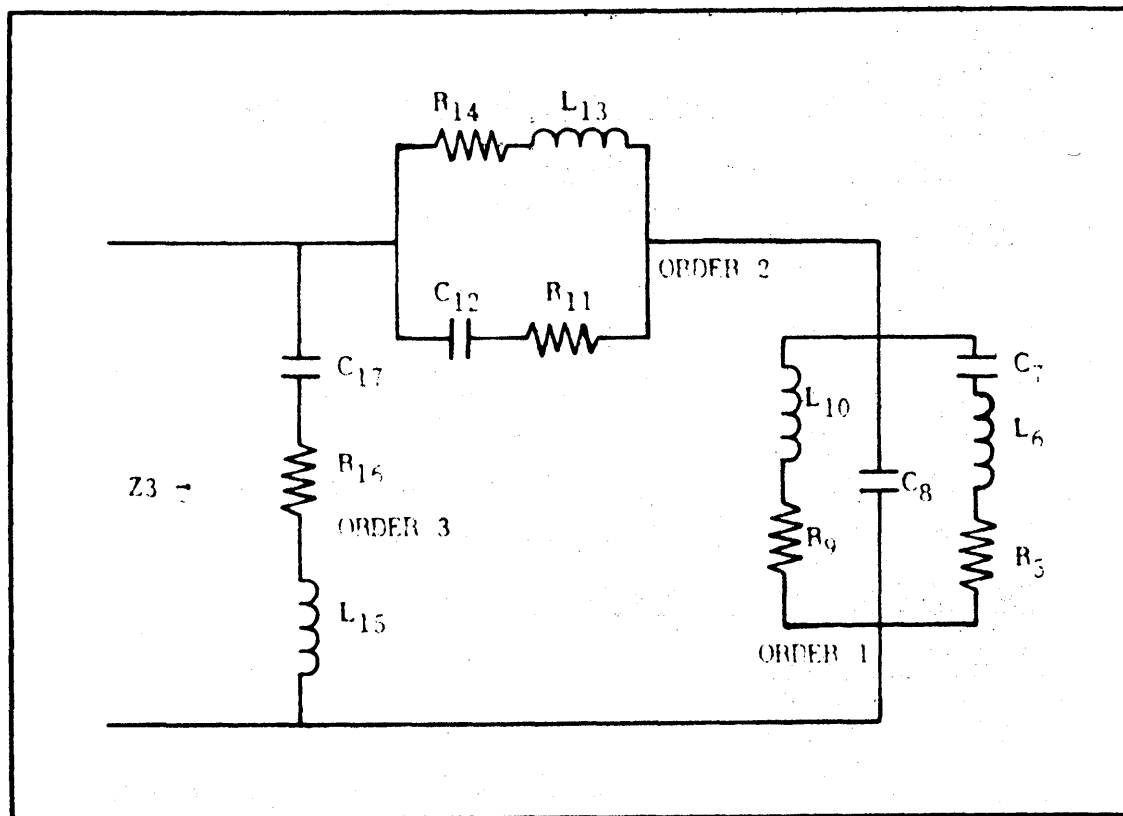


FIGURE 7 -  $Z_3$  IMPEDANCE IN SAMPLE PROBLEM

### Sample Problem 3 (Continued)

For simplicity the output voltage will be assumed 1.0 volts at zero phase as before. The input to this problem is fairly complex but the fact that the first two sections are repeated in the next two sections means that only the first two sections must be input. Setting the NCYCL input variable equal to two (2) causes the program to process the two-section group twice before starting on the input termination. NCYCL can be set to any positive integer (less than 999). This feature can be used with either inverted -L or symmetrical lattice sections. As the program is written, the group of sections which is repeated can include a maximum of four (4) different sections. Only a single group can be in any one filter and the repeated group cannot be preceded or followed by any other sections excepting the input and output terminations. The three-order impedance Z<sub>3</sub> in Figure 7 illustrates the complexity permitted by the program for all of the Z<sub>A</sub> or Z<sub>B</sub> type impedances in any or all basic sections. The complex input termination could have included three elements in each parallel branch or resonator but cannot exceed a single order.

### Data Preparation for Sample Problem 3

#### Title Card

Sample Problem 3, Network Analysis Program - Date

#### Frequency and Output Termination Card

<u>Cols.</u>	<u>Value</u>	
1 - 11	5.0	Min. Freq. Mc
12 - 22	5.0	Delta Freq. Mc
23 - 33	100.0	Max. Freq. Mc
34 - 44	100.0	R <sub>o</sub> in Ohms
45 - 55	1.0	V <sub>o</sub> (Real Part) in Volt
56 - 66	0.0	V <sub>o</sub> (Imag. Part) in Volt

#### Basic Ladder Section Data Cards

##### Preliminary Card

<u>Col</u>	<u>Value</u>	
5	2	There are two different sections to be repeated.
10	2	Ladder Section Signal
15	2	Cycle the two sections twice.

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Order Card (First Section)

<u>Col.</u>	<u>Value</u>	
5	1	Max. Number of Orders.

Resonator Card (First Order)

<u>Cols.</u>	<u>Value</u>	
5	1	Number of Resonators in $Z_A$
10	1	Number of Resonators in $Z_B$

$Z_A$  Parameter Card

<u>Col.</u>	<u>Value</u>	
1 - 14	+.10000000 E-05	$C_1$ Farads.

$Z_B$  Parameter Card

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.00000000 E-99	
15 - 28	+.00000000 E-99	
29 - 42	+.10000000 E+04	$R_2$ in Ohms.

Order Card (Second Section)

<u>Col</u>	<u>Value</u>	
5	3	Max. Number of Orders in a Section Impedance.

Resonator Card (First Order)

<u>Col</u>	<u>Value</u>	
5	1	Number of $Z_A$ Resonators
10	3	Number of $Z_B$ Resonators

$Z_A$  Parameter Card

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.10000000 E-11	$C_3$ in Farads
15 - 28	+.00000000 E-99	
29 - 42	+.10000000 E+07	$R_4$ in Ohms.

Z<sub>B</sub> Parameter Cards

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.30000000 E-06	C <sub>7</sub> in Farads.
15 - 28	+.20000000 E-12	L <sub>6</sub> in Henries.
29 - 42	+.10000000 E+04	R <sub>5</sub> in Ohms.

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.10000000 E-11	C <sub>8</sub> in Farads.

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.00000000 E-99	
15 - 28	+.10000000 E-05	L <sub>10</sub> in Henries.
29 - 42	+.10000000 E+03	R <sub>9</sub> in Ohms.

Resonator Card (Second Order)

<u>Col.</u>	<u>Value</u>	
5	0	No Z <sub>A</sub> Resonators
10	2	Two Resonators in Z <sub>B</sub> Impedance.

Z<sub>B</sub> Parameter Cards

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.10000000E-12	C <sub>12</sub> in Farads.
15 - 28	+.00000000E-99	
29 - 42	+.50000000E+02	R <sub>11</sub> in Ohms.

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.00000000E-99	
15 - 28	+.20000000E-12	L <sub>13</sub> in Henries.
29 - 42	+.10000000E+03	R <sub>14</sub> in Ohms.

Resonator Card (Third Order)

<u>Cols.</u>	<u>Value</u>	
5	0	No Z <sub>A</sub> Resonator
10	1	1 Z <sub>B</sub> Resonator

Z<sub>B</sub> Parameter Card

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.50000000E-08	C <sub>17</sub> in Farads.
15 - 28	+.10000000E-07	L <sub>15</sub> in Henries.
29 - 42	+.10000000E+03	R <sub>16</sub> in Ohms.

Order Card (Input Termination)

<u>Col.</u>	<u>Value</u>	
5	1	One Order

Resonator Card (First Order)

<u>Col.</u>	<u>Value</u>	
5	3	Three Resonators in Z <sub>A</sub> .
10	0	No Z <sub>B</sub> Resonators.

Z<sub>A</sub> Parameter Cards

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.00000000E-99	
15 - 28	+.50000000E-06	L <sub>18</sub> in Henries.

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.00000000E-99	
15 - 28	+.00000000E-99	
29 - 42	+.10000000E+03	R <sub>19</sub> in Ohms.

<u>Cols.</u>	<u>Value</u>	
1 - 14	+.10000000E-05	C <sub>21</sub> in Farads.
15 - 28	+.00000000E-99	
29 - 42	+.10000000E+06	R <sub>20</sub> in Ohms.

This completes the input to sample problem 3. A listing of these cards is found on page 27. A listing of the results is found on Page 28.



SAMPLE PROBLEM 1 FOR GENERAL FILTER NETWORK ANALYSIS PROGRAM 29APR64

1. 2 2 1 1. 10. 50. 1. 0.0

1 1 1

+ .10000000E-05 + .00000000E-99 + .10000000E+04  
+ .00000000E-99 + .20000000E-05

1 2 1

+ .10000000E-07  
+ .00000000E-99 + .00000000E-99 + .10000000E+06  
+ .00000000E-99 + .10000000E-05

1 1 0

+ .00000000E-99 + .00000000E-99 + .50000000E+02

GENERAL NETWORK ANALYSIS PROGRAM

SAMPLE PROBLEM 1 FOR GENERAL FILTER NETWORK ANALYSIS PROGRAM 29APR64

ZDR= 50.000 VDR= 1.000 VOI= 0.000

SECTIONS= 2 LETTR= 2 NCYCL= 1

ORDER 1IS 1

K= 1 I= 1 MP= .1 MQ= 1  
 C( 1 , 1 , 1 ) = .10000000E-05 L = .00000000E-99 R = .10000000E+04  
 C( 2 , 1 , 1 ) = .00000000E-99 L = .20000000E-05 R = .00000000E-99

ORDER 2IS 1

K= 2 I= 1 MP= 2 MQ= 1  
 C( 1 , 2 , 1 ) = .10000000E-07 L = .00000000E-99 R = .00000000E-99  
 C( 3 , 2 , 1 ) = .00000000E-99 L = .00000000E-99 R = .10000000E+06  
 C( 2 , 2 , 1 ) = .00000000E-99 L = .10000000E-05 R = .00000000E-99

ORDER 3IS 1

K= 3 I= 1 MP= 1 MQ= 0  
 C( 1 , 3 , 1 ) = .00000000E-99 L = .00000000E-99 R = .50000000E+02

F	DB	BETA	ZREAL	ZIMAG	T(F)
1.0000	4.9387	76.462	5.2953	-9.2583	
1.0010	4.9450	76.513	5.2851	-9.2400	.14120277E-06
2.0000	10.7053	114.502	1.1774	4.1843	
2.0010	10.7105	114.531	1.1760	4.1947	.81361108E-07
3.0000	15.4775	138.972	.4043	13.2695	
3.0010	15.4819	138.993	.4039	13.2777	.57249998E-07
4.0000	19.5824	156.885	.1711	20.9877	
4.0010	19.5862	156.901	.1709	20.9950	.43333332E-07
5.0000	23.2046	170.717	.0829	28.1290	
5.0010	23.2081	170.729	.0828	28.1360	.34055554E-07
6.0000	26.4500	181.741	.0443	34.9788	
6.0010	26.4531	181.751	.0443	34.9855	.27527777E-07
7.0000	29.3887	190.725	.0256	41.6625	
7.0010	29.3915	190.733	.0256	41.6691	.22638888E-07
8.0000	32.0716	198.173	.0157	48.2434	
8.0010	32.0741	198.180	.0157	48.2500	.18916666E-07
9.0000	34.5371	204.434	.0101	54.7565	
9.0010	34.5395	204.440	.0101	54.7630	.15999999E-07
10.0000	36.8156	209.760	.0068	61.2225	
10.0010	36.8178	209.765	.0068	61.2289	.13666666E-07

SAMPLE PROBLEM 2 NETWORK ANALYSIS PROGRAM

1MAY64

50.            2.            100.            100.            1.            0.0

  2    1    1  
+.10000000E+05+.50000000E+01

  .1  
+.10000000E+05+.50000000E+01

  1    2  
+.30000000E-10

+.20000000E-07

+.00000000E-99+.10000000E-06

  1  
+.10000000E+05+.50000000E+01

  1    2  
+.30000000E-09

+.40000000E-06

+.00000000E-99+.15000000E-06

  1,  
  1    0  
+.00000000E-99+.00000000E-99+.10000000E+03

GENERAL NETWORK ANALYSIS PROGRAM

SAMPLE PROBLEM 2 NETWORK ANALYSIS PROGRAM 1MAY64

ZOR= 100.000 VOR= 1.000 VDI= 0.000  
 RAO= .10000000E+05 RBO= .50000000E+01

SECTIONS= 2 LETTR= 1 NCYCL= 1  
 R(1, 1) = .10000000E+05 R(2, 1) = .50000000E+01 1 ORDERS  
 K= 1 I= 1 MP= 1 MQ= 2  
 C( 1, 1, 1 ) = .30000000E-10 L = .00000000E-99 R = .00000000E-99  
 C( 2, 1, 1 ) = .20000000E-07 L = .00000000E-99 R = .00000000E-99  
 C( 4, 1, 1 ) = .00000000E-99 L = .10000000E-06 R = .00000000E-99  
 R(1, 2) = .10000000E+05 R(2, 2) = .50000000E+01 1 ORDERS  
 K= 2 I= 1 MP= 1 MQ= 2  
 C( 1, 2, 1 ) = .30000000E-09 L = .00000000E-99 R = .00000000E-99  
 C( 2, 2, 1 ) = .40000000E-06 L = .00000000E-99 R = .00000000E-99  
 C( 4, 2, 1 ) = .00000000E-99 L = .15000000E-06 R = .00000000E-99

ORDER 3IS 1

K= 3 I= 1 MP= 1 MQ= 0  
 C( 1, 3, 1 ) = .00000000E-99 L = .00000000E-99 R = .10000000E+03

F	DB	BETA	ZREAL	ZIMAG	T(F)
50.0000	21.6168	89.216	5.2596	-4.8518	
50.0010	21.6170	89.216	5.2596	-4.8517	.55555554E-09
52.0000	21.9569	89.613	5.2426	-4.6662	
52.0010	21.9571	89.613	5.2425	-4.6661	.55833331E-09
54.0000	22.2846	89.994	5.2274	-4.4943	
54.0010	22.2848	89.994	5.2274	-4.4943	.52499998E-09
56.0000	22.6008	90.361	5.2138	-4.3347	
56.0010	22.6009	90.361	5.2138	-4.3346	.52777776E-09
58.0000	22.9061	90.716	5.2016	-4.1859	
58.0010	22.9063	90.716	5.2016	-4.1859	.47222220E-09
60.0000	23.2015	91.060	5.1906	-4.0471	
60.0010	23.2016	91.060	5.1906	-4.0470	.47222220E-09
62.0000	23.4874	91.393	5.1806	-3.9172	
62.0010	23.4875	91.394	5.1806	-3.9171	.47222220E-09
64.0000	23.7646	91.718	5.1715	-3.7953	
64.0010	23.7647	91.718	5.1715	-3.7953	.41666665E-09
66.0000	24.0334	92.033	5.1633	-3.6809	
66.0010	24.0336	92.034	5.1633	-3.6808	.41666665E-09
68.0000	24.2946	92.341	5.1558	-3.5731	
68.0010	24.2947	92.342	5.1557	-3.5731	.44444443E-09
70.0000	24.5484	92.642	5.1488	-3.4715	
70.0010	24.5485	92.643	5.1488	-3.4714	.41666665E-09
72.0000	24.7953	92.937	5.1425	-3.3755	
72.0010	24.7954	92.937	5.1425	-3.3755	.41666665E-09
74.0000	25.0356	93.225	5.1367	-3.2847	
74.0010	25.0358	93.226	5.1367	-3.2847	.38888887E-09
76.0000	25.2698	93.508	5.1313	-3.1987	
76.0010	25.2699	93.509	5.1313	-3.1986	.38888887E-09
78.0000	25.4982	93.786	5.1263	-3.1170	

78.0010	25.4983	93.786	5.1263	-3.1170	.38888887E-09
80.0000	25.7209	94.059	5.1217	-3.0395	
80.0010	25.7210	94.059	5.1217	-3.0394	.36111110E-09
82.0000	25.9384	94.328	5.1174	-2.9657	
82.0010	25.9385	94.328	5.1174	-2.9657	.36111110E-09
84.0000	26.1509	94.593	5.1134	-2.8954	
84.0010	26.1510	94.593	5.1134	-2.8954	.36111110E-09
86.0000	26.3585	94.853	5.1097	-2.8284	
86.0010	26.3586	94.853	5.1097	-2.8284	.38888887E-09
88.0000	26.5616	95.110	5.1062	-2.7645	
88.0010	26.5617	95.110	5.1062	-2.7644	.36111110E-09
90.0000	26.7604	95.364	5.1030	-2.7033	
90.0010	26.7605	95.364	5.1030	-2.7033	.36111110E-09
92.0000	26.9549	95.614	5.1000	-2.6449	
92.0010	26.9550	95.615	5.1000	-2.6449	.36111110E-09
94.0000	27.1455	95.862	5.0971	-2.5889	
94.0010	27.1456	95.862	5.0971	-2.5889	.33333332E-09
96.0000	27.3323	96.107	5.0944	-2.5353	
96.0010	27.3324	96.107	5.0944	-2.5352	.36111110E-09
98.0000	27.5154	96.349	5.0919	-2.4838	
98.0010	27.5155	96.349	5.0919	-2.4838	.33333332E-09
100.0000	27.6950	96.588	5.0896	-2.4344	
100.0010	27.6951	96.588	5.0896	-2.4344	.33333332E-09

SAMPLE PROBLEM 3 NETWORK ANALYSIS PROGRAM

4MAY64

5. 2 2 5. 100. 100. 1. 0.0

1  
1 1  
+.10000000E-05  
+.00000000E-99+.00000000E-99+.10000000E+04  
3  
1 3  
+.10000000E-11+.00000000E-99+.10000000E+07  
+.30000000E-06+.20000000E-12+.10000000E+04  
+.10000000E-11  
+.00000000E-99+.10000000E-05+.10000000E+03  
0 2  
+.10000000E-12+.00000000E-99+.50000000E+02  
+.00000000E-99+.20000000E-12+.10000000E+03  
0 1  
+.50000000E-08+.10000000E-07+.10000000E+03  
1  
3 0  
+.00000000E-99+.50000000E-06  
+.00000000E-99+.00000000E-99+.10000000E+03  
+.10000000E-05+.00000000E-99+.10000000E+06

GENERAL NETWORK ANALYSIS PROGRAM

SAMPLE PROBLEM 3 NETWORK ANALYSIS PROGRAM 4MAY64

ZOR= 100.000 VOR= 1.000 VOI= 0.000

SECTIONS= 2 LETTR= 2 NCYCL= 2

ORDER 1IS 1

K= 1 I= 1 MP= 1 MQ= 1  
 C( 1 , 1 , 1 ) = .10000000E-05 L = .00000000E-99 R = .00000000E-99  
 C( 2 , 1 , 1 ) = .00000000E-99 L = .00000000E-99 R = .10000000E+04

ORDER 2IS 3

K= 2 I= 1 MP= 1 MQ= 3  
 C( 1 , 2 , 1 ) = .10000000E-11 L = .00000000E-99 R = .10000000E+07  
 C( 2 , 2 , 1 ) = .30000000E-06 L = .20000000E-12 R = .10000000E+04  
 C( 4 , 2 , 1 ) = .10000000E-11 L = .00000000E-99 R = .00000000E-99  
 C( 6 , 2 , 1 ) = .00000000E-99 L = .10000000E-05 R = .10000000E+03  
 K= 2 I= 2 MP= 0 MQ= 2  
 C( 2 , 2 , 2 ) = .10000000E-12 L = .00000000E-99 R = .50000000E+02  
 C( 4 , 2 , 2 ) = .00000000E-99 L = .20000000E-12 R = .10000000E+03  
 K= 2 I= 3 MP= 0 MQ= 1  
 C( 2 , 2 , 3 ) = .50000000E-08 L = .10000000E-07 R = .10000000E+03

ORDER 3IS 1

K= 3 I= 1 MP= 3 MQ= 0  
 C( 1 , 3 , 1 ) = .00000000E-99 L = .50000000E-06 R = .00000000E-99  
 C( 3 , 3 , 1 ) = .00000000E-99 L = .00000000E-99 R = .10000000E+03  
 C( 5 , 3 , 1 ) = .10000000E-05 L = .00000000E-99 R = .10000000E+06

F	DB	BETA	ZREAL	ZIMAG	T(F)
5.0000	180.7194	180.838	1065.2886	.2987	
5.0010	180.7228	180.838	1065.2887	.2999	.77777775E-09
10.0000	192.0992	182.003	1066.2497	4.5705	
10.0010	192.1008	182.004	1066.2501	4.5712	.55555554E-09
15.0000	198.3916	182.824	1067.7597	7.4989	
15.0010	198.3926	182.825	1067.7601	7.4994	.38888887E-09
20.0000	202.7228	183.368	1069.6027	9.6923	
20.0010	202.7236	183.368	1069.6031	9.6927	.22222221E-09
25.0000	206.0672	183.685	1071.5953	11.2810	
25.0010	206.0678	183.685	1071.5957	11.2812	.13888888E-09
30.0000	208.8277	183.835	1073.5931	12.3601	
30.0010	208.8282	183.835	1073.5935	12.3602	.55555554E-10
35.0000	211.1986	183.869	1075.4981	13.0270	
35.0010	211.1991	183.869	1075.4985	13.0271	.00000000E-99
40.0000	213.2865	183.828	1077.2542	13.3759	
40.0010	213.2869	183.828	1077.2546	13.3760	-.27777777E-10
45.0000	215.1560	183.741	1078.8370	13.4893	
45.0010	215.1563	183.741	1078.8373	13.4893	-.55555554E-10
50.0000	216.8501	183.628	1080.2430	13.4347	

50.0010	216.8505	183.628	1080.2433	13.4347	-.55555554E-10
55.0000	218.3995	183.501	1081.4807	13.2648	
55.0010	218.3998	183.501	1081.4809	13.2647	-.55555554E-10
60.0000	219.8269	183.369	1082.5648	13.0190	
60.0010	219.8272	183.369	1082.5651	13.0189	-.55555554E-10
65.0000	221.1500	183.237	1083.5125	12.7260	
65.0010	221.1502	183.237	1083.5127	12.7260	-.83333331E-10
70.0000	222.3827	183.108	1084.3404	12.4067	
70.0010	222.3829	183.108	1084.3405	12.4066	-.55555554E-10
75.0000	223.5363	182.985	1085.0647	12.0753	
75.0010	223.5365	182.985	1085.0649	12.0752	-.55555554E-10
80.0000	224.6202	182.867	1085.6997	11.7419	
80.0010	224.6204	182.867	1085.6998	11.7418	-.55555554E-10
85.0000	225.6422	182.757	1086.2579	11.4133	
85.0010	225.6424	182.757	1086.2580	11.4132	-.55555554E-10
90.0000	226.6089	182.653	1086.7501	11.0938	
90.0010	226.6091	182.653	1086.7502	11.0937	-.55555554E-10
95.0000	227.5257	182.555	1087.1857	10.7862	
95.0010	227.5259	182.555	1087.1857	10.7862	-.55555554E-10
100.0000	228.3976	182.464	1087.5724	10.4922	
100.0010	228.3977	182.464	1087.5725	10.4922	-.55555554E-10



APPENDIX A

The following paragraphs will outline the equations programmed for the various subprograms. These assume steady state conditions on the imaginary axis ( $S = j\omega$ ) with linear, lumped, bilateral, passive elements. The program processes basic sections starting with that section nearest the output termination. Regardless of section type, this section will have  $Z_A$  and  $Z_B$  impedances.

The DRPTZ subprogram calculates the impedances and admittances for a single order as in Figure 8 for both the  $Z_A$  and  $Z_B$  impedance.

The number of parallel branches, or resonators as they are called in the Westinghouse report, allowed is a function of available core storage. As written, three resonators per order are allowed although this could be increased by reallocating storage, i.e., reducing the number of orders per basic section impedance or the number of basic sections. The impedance for a single branch or resonator is:

$$Z_1 = R_1 + j\omega L_1 + \frac{1}{j\omega C_1} = R_1 + j\left(\omega L_1 - \frac{1}{\omega C_1}\right)$$

$$Y_1 = 1/Z_1 = \frac{1}{R_1 + j\left(\omega L_1 - \frac{1}{\omega C_1}\right)}$$
$$= \frac{R_1 + j\left(\frac{1}{\omega C_1} - \omega L_1\right)}{R_1^2 + \left(\omega L_1 - \frac{1}{\omega C_1}\right)^2}$$

$$Y_1 = Y_{RL1} + jY_{IM1}$$

$$Y_{RL1} = \frac{R_1}{R_1^2 + \left(\omega L_1 - \frac{1}{\omega C_1}\right)^2}$$

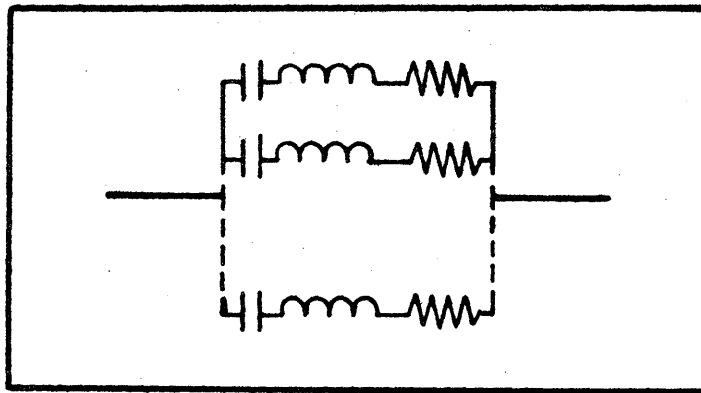


FIGURE 8 - A SINGLE ORDER

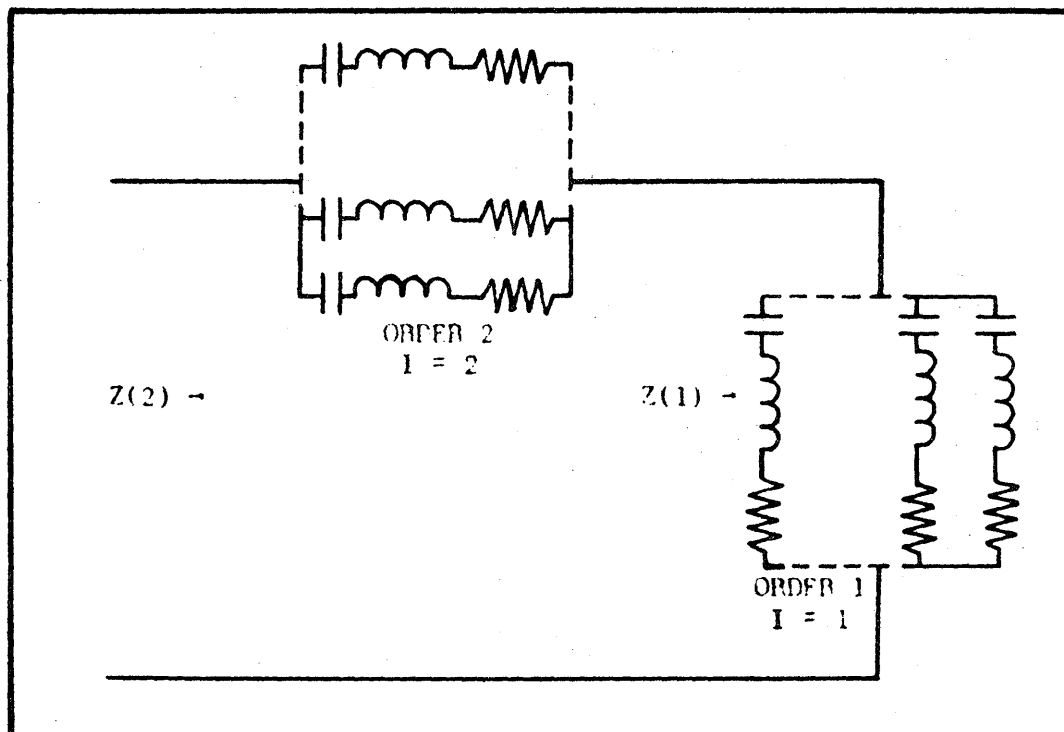


FIGURE 9 - TWO ORDERS OF A BASIC SECTION IMPEDANCE

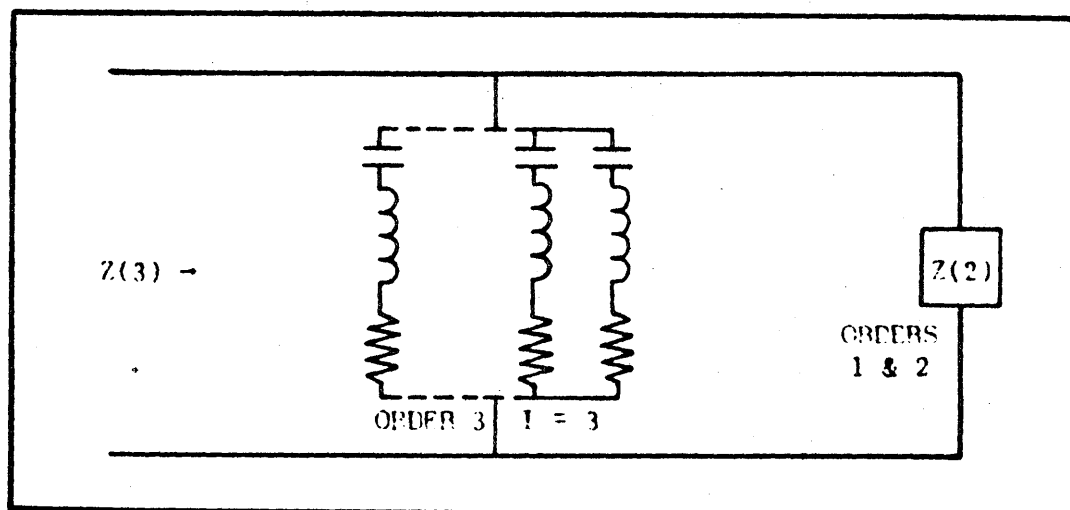


FIGURE 10 - THREE ORDERS OF A BASIC SECTION IMPEDANCE

244

$$YIM_1 = \frac{\frac{1}{\omega C_1} - \omega L_1}{R_1^2 + (\omega L_1 - \frac{1}{\omega C_1})^2}$$

Summing admittances for all k parallel branches in an order gives:

$$YRL_T = \sum_{i=1}^k YRL_i$$

$$YIM_T = \sum_{i=1}^k YIM_i$$

$$Y_T = YRL_T + jYIM_T$$

$$Z_T = 1/Y_T = ZRL_T + jZIM_T.$$

The DRPTZ subprogram stores values of  $ZRL_T$ , and  $ZIM_T$  impedances. Program control is then given to the ORDER subprogram.

The ORDER subprogram first stores the impedance values just computed by DRPTZ as subscripted impedances  $Z(I)$ , with the order counter I set at 1. Control is returned to the main line which determines if either the  $Z_A$  or  $Z_B$  basic section impedance consists of more than one order. If not, control is given to the proper recursion subprogram. If there are more than one order in either  $Z_A$  or  $Z_B$ , the counter is incremented to 2 and control is returned to DRPTZ.

The DRPTZ subprogram then zeros the variables  $YRL_T$  and  $YIM_T$  and repeats the process described above using the element values for the second orders of the appropriate basic section impedances. Control is then given to the ORDER subprogram.

The ORDER subprogram determines if the order counter I equals 2 or 3. If I is 2, ORDER simply sums the impedances  $Z_T$  just computed by DRPTZ with those previously stored as subscripted impedances when I equalled 1. Figure 9 illustrates the situation.

$$Z(1) = ZTRL(1) + jZTIM(1)$$

$$Z(2) = Z(1) + Z_T = ZTRL(1) + jZTIM(1) + ZRL_T + jZIM_T = ZTRL(2) + jZTIM(2)$$

$$ZTRL(2) = ZTRL(1) + ZRL_T$$

$$ZTIM(2) = ZTIM(1) + ZIM_T$$

If the order counter I is 3, the ORDER subprogram has the situation in Figure 10. DRPTZ again has computed the impedance  $Z_T$  and admittance  $Y_T$  for the third order.

$$Y_T = YRL_T + jYIM_T \text{ (for Order 3)}$$

$$Y(3) = 1/Z(3) = \frac{1}{Z(2)} + Y_T = \frac{1}{ZTRL(2) + jZTIM(2)} + Y_T$$

$$= \frac{ZTRL(2) - jZTIM(2)}{A} + YRL_T + jYIM_T$$

where  $A = [ZTRL(2)]^2 + [ZTIM(2)]^2$

$$Y(3) = \frac{1}{A} \left\{ [ZTRL(2) + A(YRL_T)] + j [A(YIM_T) - ZTIM(2)] \right\}$$

Let  $B = ZTRL(2) + A(YRL_T)$

$C = A(YIM_T) - ZTIM(2)$

$$Z(3) = 1/Y(3) = \frac{A}{B + jC} = \frac{A(B - jC)}{B^2 + C^2}$$

Control is returned to the main line. If there are no further orders to be processed in the particular basic section impedances, control is given to the appropriate recursion subprogram LADDER or LATTICE.

The LADDER subprogram has the situation in Figure 11. As mentioned previously computation proceeds from right to left. The first step is to pick up the complex values of impedance just computed by the DRPTZ and ORDER subprograms, and the voltage and impedance existing at the output of the particular section being considered. If the first section to the left of the output termination is being processed, these latter two values are  $V_0$  and  $R_0$ . In the general case they are  $V_{K-1}$  and  $Z_{K-1}$ .

$$V_{K-1} = VRL_{K-1} + jVIM_{K-1}$$

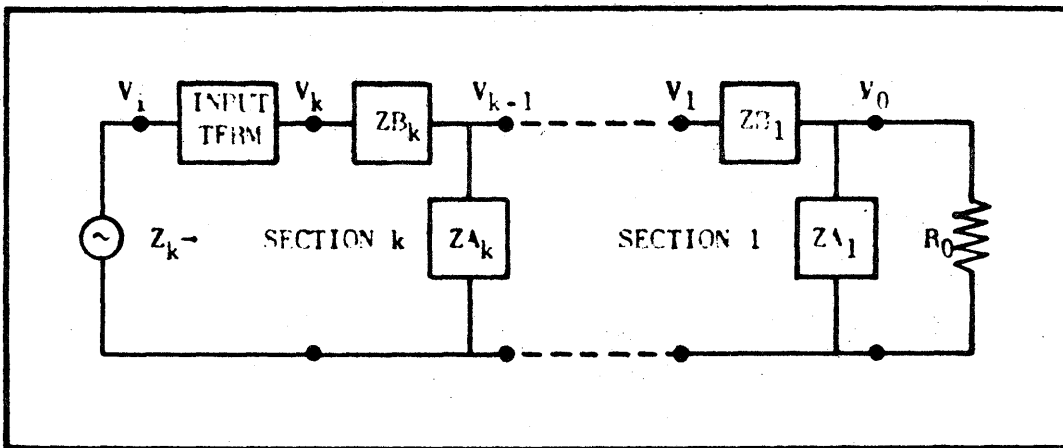


FIGURE 11 - CASCADED LADDER SECTIONS

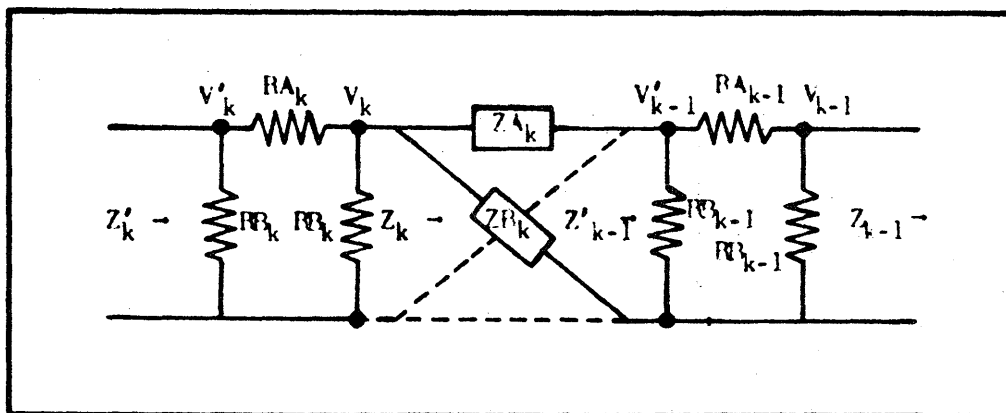


FIGURE 12 - SYMMETRICAL LATTICE SECTION

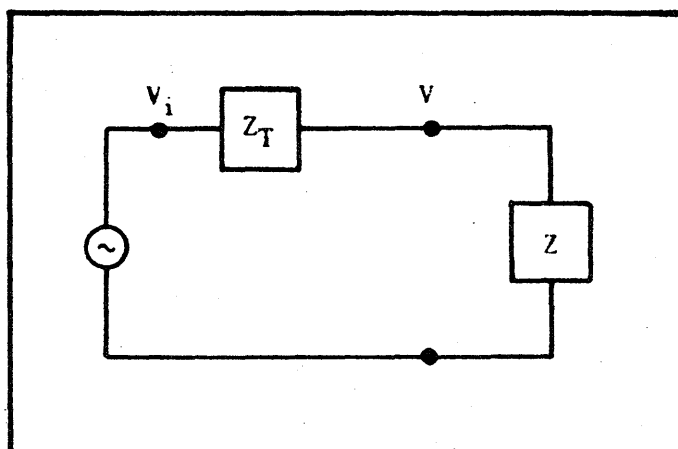


FIGURE 13 - FILTER NETWORK  $Z$  AND INPUT TERMINATION  $Z_T$

$$Z_{K-1} = ZR_{K-1} + jZIM_{K-1}$$

Note that  $ZA_K$  and  $ZB_K$  are equal to the appropriate  $Z(I)$  values determined in DRPTZ and ORDER where  $I$  is the maximum order of the  $k^{\text{th}}$  section.

$$Z_K = ZB_K + \frac{ZA_K Z_{K-1}}{ZA_K + Z_{K-1}} = ZR_{K-1} + jZIM_{K-1}$$

$$\frac{V_K - V_{K-1}}{ZB_K} = V_{K-1} \left( \frac{1}{ZA_K} + \frac{1}{Z_{K-1}} \right)$$

$$V_K = V_{K-1} \left\{ 1 + ZB_K \left( \frac{1}{ZA_K} + \frac{1}{Z_{K-1}} \right) \right\}$$

The LATTICE subprogram has the situation presented in Figure 12.

It will be noted that symmetrical resistive pi sections are placed at each end of the basic symmetrical lattice section to permit a definite amount of attenuation or isolation between sections. This feature can effectively be eliminated by inserting zero values for the series elements  $R_B$  and very large values for the shunt  $R_A$  elements. Again  $ZA_K$  and  $ZB_K$  have been computed by the DRPTZ and ORDER subprograms and are equal to the appropriate  $Z(I)$  values where  $I$  is the maximum order of the  $K^{\text{th}}$  section. The recursion equations for the LATTICE subprogram follow. Computation proceeds from right to left as before. Note that the following equations hold for Fig. 12, in which the pi section shunt  $RA$  and series  $RB$  have been interchanged.

$$Z'_{K-1} = \frac{RB_{K-1} \left\{ RA_{K-1} + \frac{(RB_{K-1})(Z_{K-1})}{RB_{K-1} + Z_{K-1}} \right\}}{RB_{K-1} + RA_{K-1} + \frac{(RB_{K-1})(Z_{K-1})}{(RB_{K-1} + Z_{K-1})}}$$

$$V'_{K-1} = V_{K-1} \left\{ 1 + RA_{K-1} \left( \frac{1}{RB_{K-1}} + \frac{1}{Z_{K-1}} \right) \right\}$$

$$Z_K = \frac{Z_{K-1}' (Z_{A_K} + Z_{B_K}) + 2(Z_{A_K}) (Z_{B_K})}{2 Z_{K-1}' + Z_{A_K} + Z_{B_K}}$$

$$V_K = \frac{V_{K-1}' \{ Z_{K-1}' (Z_{A_K} + Z_{B_K}) + 2(Z_{A_K}) (Z_{B_K}) \}}{Z_{K-1}' (Z_{B_K} - Z_{A_K})}$$

$$Z_K' = \frac{(1/Z_K + 1/R_{B_K}) R_{A_K} + 1}{(1/R_{B_K}) \{ (1/Z_K + 1/R_{B_K}) R_{A_K} + 1 \} + 1/Z_K + 1/R_{B_K}}$$

$$V_K' = V_K \{ (1/Z_K + 1/R_{B_K}) R_{A_K} + 1 \}$$

The procedures described are repeated section-by-section until the last filter section before the input termination has been processed. The impedance of the entire filter plus the output termination is stored for subsequent output as ZREAL and ZIMAG. The Termin subprogram is called to monitor calculation of the input termination impedance and input voltage  $V_i$  required to produce the specified output voltage  $V_o$ . The subprogram uses the DRPTZ subprogram to determine the single-ordered input termination impedance, subsequently computing  $V_i$ . Figure 13 illustrates the situation with Z representing the input impedance of the filter plus the output resistance and  $Z_T$  the input termination. V is the voltage at the filter input.

$$\frac{V_i - V}{Z_T} = \frac{V}{Z}$$

$$\frac{V_i}{V} = \frac{Z_T}{Z} + 1$$

$$V_i = V \left( \frac{Z_T}{Z} + 1 \right) = \frac{V}{Z} (Z_T + Z)$$

$Z_T$  is stored as follows:

$$Z_T = RK + jXK$$

and control is returned to the main line which calls the BETADB subprogram. BETADB calculates the phase angle  $\beta$  of the voltage  $V_1$  and the insertion loss A. These equations follow:

$$V_1 = VPRL + jVPII$$

$$\beta = \tan^{-1} (VPII/VPRL)$$

$$A = 20 \log_{10} \left| \frac{R_o}{(RK + R_o + jXK)} \cdot \frac{V_1}{V_o} \right|$$

$\beta$  is reported as a positive angle.

Following calculation of the voltage angle and insertion loss by the BETADB subprogram, control is returned to the main line which outputs the following information:

Frequency (in megacycles)  
 Insertion loss (in decibels)  
 Input voltage angle (in degrees)  
 Filter Impedance including output termination  
 (in ohms - rectangular form).

This completes one full pass thru the program.

Time delay is approximated by repeating the complete procedure with the frequency incremented by 1000 cycles, finding the change in input voltage angle, and dividing by a constant proportional to the increment in frequency.

$$T = \frac{\Delta\beta}{1000 C}$$

As written,  $C = 360$ . This result approximates the slope of the phase versus frequency curve at the frequency point 500 cycles above the base frequency. The mainline then outputs the same information as above at the incremented frequency, as well as T. The frequency value is decremented by 1000 cycles, and tested against the maximum frequency value, FMAX. If less than that value, frequency is incremented by FDEL and the entire procedure repeated. Eventually frequency reaches FMAX and the program returns to its beginning point ready to read in another title card and complete set of frequency and element cards representing a new filter. If there are no cards in the input hopper the Reader-No-Feed light is turned on.



C GENERAL FILTER NETWORK ANALYSIS PROGRAM FOR LADDER AND LATTICE  
 C SECTIONS.  
 C Z = COMPLEX DRIVING - POINT IMPEDANCE FUNCTION  
 C Y = COMPLEX DRIVING - POINT ADMITTANCE FUNCTION  
 C SUBSCRIPT (LKI) = (L) STANDARDIZED POSITION OF (K)  
 C SECTION OF (I) ORDER  
 C PROGRAM BASED UPON WESTINGHOUSE REPORT TITLED ELECTRONIC DIGITAL  
 C COMPUTER ANALYSIS OF CASCADED NETWORKS BY RICHARD TUZNIK  
 C AND DEAN H. WOOD.  
 C PROGRAM CAN BE EXPANDED TO HANDLE PI, TEE, AND BRIDGED-TEE  
 C SECTIONS ADDING APPROPRIATE SUBPROGRAMS AND MAKING  
 C NECESSARY REVISIONS IN EXISTING ROUTINES.

DIMENSION C(6,5,3), R(6,5,3), AL(6,5,3), MP(5,3), MQ(5,3),  
 IRA(5), RB(5), IORDR(5), ZAKRL(3), ZBKRL(3), ZAKIM(3), ZBKIM(3)  
 COMMON C, R, AL, MP, MQ, RA, RB, ZAKRL, ZBKRL, ZAKIM, ZBKIM, YARL,  
 IYBRL, YAIM, YBIM, ZARL, ZBRL, ZAIM, ZBIM, F, F2, IREC, ZPRL, ZPIM,  
 ZVPRL, VPIM, ZRL, ZIM, VRL, VIM, K, I, IPUNC, RAO, RBO, DEN, L,  
 3INDX, REAL, AIMG, RLDEN, ANMIN, RLNUM, DENIM, A, B, QUAN, IORDR,  
 4LETT, VAR1, IORD, ZARLK, ZBRLK, ZAIMK, ZBIMK, NSECT, NT  
 COMMON XK, RK, ZOR, ZOI, VOR, VOI, BETA, DB, XIMAG  
 PI = 3.14159265  
 PI2 = 6.2831853  
 CONST = 1.0/360000.  
 CONV = 180.0/PI

100 FORMAT (/32HGENERAL NETWORK ANALYSIS PROGRAM/)  
 101 FORMAT (6F11.5)  
 102 FORMAT (3I5)

103 FORMAT (3E14.8)

C  
C  
C

12 READ 116

C READ TITLE CARD WITH UP TO 80 CHARACTERS.

PUNCH 100

PUNCH 116

116 FORMAT(40H

,40H

1

C  
C  
C

READ 101, FMIN, FDEL, FMAX, ZOR, VOR, VOI

C ZOR IS OUTPUT TERMINAL RESISTANCE IN OHMS, FMIN, FDEL, AND  
C FMAX ARE FREQUENCIES IN MC.

C VOR AND VOI ARE OUTPUT TERMINAL VOLTAGES. FOR REFERENCE,

C VOR = 1.0, VOI = 0.0

C  
C  
C

READ 102, NSECT, LETTR, NCYCL

C NSECT IS NUMBER OF INVERTED L-SECTIONS OR SYMMETRICAL LATTICE  
C SECTIONS.

C LETTR SPECIFIES TYPE OF BASIC SECTION, = 1 FOR SYM LATTICE,  
C = 2 FOR LADDER.

C NCYCL = NO OF TIMES GROUP OF SECTIONS IS USED BEFORE PROCEEDING  
C TO FINAL PORTION OF PROGRAM.

C  
C  
C

GO TO (45,54), LETTR

45 READ 103, RAO, RBO

C READ INITIAL PI PAD VALUES, RAO IS SHUNT BRANCH, RBO IS SERIES  
C BRANCH. THESE VALUES ARE REQUIRED FOR LATTICE SECTIONS ONLY.

C  
C  
C  
C  
C

PUNCH OUT INPUT VALUES

55 PUNCH 105, ZOR, VOR, VOI, RAO, RBO

105 FORMAT (/4HZOR=,F9.3, 3X4HVOR=,F9.3, 3X4HVOI=,F9.3/ 4HRAO=,E14.8,  
13X4HRBO=,E14.8/)

GO TO 56

54 PUNCH 104, ZOR, VOR, VOI

104 FORMAT (/4HZOR=,F9.3, 3X4HVOR=,F9.3, 3X4HVOI=,F9.3/)

56 PUNCH 111, NSECT, LETTR, NNCYCL

111 FORMAT (9HSECTIONS=,I4, 3X6HLETTR=,I4, 3X6HNCYCL=,I4)

NSECT = NSECT+1

DO 2 K = 1, NSECT

C  
C

READ 102, IORDR(K)

C IORDR IS NUMBER OF ORDERS PER SECTION IMPEDANCE ARM OF A SECTION.  
C IF SECTION IMPEDANCE ARMS HAVE DIFFERENT ORDERS, USE LARGER  
C VALUE.

C

C

C DO NOT READ PI PAD VALUES FOLLOWING TERMINATION SECTION.

GO TO (47,48),LETTR

47 IF(NSECT-K) 48, 48, 64

C

C

64 READ 103, RA(K), RB(K)

C RA AND RB ARE PI PAD VALUES FOLLOWING SECTION K. RA - SHUNT

C BRANCH. RB - SERIES BRANCH.

C

C

PUNCH 106, K, RA(K), K, RB(K), IORDR(K)

GO TO 49

48 PUNCH 110, K, IORDR(K)

110 FORMAT (/5HORDER, I4, 2HIS, I4/)

106 FORMAT(4HR(1,,I2,5H ) = ,E14.8,2X4HR(2,,I2,5H ) = ,E14.8,2X,I2,  
17H ORDERS)

49 IORD = IORDR(K)

DO 2 I = 1, IORD

C

C

READ 102, MP(K,I), MQ(K,I)

C MQ = NO OF RESONATORS IN (I)TH ORDER IMPEDANCE ZB IN (K)TH SECTION

C MP = NO OF RESONATORS IN (I)TH ORDER IMPEDANCE ZA IN (K)TH SECTION

C RESONATORS ARE NUMBER OF PARALLEL BRANCHES.

C

C

PUNCH 112, K, I, MP(K,I), MQ(K,I)

112 FORMAT (2HK=, I4, 3X2HI=, I4, 3X3HMP=, I4, 3X3HMQ=, I4)

IREC = MP(K,I)

L = 1

9 IF (IREC-L) 7, 8, 8

8 INDX = 2\*L-1

C

C

READ 103, C(INDX,K,I), AL(INDX,K,I), R(INDX,K,I)

C

READ C(FARADS),L(HENRIES),AND R(OHMS) VALUES OF ZA TYPE IMPEDANCE

C

WHICH IS PARALLEL ARM FOR INVERTED-L SECTIONS AND IS SERIES

C

ARM FOR LATTICE SECTIONS.

C

C

PUNCH 107, INDX, K, I, C(INDX,K,I), AL(INDX,K,I), R(INDX,K,I)

107 FORMAT (2HC(, I2, 2H ,, I2, 2H ,, I2, 5H ) = , E14.8, 2X4HL = ,

1E14.8, 2X4HR = , E14.8)

L = L+1

GO TO 9

7 IREC = MQ(K,I)

L = 1

10 IF (IREC-L) 2, 11, 11

11 INDX = 2\*L

C

C

READ 103, C(INDX,K,I), AL(INDX,K,I), R(INDX,K,I)

C

READ C(FARADS),L(HENRIES),AND R(OHMS) OF ZB TYPE IMPEDANCE WHICH

C

IS SERIES ARM FOR INVERTED-L SECTIONS AND IS PARALLEL ARM FOR

C

LATTICE SECTIONS.

255

PUNCH 107, INDX, K, I, C(INDX,K,I), AL(INDX,K,I), R(INDX,K,I)

L = L+1

GO TO 10

2 CONTINUE

FCOUN = FMIN

KKLL = 1

JFREQ = 1

NSECT = NSECT-1

20 ZRL = ZOR

ZOI=0.

ZIM = 0.0

VRL = VOR

VIM = VOI

IREC = 0

K = 1

I = 1

N = 1

GO TO (42,43),LETTR

C

C CALCULATE DROP FOR PI PAD PRECEDING FIRST SECTION.

C

42 CALL LATTIS

43 IREC = 2

F = (FCOUN\*1.0E+6)\*PI2

F2 = F\*F

N = 1

```
K = 1
25 IORD = IORDR(K)
    I = 1
22 CALL DRPTZ
    CALL ORDER
    IF(IORD-I) 3, 3, 21
21 I = I+1
    GO TO 22
    3 GO TO (36,37),LETTR
36 CALL LATTIS
    GO TO 38
37 CALL LADDER
38 IF(NSECT-K) 27, 27, 24
24 K = K+1
    GO TO 25
27 IF(NCYCL-N) 30, 30, 29
29 N = N+1
    K = 1
    GO TO 25
30 CONTINUE
    GO TO (99,44), LETTR
99 ZRLN = ZPRL
    ZIMN = ZPIM
    GO TO (97,98), KKLL
44 ZRLN = ZRL
    ZIMN = ZIM
    GO TO (97,98), KKLL
97 PUNCH 109
```

KKLL = 2

C  
C  
C

98 K = NSECT+1

C SET K AND I VALUES PERTAINING TO INPUT TERMINAL IMPEDANCE. K WILL  
C BE EQUAL TO NSECT+1, I = 1, WHERE THE IMPEDANCE IS CONSIDERED  
C ZA TYPE.

C

I = 1

CALL TERMIN

32 GO TO (33,34), JFREQ

C

33 JFREQ = 2

CALL BETADB

C CALCULATE DB AND BETA.

C

C

BETA1 = BETA

C ZRLN AND ZIMN ARE IMPEDANCE VALUES OF INSERTED NETWORK WITH OUTPUT  
C TERMINATION ZOR. THESE INCLUDE PI PADS FOR LATTICE SECTIONS.  
PUNCH 108, FCOUN, DB, BETA, ZRLN, ZIMN

C

C INCREMENT F BY 1000 CYCLES AND REPEAT ENTIRE PROCEDURE IN ORDER  
C TO OBTAIN TIME DELAY.

C

FCOUN = FCOUN+0.001



GO TO 20

34 CALL BETADB

T = (BETA-BETA1)\*CONST

JFREQ = 1

PUNCH 108, FCOUN, DB, BETA, ZRLN, ZIMN, T

C

C DECREMENT FREQUENCY BY 1000 CYCLES.

C

FCOUN = FCOUN-0.001

C

C IS FREQUENCY AT MAXIMUM VALUE.

C

IF (FMAX-FCOUN) 12, 12, 35

C

C NO. INCREMENT BY AMOUNT AT INPUT AND REPEAT CALCULATIONS.

C

35 FCOUN = FCOUN+FDEL

GO TO 20

C

C YES. GO TO START TO READ IN NEXT COMPLETE FILTER PROBLEM.

C

108 FORMAT (F9.4, 2XF9.4, 2XF10.3, 2XF10.4, 2XF10.4, 2XE14.8)

109 FORMAT (/4X1HF, 10X2HDB, 8X4HBETA, 8X5HZREAL, 7X5HZIMAG, 9X4HT(F)/)

END

SUBROUTINE BETADB

C CALCULATES BETA(RADIANS) AND DB FOR FREQUENCY

DIMENSION C(6,5,3), R(6,5,3), AL(6,5,3), MP(5,3), MQ(5,3),  
1RA(5), RB(5), IORDR(5), ZAKRL(3), ZBKRL(3), ZAKIM(3), ZBKIM(3)

COMMON C, R, AL, MP, MQ, RA, RB, ZAKRL, ZBKRL, ZAKIM, ZBKIM, YARL,  
1YBRL, YAIM, YBIM, ZARL, ZBRL, ZAIM, ZBIM, F, F2, IREC, ZPRL, ZPIM,

2VPRL, VPIM, ZRL, ZIM, VRL, VIM, K, I, IPUNC, RAO, RBO, DEN, L,

3INDX, REAL, AIMG, RLDEN, ANMIN, RLNUM, DENIM, A, B, QUAN, IORDR,

4LETT, VAR1, IORD, ZARLK, ZBRK, ZAIMK, ZBIMK, NSECT, NT

COMMON XK, RK, ZOR, ZOI, VOR, VOI, BETA, DB, XIMAG

PI = 3.14159265

PI2 = 6.2831853

CONV = 180.0/PI

BETA = (ATANF(VPIM/VPRL))\*CONV

C

C

MAKE BETA POSITIVE ANGLE BETWEEN 0 AND 360 DEGREES.

C

IF(VPRL) 1, 2, 2

1 BETA = BETA+180.0

GO TO 4

2 IF(BETA) 3, 4, 4

3 BETA = BETA+360.0

4 RNUM = ZOR\*VPRL-ZOI\*VPIM

XIMAG = VPRL\*ZOI+ZOR\*VPIM

XNUM = SQRTF(RNUM\*RNUM+XIMAG\*XIMAG)

RLDEN = RK\*VOR+ZOR\*VOR-XK\*VOI-ZOI\*VOI

XIMAD = XK\*VOR+ZOI\*VOR+RK\*VOI+ZOR\*VOI

DEN = SQRTF(RLDEN\*RLDEN+XIMAD\*XIMAD)

260

DB = 2.0\*4.3429448\*LOGF(XNUM /DEN)

RETURN

END

SUBROUTINE LATTIS

RECUSION FORMULAE FOR SYMMETRICAL LATTICE SECTION

CALCULATES ZRL, ZIM, ZPRL, ZPIM, VRL, VIM, VPRL, VPIM

DIMENSION C(6,5,3), R(6,5,3), AL(6,5,3), MP(5,3), MQ(5,3),

1RA(5), RB(5), IORDR(5), ZAKRL(3), ZBKRL(3), ZAKIM(3), ZBKIM(3)

COMMON C, R, AL, MP, MQ, RA, RB, ZAKRL, ZBKRL, ZAKIM, ZBKIM, YARL,

1YBRL, YAIM, YBIM, ZARL, ZBRL, ZAIM, ZBIM, F, F2, IREC, ZPRL, ZPIM,

2VPRL, VPIM, ZRL, ZIM, VRL, VIM, K, I, IPUNC, RAO, RBO, DEN, L,

3INDX, REAL, AIMG, RLDEN, ANMIN, RLNUM, DENIM, A, B, QUAN, IORDR,

4LETR, VAR1, IORD, ZARLK, ZBRLK, ZAIMK, ZBIMK, NSECT, NT

COMMON XK, RK, ZOR, ZOI, VOR, VOI, BETA, DB, XIMAG

REAL = 1.0/RAO

AIMG = RBO

IF FIRST SECTION, GO TO 4022 TO CALCULATE INTIIAL PI PAD DROP.

COME BACK TO CALCULATE FIRST SECTION AND FOLLOWING PI

PAD DROP.

IF(IREC) 4021, 4022, 4021

4021 ZARLK = ZAKRL(I)

ZBRLK = ZBKRL(I)

ZAIMK = ZAKIM(I)

ZBIMK = ZBKIM(I)

REAL = ZARLK+ZBRLK

AIMG = ZAIMK+ZBIMK

RLNUM = ZPRL\*REAL- ZPIM\*AIMG

RLNUM = RLNUM+2.0\*(ZARLK\*ZBRLK-ZAIMK\*ZBIMK)

RLDEN = 2.0\* ZPRL+REAL

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DENIM = 2.0\* ZPIM+AIMG  
ANMIN = ZPIM\*REAL+ ZPRL\*AIMG  
ANMIN = ANMIN+2.0\*(ZARLK\*ZBIMK+ZAIMK\*ZBRLK)  
DEN = 1.0/(RLDEN\*RLDEN+DENIM\*DENIM)  
ZRL = (RLNUM\*RLDEN+DENIM\*ANMIN)\*DEN  
ZIM = (RLDEN\*ANMIN-RLNUM\*DENIM)\*DEN  
REAL = ZBRLK-ZARLK  
AIMG = ZBIMK-ZAIMK

RLDEN = ZPRL\*REAL- ZPIM\*AIMG  
DENIM = ZPIM\*REAL+ ZPRL\*AIMG  
DEN = 1.0/(RLDEN\*RLDEN+DENIM\*DENIM)

A = VPRL\*RLNUM-VPIM\*ANMIN  
B = VPIM\*RLNUM+VPRL\*ANMIN  
VRL = (A\*RLDEN+DENIM\*B)\*DEN  
VIM = (RLDEN\*B-A\*DENIM)\*DEN

REAL = 1.0/RA(K)  
AIMG = RB(K)

4022 DEN = 1.0/(ZRL\*ZRL+ZIM\*ZIM)

RLNUM = 1.0+REAL\*AIMG+AIMG\*ZRL\*DEN

ANMIN = AIMG\*ZIM\*DEN

RLDEN = 2.0\*REAL+AIMG\*REAL\*REAL+(ZRL\*REAL\*AIMG+ZRL)\*DEN

DENIM = (REAL\*AIMG\*ZIM+ZIM)\*DEN

DEN = 1.0/(RLDEN\*RLDEN+DENIM\*DENIM)

ZPRL = (RLNUM\*RLDEN+DENIM\*ANMIN)\*DEN

ZPIM = (-RLDEN\*ANMIN+RLNUM\*DENIM)\*DEN

DEN = 1.0/ (ZRL\*ZRL + ZIM\*ZIM)

VPRL = VRL+VRL\*REAL\*AIMG+(VRL\*AIMG\*ZRL+VIM\*ZIM\*AIMG)\*DEN

VPIM = VIM+VIM\*AIMG\*REAL+(VIM\*AIMG\*ZRL-VRL\*ZIM\*AIMG)\*DEN

4002 RETURN

END

SUBROUTINE ORDER

C COMPUTES TOTAL COMPLEX IMPEDANCE FOR BASIC SECTION IMPEDANCES

C ZA AND ZB FROM INDIVIDUAL ORDERS.

DIMENSION C(6,5,3), R(6,5,3), AL(6,5,3), MP(5,3), MQ(5,3),  
1RA(5), RB(5), IORDR(5), ZAKRL(3), ZBKRL(3), ZAKIM(3), ZBKIM(3)  
COMMON C, R, AL, MP, MQ, RA, RB, ZAKRL, ZBKRL, ZAKIM, ZBKIM, YARL,  
1YBRL, YAIM, YBIM, ZARL, ZBRL, ZAIM, ZBIM, F, F2, IREC, ZPRL, ZPIM,  
2VPRL, VPIM, ZRL, ZIM, VRL, VIM, K, I, IPUNC, RAO, RBO, DEN, L,  
3INDX, REAL, AIMG, RLDEN, ANMIN, RLNUM, DENIM, A, B, QUAN, IORDR,  
4LETT, VAR1, IORD, ZARLK, ZBRLK, ZAIMK, ZBIMK, NSECT, NT  
COMMON XK, RK, ZOR, ZOI, VOR, VOI, BETA, DB, XIMAG

IVAL = I-1

IF(IVAL) 5001, 5002, 5001

5002 ZAKRL(I) = ZARL

ZAKIM(I) = ZAIM

ZBKRL(I) = ZBRL

ZBKIM(I) = ZBIM

GO TO 5006

5001 IF(I-(I/2)\*2) 5007, 5004, 5007

5004 ZAKRL(I) = ZARL+ZAKRL(IVAL)

ZAKIM(I) = ZAIM+ZAKIM(IVAL)

ZBKRL(I) = ZBRL+ZBKRL(IVAL)

ZBKIM(I) = ZBIM+ZBKIM(IVAL)

GO TO 5006

5007 RLNUM = ZAKRL(IVAL)

DENIM = ZAKIM(IVAL)

QUAN =RLNUM\*RLNUM+DENIM\*DENIM

B = YARL\*QUAN+RLNUM

```
A = YAIM*QUAN-DENIM
DEN = 1.0/(B*B+A*A)
ZAKRL(I) = B*QUAN*DEN
ZAKIM(I) = -A*QUAN*DEN
RLNUM = ZBKRL(IVAL)
DENIM = ZBKIM(IVAL)
QUAN =RLNUM*RLNUM+DENIM*DENIM
B = YBRL*QUAN+RLNUM
A = YBIM*QUAN-DENIM
DEN = 1.0/(B*B+A*A)
ZBKRL(I) = B*QUAN*DEN
ZBKIM(I) = -A*QUAN*DEN
```

```
5006 RETURN
```

```
END
```



SUBROUTINE TERMIN

C CALCULATES IMPEDANCE AND VOLTAGE AT INPUT TO INPUT TERMINATION.  
DIMENSION C(6,5,3), R(6,5,3), AL(6,5,3), MP(5,3), MQ(5,3),  
1RA(5), RB(5), IORDR(5), ZAKRL(3), ZBKRL(3), ZAKIM(3), ZBKIM(3)  
COMMON C, R, AL, MP, MQ, RA, RB, ZAKRL, ZBKRL, ZAKIM, ZBKIM, YARL,  
1YBRL, YAIM, YBIM, ZARL, ZBRL, ZAIM, ZBIM, F, F2, IREC, ZPRL, ZPIM,  
2VPRL, VPIM, ZRL, ZIM, VRL, VIM, K, I, IPUNC, RAO, RBO, DEN, L,  
3INDX, REAL, AIMG, RLDEN, ANMIN, RLNUM, DENIM, A, B, QUAN, IORDR,  
4LETR, VAR1, IORD, ZARLK, ZBRLK, ZAIMK, ZBIMK, NSECT, NT  
COMMON XK, RK, ZOR, ZOI, VOR, VOI, BETA, DB, XIMAG  
GO TO (5,6), LETTR  
5 ZTRL = ZPRL  
ZTIM = ZPIM  
GO TO 7  
6 ZTRL = ZRL  
ZTIM = ZIM  
C  
C USE DRPTZ TO CALCULATE INPUT TERMINATION IMPEDANCE.  
C  
7 CALL DRPTZ  
ZBTRL = ZTRL+ZARL  
ZBTIM = ZTIM+ZAIM  
RK = ZARL  
XK = ZAIM  
REAL = VPRL\*ZBTRL-VPIM\*ZBTIM  
XIMAG= VPIM\*ZBTRL+VPRL\*ZBTIM  
DEN = ZTRL\*ZTRL+ZTIM\*ZTIM  
IF(DEN) 1, 2, 1

, 2 PUNCH 101

101 FORMAT (41HERROR,DENOMINATOR IN TERMIN EQUAL TO ZERO)

PAUSE

C

C CALCULATE INPUT VOLTAGE REQUIRED TO PRODUCE SPECIFIED OUTPUT

C TERMINATION VOLTAGE.

C

1 VKRL = (REAL\*ZTRL+XIMAG\*ZTIM)/DEN

VKIM = (ZTRL\*XIMAG-ZTIM\*REAL)/DEN

VPRL = VKRL

VPIM = VKIM

RETURN

END

SUBROUTINE LADDER

DIMENSION C(6,5,3), R(6,5,3), AL(6,5,3), MP(5,3), MQ(5,3),  
IRA(5), RB(5), IORDR(5), ZAKRL(3), ZBKRL(3), ZAKIM(3), ZBKIM(3)  
COMMON C, R, AL, MP, MQ, RA, RB, ZAKRL, ZBKRL, ZAKIM, ZBKIM, YARL,  
1YBRL, YAIM, YBIM, ZARL, ZBRL, ZAIM, ZBIM, F, F2, IREC, ZPRL, ZPIM,  
2VPRL, VPIM, ZRL, ZIM, VRL, VIM, K, I, IPUNC, RAO, RBO, DEN, L,  
3INDX, REAL, AIMG, RLDEN, ANMIN, RLNUM, DENIM, A, B, QUAN, IORDR,  
4LETR, VAR1, IORD, ZARLK, ZBRLK, ZAIMK, ZBIMK, NSECT, NT

COMMON XK, RK, ZOR, ZOI, VOR, VOI, BETA, DB, XIMAG

C COMPUTES RECURSION EQUATIONS FOR LADDER SECTIONS.

C CALCULATES ZRL, ZIM, ZPRL, ZPIM, VRL, VIM, VPRL, VPIM.

ZARLK = ZAKRL(I)

ZAIMK = ZAKIM(I)

ZBRLK = ZBKRL(I)

ZBIMK = ZBKIM(I)

VPRL = VRL

VPIM = VIM

RLDEN = ZARLK+ZRL

DENIM = ZAIMK+ZIM

DEN = RLDEN\*RLDEN+DENIM\*DENIM

RLNUM = ZARLK\*ZRL-ZAIMK\*ZIM

XIMAG = ZAIMK\*ZRL+ZIM\*ZARLK

RNUM = RLNUM\*RLDEN+XIMAG\*DENIM

XMAG = XIMAG\*RLDEN-RNUM\*DENIM

IF(DEN) 1, 2, 1

2 PUNCH 101

101 FORMAT (32HERROR,ZERO DENOMINATOR IN LADDER)

PAUSE

```
1 ZPRL = RNUM/DEN
  ZPIM = XMAG/DEN
  ZRL = ZPRL+ZBRLK
  ZIM = ZPIM+ZBIMK
  DEN = ZPRL*ZPRL+ZPIM*ZPIM
  IF(DEN) 3, 2, 3
3 REAL = 1.0+(ZBRLK*ZPRL+ZBIMK*ZPIM)/DEN
  XIMAG = (ZBIMK*ZPRL-ZPIM*ZBRLK)/DEN
  VRL = VPRL*REAL-VPIM*XIMAG
  VIM = VPRL*XIMAG+VPIM*REAL
  VPRL = VRL
  VPIM = VIM
  RETURN
  END
```

SUBROUTINE DRPTZ

```

C   CALCULATES IMPEDANCE AND ADMITTANCE FOR A TWO-TERMINAL IMPEDANCE
C       (ORDER). IE. YARL, YAIM, YBRL, YBIM, ZARL, ZAIM, ZBRL, ZBIM
C       IE. YARL, YAIM, YBRL, YBIM, ZARL, ZAIM, ZBRL, ZBIM
    DIMENSION C(6,5,3), R(6,5,3), AL(6,5,3), MP(5,3), MQ(5,3),
    IRA(5), RB(5), IORDR(5), ZAKRL(3), ZBKRL(3), ZAKIM(3), ZBKIM(3)
    COMMON C, R, AL, MP, MQ, RA, RB, ZAKRL, ZBKRL, ZAKIM, ZBKIM, YARL,
    IYBRL, YAIM, YBIM, ZARL, ZBRL, ZAIM, ZBIM, F, F2, IREC, ZPRL, ZPIM,
    ZVPRL, VPIM, ZRL, ZIM, VRL, VIM, K, I, IPUNC, RAO, RBO, DEN, L,
    3INDX, REAL, AIMG, RLDEN, ANMIN, RENUM, DENIM, A, B, QUAN, IORDR,
    4LETR, VAR1, IORD, ZARLK, ZBRK, ZAIMK, ZBIMK, NSECT, NT
    COMMON XK, RK, ZOR, ZOI, VOR, VOI, BETA, DB, XIMAG
    IVAL = MP(K,I)
    YARL = 0.0
    YAIM = 0.0
    L = 1
11 IF (IVAL-L) 9, 10, 10
10 INDX = L*2-1
    VAR1 = R(INDX,K,I)
    QUAN = F2*C(INDX,K,I)
    IF (QUAN) 3, 4, 3
4  QUAN = AL(INDX,K,I)*F
    GO TO 5
3  QUAN = (AL(INDX,K,I)-1.0/QUAN)*F
5  DEN = 1.0/(VAR1*VAR1+QUAN*QUAN)
    YARL = YARL+VAR1*DEN
    YAIM = YAIM-QUAN*DEN
    L = L+1

```

```

GO TO 11
9 DEN = YARL*YARL + YAIM*YAIM
  IF (DEN) 12, 13, 12
12 DEN = 1.0/DEN
13 ZARL = YARL*DEN
  ZAIM = -YAIM*DEN
  IVAL = MQ(K,I)
  YBRL = 0.0
  YBIM = 0.0
  L = 1
16 IF (IVAL-L) 14, 15, 15
15 INDX = 2*L
  VAR1 = R(INDX,K,I)
  QUAN = F2*C(INDX,K,I)
  IF (QUAN) 6, 7, 6
7 QUAN = AL(INDX,K,I)*F
  GO TO 8
6 QUAN = (AL(INDX,K,I)-1.0/QUAN)*F
8 DEN = 1.0/(VAR1*VAR1+QUAN*QUAN)
  YBRL = YBRL+VAR1*DEN
  YBIM = YBIM-QUAN*DEN
  L = L+1
  GO TO 16
14 DEN = YBRL*YBRL + YBIM*YBIM
  IF (DEN) 17, 18, 17
17 DEN = 1.0/DEN
18 ZBRL = YBRL*DEN
  ZBIM = -YBIM*DEN

```

RETURN

END

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1620 USERS GROUP  
WESTERN REGION MEETING

June 18, 1964

FORTRAN II - DEBUGGING TECHNIQUES AND AIDS

Leon P. Goldberg  
Technical Staff  
Princeton University



FORTRAN II

FLOATING HARDWARE  
NON RELOCABLE SUBROUTINES

<u>Address</u>	<u>Subroutine name</u>	<u>Function</u>
01510	SWC	I/O
01768	COMPLT	I/O
03158	RATY	I/O
03182	RAPT	I/O
01418	RACD	I/O
01574	SLASH	I/O
00986	WATY	I/O
01022	WAPT	I/O
01058	WACD	I/O
01800	HTYPE	Hollerith conversion
02052	REDO	Multiple field specs.
02152	REP	Multiple parenthesized specs.
02380	ITYPE	I specification
03280	FTYPE	F specification
03300	EITYPE	E specification
06020	XTYPE	X specification
06052	ATYPE	A specification
06528	MATRIX	Reading arrays
07316	FXA	I+J
07348	FXSR	-(J*K)+I
07416	FXS	I-J
07440	FXM	I*J
07484	FXD	I/J
07570	FXDR	1/(I/J)
07604	RSGN	-I or -A
07698	FLOAT	A=I
07932	FLX	I=A
08152	FLXI	I**J
08586	FAXI	A**I
09044	FAXB	A**B
09356	FMFAC	FAC to A
09504	FSB	A-B
09528	FAD	A+B
09740	FSBR	-(A*B)+C
09808	FMP	A*B
09856	FD	A/B
09952	FDVR	1/(A/B)
10000	TOPAC	A to FAC
06710	TRACE	
02256	REPEAT (REP)	
00485	FAC	

C INTERACTION OF COMMON AND EQUIVALENCE IN UNDIMENSIONED VARIABLES.  
 COMMON X, Y, Z  
 EQUIVALENCE (X,A), (Y,B,C), (Z,L)  
 A=4.1527341  
 B=2.\*A  
 R=X+2.  
 S=B+R  
 Z=X+Y+S  
 STOP  
 END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

11043 41527341  
 11051 20000001  
 59999 X  
 59999 A  
 59991 Y  
 59991 B  
 59991 C  
 59983 Z  
 59983 L  
 11059 R  
 11067 S  
 END OF PASS I

171000011043  
 170671059999  
 171000011051  
 170980859999  
 170671059991  
 171000059999  
 170952811051  
 170671011059  
 171000059991  
 170952811059  
 170671011067  
 171000059999  
 170952859991  
 170952811067  
 170671059983  
 340000000102  
 390041700100  
 480000000000

C COMMON STORAGE IN DIMENSIONED VARIABLES.  
DIMENSION X(5), Y(2,5), Z(3,4,5)  
COMMON X,Y  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

59959 X 59999

59859 Y 59949

11045 Z 11635

END OF PASS 1

```

C   EXAMPLES TO DEMONSTRATE THE COMPLEX INTERACTION OF EQUIVALENCE AND
C   COMMON STORAGE ASSIGNMENT IN DIMENSIONED VARIABLES.
    DIMENSION X(5), Y(2,5), Z(3,4,5)
    COMMON X
    EQUIVALENCE (Y,Z)
    STOP
    END
TURN SW 1 ON FOR SYMBOL TABLE, PRESS START
59959      X 59999
11045      Y 11135
11045      Z 11635
END OF PASS 1

```

```

ENTER SOURCE PROGRAM, PRESS START
    DIMENSION X(5), Y(2,5), Z(3,4,5)
    COMMON X
    EQUIVALENCE (X(5),Y(10),Z(60))
    STOP
    END
TURN SW 1 ON FOR SYMBOL TABLE, PRESS START
59959      X 59999
59909      Y 59999
59409      Z 59999
END OF PASS 1

```

```

ENTER SOURCE PROGRAM, PRESS START
    DIMENSION X(5), Y(2,5), Z(3,4,5)
    EQUIVALENCE (Y(10),X(5),Z(60))
    STOP
    END
TURN SW 1 ON FOR SYMBOL TABLE, PRESS START
11045      Z 11635
11545      Y 11635
11595      X 11635
END OF PASS 1

```

C EXAMPLES TO DEMONSTRATE THE COMPLEX INTERACTION OF EQUIVALENCE AND  
C COMMON STORAGE ASSIGNMENT IN DIMENSIONED VARIABLES.  
DIMENSION X(5), Y(2,5), Z(3,4,5)  
EQUIVALENCE (Z(60),X(5), Y(10))  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

11045        Z 11635  
11595        X 11635  
11545        Y 11635  
END OF PASS 1

ENTER SOURCE PROGRAM, PRESS START  
DIMENSION X(5), Y(2,5), Z(3,4,5)  
COMMON X  
EQUIVALENCE (X,Y,Z)  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

59959        X 59999  
ERROR 55  
59959        Y 60049  
ERROR 55  
59959        Z 60549  
END OF PASS 1

ENTER SOURCE PROGRAM, PRESS START  
DIMENSION X(5), Y(2,5), Z(3,4,5)  
COMMON Z  
EQUIVALENCE (X,Y,Z)  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

59409        Z 59999  
59409        X 59449  
59409        Y 59499  
END OF PASS 1

ENTER SOURCE PROGRAM, PRESS START  
DIMENSION X(5), Y(2,5), Z(3,4,5)  
COMMON Y  
EQUIVALENCE (X(3),Y(6),Z(15))  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

59909        Y 59999  
59909        X 59979  
ERROR 55  
59819        Z 60409  
END OF PASS 1

C EQUIVALENCE STORAGE ASSIGNMENT.  
 EQUIVALENCE (R,S),(T,U,V).(R,L)  
 EQUIVALENCE (A,B,C,D)  
 R=1.5  
 A=R+3.  
 L=B+1.  
 STOP  
 END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

T1043	T5000001
T1051	30000001
T1059	10000001
T1067	R
T1067	S
T1067	L
T1075	T
T1075	U
T1075	V
T1083	A
T1083	B
T1083	C
T1083	D

END OF PASS 1

171000011043
170671011067
171000011067
170952811051
170671011083
171000011083
170952811059
170793200000
170671011067
340000000102
390041700100
480000000000

C COMMON STORAGE ASSIGNMENT.

COMMON W,X,Y,Z

W=1.5

X=2.0

Z=3.0

Y=2.5

A=W+X+Y+Z+0.5

L=4. + 5.

STOP

END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

11047	150000000001
11059	200000000001
11071	300000000001
11083	250000000001
11095	500000000000
11107	400000000001
11119	500000000001
59999	W
59987	X
59975	Y
59963	Z
11131	A
11141	L

END OF PASS I

171000011047
170671059999
171000011059
170671059987
171000011071
170671059963
171000011083
170671059975
171000059999
170952859987
170952859975
170952859963
170952811095
170671011131
171000011107
170952811119
170793200000
170671011141
340000000102
390041700100
480000000000

C GENERAL I/O  
READ 1, X, Y, L  
1 FORMAT(F10.3, E15.8, I5, 5X, 3HXY=)  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

11045 0001  
11055 X  
11065 Y  
11070 L  
END OF PASS 1

270141811040  
170151011055  
170151011065  
170151011070  
170176800000  
491104500328  
001003033000  
150802380010  
060200100180  
000607083302  
052111393400  
000001023900  
417001004800  
00000000491



READ I, X, Y1, Y2, Y3  
1 FORMAT (A4, 4(F10.3))  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

T1045 0001  
T1055 X  
T1065 Y1  
T1075 Y2  
T1085 Y3  
END OF PASS 1

2701418T1040  
1701510T1055  
1701510T1065  
1701510T1075  
1701510T1085  
170176800000  
49T104500605  
200803280010  
0302256T1178  
0402052T1173  
340000000102  
390041700100  
480000000000

C I/O OF MATRICES WITHOUT THE IMPLIED DO LOOP.  
DIMENSION A(15)  
READ 1, A  
1 FORMAT (F10.3)  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

T1045 0001  
T1055 A T1195  
END OF PASS 1

2701418T1040  
160620300015  
160626300000  
1706528T1045  
170176800000  
49T104500328  
00100302052T  
126334000000  
010239004170  
010048000000  
0000

```

C      I/O OF MATRICES WITH THE IMPLIED DO LOOP.
      DIMENSION X(20)
      READ 1, (X(I),I=1,20)
      1 FORMAT (8F10.3)
      STOP
      END

```

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

```

T1040 00001
T1045 00020
T1055 00001
T1065      X T1255
T1260      I
END OF PASS 1

```

```

2701418T1050
26T1260T1040
T3T1260000T0
1T00099T1055
320009500000
270151000099
T1T126000001
24T1260T1045
47T128601100
170176800000
49T105500328
00100302152T
13940802052T
138934000000
010239004170
010048000000
0000

```

C      SAMPLE PROBLEM.  
X = A + B + C\*D  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

11045	X
11055	A
11065	B
11075	C
11085	D

END OF PASS I

171000011055  
170952811065  
170935611139  
491114000000  
000000171000  
011075170980  
811085170952  
811139170671  
011045340000  
000102390041  
700100480000  
0000004

ENTER SOURCE PROGRAM, PRESS START  
C      SAMPLE PROBLEM.  
X = C\*D + A + B  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

11045	X
11055	C
11065	D
11075	A
11085	B

END OF PASS I

171000011055  
170980811065  
170952811075  
170952811085  
170671011045  
340000000102  
390041700100  
480000000000

C SAMPLE PROBLEM.  
X = A\*B\*C\*\*4  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

T1040 00004  
T1050 X  
T1060 A  
T1070 B  
T1080 C  
END OF PASS I

17T0000T1060  
1709808T1070  
1709356T1135  
49T113600000  
00000017T1000  
0T1080170858  
6T1040170980  
8T1135170671  
0T1050340000  
000102390041  
700100480000  
00000049T

C SAMPLE PROBLEM.  
X = C\*\*4 \*A\*B  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

I1040 00004  
I1050 X  
I1060 C  
I1070 A  
I1080 B  
END OF PASS 1

17I0000I1060  
1708586I1040  
1709808I1070  
1709808I1080  
1706710I1050  
340000000102  
390041700100  
480000000000

C SAMPLE PROBLEM.  
X = C\*\*4. \*A\*B  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

T1045 4000000001  
T1055 X  
T1065 C  
T1075 A  
T1085 B  
END OF PASS I

17T0000T1065  
1709044T1045  
1709808T1075  
1709808T1085  
1706710T1055  
340000000102  
390041700100  
480000000000

C SAMPLE PROBLEM.  
X = A\*B - C\*D  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

T1045 X  
T1055 A  
T1065 B  
T1075 C  
T1085 D  
END OF PASS I

17T0000T1055  
1709808T1065  
1709356T1139  
49T114000000  
17067117T000  
0T1075170980  
8T1085170974  
0T1139170671  
0T1045340000  
000102390041  
700100480000  
000000491



C SAMPLE PROBLEM.  
TEMP = A \* B  
X = C \* D - TEMP  
STOP  
END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

T1045 TEMP  
T1055 A  
T1065 B  
T1075 X  
T1085 C  
T1095 D  
END OF PASS I

1710000T1055  
1709808T1065  
1706710T1045  
1710000T1085  
1709808T1095  
1709504T1045  
1706710T1075  
340000000102  
390041700100  
480000000000

C SIMPLE SINGLE, DOUBLE AND TRIPLE SUBSCRIPTION EXAMPLE.  
 DIMENSION A(5), B(2,4), C(2,2,2)  
 EQUIVALENCE (A3, A(3)), (B2, B(2,3)), (C1, C(1,2,1))

ERROR 05  
 B2 = 23.  
 C1 = 121.  
 A2 = B2 + C1  
 STOP  
 END

END OF PASS I  
 ENTER SOURCE PROGRAM, PRESS START

C SIMPLE SINGLE, DOUBLE AND TRIPLE SUBSCRIPTION EXAMPLE.  
 DIMENSION A(5), B(2,4), C(2,2,2)  
 EQUIVALENCE (A3, A(3)), (B2, B(7)), (C1, C(3))

B2 = 23.  
 C1 = 121.  
 A2 = B2 + C1  
 STOP  
 END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

T1045 2300000002  
 T1055 1210000003  
 T1065 A T1105  
 T1085 A3  
 T1115 B T1185  
 T1175 B2  
 T1195 C T1265  
 T1215 C1  
 T1275 A2

END OF PASS I  
 CARDS NOT IN ORDER  
 SW1 ON TO PUNCH SUBROUTINES, PRESS START  
 END OF PASS II  
 SW1 ON TO PUNCH SUBROUTINES, PRESS START  
 END OF PASS III

T3T104500002  
 210009911040  
 260003900099  
 130003900010  
 110009911100  
 320009500000  
 260048500099  
 170935611395  
 491139600000  
 000000171000  
 011055270671  
 011395131106  
 000002210009  
 911040260003  
 900099130003  
 900002210009  
 911060260003  
 900099130003  
 900010110009  
 911140320009  
 500000260048  
 500099170935  
 611395171000  
 011070270671  
 011395131104  
 500010110009  
 911070320009  
 500000260048  
 500099170935  
 611395131104  
 500002210009  
 911040260003  
 900099130003  
 900010110009  
 911100320009  
 500000271000  
 000099131106  
 000002210009  
 911040260003  
 900099130003  
 900002210009  
 911060260003  
 900099130003  
 900010110009  
 911140320009  
 500000170935  
 611857491185  
 800000000000  
 271000000099  
 170952811857  
 270671011395  
 340000000102  
 390041700100  
 480000000000

C SIMPLE SINGLE, DOUBLE AND TRIPLE SUBSCRIPTION EXAMPLE.  
 DIMENSION A(5), B(2,4), C(2,2,2)  
 B(2,3) = 23.  
 C(1,2,1) = 121.  
 A(3) = B(2,3) + C(1,2,1)  
 STOP  
 END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

T1040 00002  
 T1045 00003  
 T1055 2300000002  
 T1060 00001  
 T1070 1210000003  
 T1080 A T1120  
 T1130 B T11200  
 T1210 C T11280  
 END OF PASS I

17T0000T1045  
 1706710T11175  
 17T0000T1055  
 1706710T11215  
 17T0000T11175  
 1709528T11215  
 1706710T11275  
 340000000102  
 390041700100  
 480000000000

```

C   SAMPLE PROBLEM SHOWING POLYNOMIAL EXPRESSION W/O NESTING.
    DIMENSION X(10),A(20)
    READ 1, NPTS,(X(I),I=1,NPTS)
    DO 30 I=1,20
      AI = I
    30 A(I) = AI
      DO 20 I=1,NPTS
        SUM = 0.0
        XX = X(I)
        DO 10 K=1,20
          KK= K-1
          AK= A(K)
        10 SUM = SUM + AK *XX**KK
        20 PRINT 2, SUM,I,XX
          2 FORMAT (2X,4HSUM=E15.8,7HWHEN X(12,2H)=E15.8)
          1 FORMAT (12/8E10.3)
        STOP
      END

```

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

```

11040 00001
11045 00020
11055 0000000099
11065 00001
11070 0030
11075 0020
11080 0010
11090 0002
11100      X 11190
11200      A 11390
11395      NPTS
11400      I
11410      AI
11420      SUM
11430      XX
11435      K
11440      KK
11450      AK
END OF PASS 1

```

C SAMPLE PROBLEM SHOWING POLYNOMIAL EXPRESSION W. NESTING.

DIMENSION X(10),A(20)  
 READ 1, NPTS, (X(I), I=1, NPTS)

DO 30 I=1,20

AI = I

30 A(I) = AI

DO 20 I=1, NPTS

XX = X(I)

PROD = A(20)\*XX + A(19)

DO 10 K=3,20

MK=21-K

10 PROD = PROD\*XX+ A(MK)

20 PRINT 2, PROD, I, XX

2 FORMAT (2X, 4HSUM=E15.8, 7HWHEN X(12,2H)=E15.8)

1 FORMAT (12/8E10.3)

STOP

END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

I1040	00001	
I1045	00020	
I1050	00019	
I1055	00003	
I1060	00021	
I1070	0001	
I1075	0030	
I1080	0020	
I1085	0010	
I1095	0002	
I1105	X	I1195
I1205	A	I1395
I1400	NPTS	
I1405	I	
I1415	AI	
I1425	XX	
I1435	PROD	
I1440	K	
I1445	MK	

END OF PASS 1

C SAMPLE PROBLEM SHOWING COMPUTED AND UNCONDITIONAL GO TO @S.

L=1

3 Y=4

X=2

A=X+Y

GO TO (1,2),L

1 L=2

GO TO 3

2 STOP

END

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

I1040 00001  
I1045 00004  
I1050 00002  
I1055 00003  
I1060 00001  
I1065 00002  
I1070 L  
I1080 Y  
I1090 X  
I1100 A  
END OF PASS I

LOAD SUBROUTINES

1710000I1040  
1706710I1070  
1710000I1045  
170769800000  
1706710I1080  
1710000I1050  
170769800000  
1706710I1090  
1710000I1090  
1709528I1080  
1706710I1100  
131107000005  
1200099I1264  
4900099I1060  
I10650171000  
0I1050170671  
0I107049I105  
500000340000  
000102390041  
700100480000  
00000049

```
C    SAMPLE PROBLEM USING A DO LOOP
    DIMENSION X(10)
    DO 5 I=1,10
5    X(I) = I
    STOP
    END
```

```
TURN SW 1 ON FOR SYMBOL TABLE, PRESS
RT          T          S
```

```
T1040 00001
T1045 00010
T1050 0005
T1060      X T1150
T1155      I
END OF PASS I
```

```
26T1155T1040
13T115500010
1T00099T1050
320009500000
260048500099
1709356T11245
49T1124600000
00000517T000
0T1155170769
800000270671
0T11245T11115
50000124T1115
5T104547T1116
801100340000
000102390041
700100480000
00000049
```

```

C      SAMPLE PROBLEM USING IF STATEMENT
      IF (A) 1,2,3
1     B= -1.0
      GO TO 4
2     B= 0.0
      GO TO 4
3     B= 1.0
4     STOP
      END

```

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

```

11045 1000000001
11055 0000000099
11060 0001
11065 0002
11070 0003
11075 0004
11085      A
11095      B
END OF PASS I

```

```

171000011085
170671000485
431114400476
491106500000
441107000483
491106001710
000110451707
604000001706
710110954911
075000001710
000110551706
710110954911
075000001710
000110451706
710110953400
000001023900
417001004800
00000000491

```



C SAMPLE PROBLEM USING IF STATEMENT

```
B= 1.0  
IF (A) 1,2,4  
1 B= -1.0  
GO TO 4  
2 B= 0.0  
4 STOP  
END
```

TURN SW 1 ON FOR SYMBOL TABLE, PRESS START

```
11045 1000000001  
11055 0000000099  
11060 0001  
11065 0002  
11070 0004  
11080 B  
11090 A  
END OF PASS. I
```

```
171000011045  
170671011080  
171000011090  
170671000485  
431116400476  
491106500000  
441107000483  
491106001710  
000110451707  
604000001706  
710110804911  
070000001710  
000110551706  
710110803400  
000001023900  
417001004800  
0000000049
```

REFERENCES

Private Consultations with - - - - - L. Hoffman  
P. Larrea

FORTRAN II AND THE 1443

Lanny L. Hoffman

PRINTER FN II SUB. LOCATIONS.

---

00485	FAC	FLOATING ACCUMULATOR
01314	TUFAC	A TO FAC
01358	FMFAC	FAC TO A
01402	TRACE	PRINT FAC, GO TO FMFAC
01482	FXA	I+J
01514	FXS	I-J
01538	FXSR	-(J*K)+I
01562	XRSGN	-I
01608	FXM	I*J
01652	FXD	I/J
01738	FXDR	1/(I/J)
01772	FIX	I=A
02112	FLUAT	A=I
02336	RSGN	-A
02336	FSBR	-(A*B)+C
02430	FAD	A+B
03002	FSB	A-B
03046	FMP	A*B
03218	FD	A/B
03434	FDVR	1/(A/B)
03538	FIXI	I**J
03960	FAXI	A**I
04550	FAXB	A**B
04920	WATY	I/O
04992	WACD	I/O
05028	PRA	I/O ( PRINT ALPH )
05304	RATY	I/O
05364	RACD	I/O
05408	SLASH	I/O
05650	COMPLT	I/O
05690	REP	MULTIPLE PARENTHESIZED SPECS.
06066	SWC	I/O
06122	MATRIX	READING ARRAYS
06280	REDO	MULTIPLE FIELD SPECS.
05794	XTYPE	X SPECIFICATION
05826	HTYPE	H SPECIFICATION
06350	ATYPE	A SPECIFICATION
06842	ITYPE	I SPECIFICATION
07830	FTYPE	F SPECIFICATION
07850	ETYPE	E SPECIFICATION
01976	ERR MESS	PRINT ERROR MESSAGE
01924	PASS I	PASS I INIT.
00684	PASS II	PASS II INIT.
01938	SPS I	SPS PASS I INIT.

```
SUBROUTINE X(A,B,C,D)
C=A+B
D=A-B
RETURN
END
```

SUBPROGRAM

```
CALL X(ALFA,BETA,GAMMA,DELTA)
```

MAIN LINE

```
BTM XCELL,*,+11,6
DSA ALFA,BETA,GAMMA,DELTA
```

CALLING SEQUENCE

SUBPROGRAM  
GENERATED  
MATERIALS

```
XPGM DS 5
AM XPGM-1,5,010
TF A,XPGM-1,0111
BNF *,+36,A,01
CF A,,0
TF A,A,0111
AM XPGM-1,5,010
TF B,XPGM-1,0111
BNF *,+36,B,01
CF B,,0
TF B,B,0111
AM XPGM-1,5,010
TF C,XPGM-1,0111
BNF *,+36,C,01
CF C,,0
TF C,C,0111
AM XPGM-1,5,010
TF D,XPGM-1,0111
BNF *,+36,D,01
CF D,,0
TF D,D,0111
AM XPGM-1,1,010
BT TOFAC,A,1
BT FAD,B,1
BT FRMFAC,C,1
BT TOFAC,A,1
BT FSB,B,1
BT FRMFAC,D,1
B XPGM-1,,06
```

-----  
FLAG CONVENTION.....

IF P IS RELOCATABLE, FLAG OPERAND IS 0  
IF Q IS RELOCATABLE, FLAG OPERAND IS 1

NORMAL FLAG OPERANDS ARE STILL IN EFFECT FOR  
IMMEDIATE AND INDIRECT ADDRESSING.  
FLAGS ARE USED OVER THE OPERATION CODE TO DENOTE RELOCATION TO  
THE LOADER. THESE FLAGS DO NOT ALTER THE OPERATION OF  
THE INSTRUCTION.

```
C TEST PROGRAM FOR 1443 PRINTER PLOT, SIN (X) VS. X.  
C  
  DIMENSION X(500),Y(500)  
  PRINT 10  
10 FORMAT(1H1)  
  T=0.  
  DO 1 I=1,200  
  X(I)=T  
  Y(I)=SINF(T)  
  T=T+.01*3.14159  
1 CONTINUE  
  PAUSE  
  CALL PLOT(X,10.,0.,5,Y,1.,-1.,10,200)  
  STOP  
  END
```

```

C      FORTRAN SUBROUTINE FOR 1443 PRINTER PLOTTING,
C      BY L. HOFFMAN, GUGGENHEIM LABS.
C
SUBROUTINE PLOT(X,XMAX,XMIN,NX,Y,YMAX,YMIN,NY,N)
DIMENSION DUMMY(2),OUTPUT(102),X(2),Y(2)
XCHAR=.20
YCHAR=.71
CHAR=.14
BLANK=0.
XNO=100.
NOX=XNO+1.
YLABEL=YMAX
DX=(XMAX-XMIN)/XNO
DY=(YMAX-YMIN)/50.
C      MOVE MAX DOWN BY ONE-HALF BOX....
      YYMAX=YMAX+.5*DY
      XXMIN=XMIN-.5*DX
C
      KY=0
      NX1=NX+1
      DO 1 I=1,51
C      CALL INIT(OUTPUT,BLANK)
      DO 111 II2=1,NOX
111    OUTPUT(II2)=BLANK
C      CALL GRID(OUTPUT,DX,DY,NX,NY,KY,I,XCHAR,YCHAR,IND)
      IND=0
      IF(I-1-50*KY/NY)211,222,211
222    DO 332 JJ=1,NOX
332    OUTPUT(JJ)=XCHAR
      IND=1
      KY=KY+1
211    DO 444 JJ=1,NX1
      I2=((JJ-1)*(NOX-1))/NX
444    OUTPUT(I2+1)=YCHAR
      ZI=I
      UP=YYMAX-(ZI-1.)*DY
      DOWN=UP-DY
C      CALL FINDY(X,Y,UP,DOWN,OUTPUT,N,DX,DY,XMAX,XMIN,CHAR)
      DO 1121 IF=1,N
      IF(Y(IF)-UP)2221,1121,1121
2221    IF(Y(IF)-DOWN)1121,3331,3331
3331    CONTINUE
      JJ=(X(IF)-XXMIN)/DX
      JJ=JJ+1
      OUTPUT(JJ)=CHAR
1121    CONTINUE
      8 IF(IND)10,10,11
      10 PRINT 2,(OUTPUT(J),J=1,NOX)
      2 FORMAT(12X,50A1,51A1)
      GO TO 12
      11 PRINT 3,YLABEL,(OUTPUT(J),J=1,NOX)
      3 FORMAT(1X,E10.3,1X,50A1,51A1)
      12 YLABEL=YLABEL-DY
      1 CONTINUE
      RETURN
      END

```

ASSEMBLY AND FINAL PHASE OF SPS SUBS. FOR FN II.

- 1) USE 1620/1710 SPS TO ASSEMBLE AND COMPRESS THE SPS PROGRAM.
- 2) REMOVE THE FIRST TWO (2) AND THE LAST SEVEN (7) CARDS FROM THE COMPRESSED DECK. (THIS DOES NOT INCLUDE THE TWO BLANK CARDS AT THE END OF THE DECK)
- 3) ADD HEADER CARD AS NO. 1.
- 4) ADD TRAILER CARD TO END OF DECK.
- 5) CORRECT DSA'S , IF ANY, USED IN THE SPS PROGRAM, OTHERWISE, GO TO 7.
- 6) PUNCH A FLAGGED ZERO IN COLUMN 62 OF ALL OBJECT DECK CARDS PRODUCED BY DSA'S IN SPS PROGRAM.
- 7) CHECK FOR RELOCATABLE CONSTANTS, IF NONE, THEN GO TO 9.
- 8) PUNCH A FLAGGED 1 IN COLUMN 62 OF ALL CONSTANTS NOT TO BE RELOCATED.
- 9) PUNCH NEW CARD NO. IN TRAILER CARD TO CONTINUE SEQUENCING.
- 10) THE DECK CAN NOW BE USED WITH A FORTRAN CALL STATEMENT.

THE HEADER CARD.....

COLS. 1-12	SUBROUTINE NAME IN TWO-DIGIT ALPHANUMERIC FORM WITH FLAG OVER HIGH ORDER DIGIT AND RIGHT JUSTIFIED.
COLS. 13-20	BLANK
COLS. 21-22	FF, LENGTH OF FLOATING MANTISSA, FLAG OVER HIGH ORDER DIGIT.
COLS. 23-24	KK, LENGTH OF FIXED MANTISSA, FLAG OVER HIGH ORDER DIGIT.
COLS. 25-62	BLANK
COLS. 63	RECORD MARK (0-2-8)
COLS. 64-80	BLANK, EXCEPT FLAG IN COL. 78

THE TRAILER CARD.....

COLS. 1-62	BLANK
COLS. 63	FLAGGED 1
COLS. 64-80	BLANK, EXCEPT FOR CARD NO. IN COL. 78-80.



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\*  
\*

PLOT FOR 1443 PRINTER, L. HOFFMAN GUGGENHEIM LABS.

AN EXAMPLE OF AN SPS SUBROUTINE FOR FN II .

```
11036          4/24          DORG 11036
11040          00005        AUTO PLOT DS 5
11045          00005        AUTO X DS 5
11050          00005        AUTO XMAX DS 5
11055          00005        AUTO XMIN DS 5
11060          00005        AUTO NX DS 5
11065          00005        AUTO Y DS 5
11070          00005        AUTO YMAX DS 5
11075          00005        AUTO YMIN DS 5
11080          00005        AUTO NY DS 5
11085          00005        AUTO N DS 5
11092          00007        AUTO DS 7

*
*          END OF ARGUMENT ADDRESSES/
11094 J1 11093 000-5 AUTO START AM START-1,5,010
11106 KO 11045 1109L AUTO TF X ,START-1,0111
11118 MM 11154 11045 AUTO BNF *+36,X ,01
11130 L3 11045 00000 AUTO CF X ,,0
11142 KO 11045 1104N AUTO TF X ,X ,0111
11154 J1 11093 000-5 AUTO AM START-1,5,010
11166 KO 11050 1109L AUTO TF XMAX ,START-1,0111
11178 MM 11214 11050 AUTO BNF *+36,XMAX ,01
11190 L3 11050 00000 AUTO CF XMAX ,,0
11202 KO 11050 1105- AUTO TF XMAX ,XMAX ,0111
11214 J1 11093 000-5 AUTO AM START-1,5,010
11226 KO 11055 1109L AUTO TF XMIN ,START-1,0111
11238 MM 11274 11055 AUTO BNF *+36,XMIN ,01
11250 L3 11055 00000 AUTO CF XMIN ,,0
11262 KO 11055 1105N AUTO TF XMIN ,XMIN ,0111
11274 J1 11093 000-5 AUTO AM START-1,5,010
11286 KO 11060 1109L AUTO TF NX ,START-1,0111
11298 MM 11334 11060 AUTO BNF *+36,NX ,01
11310 L3 11060 00000 AUTO CF NX ,,0
11322 KO 11060 1106- AUTO TF NX ,NX ,0111
11334 J1 11093 000-5 AUTO AM START-1,5,010
11346 KO 11065 1109L AUTO TF Y ,START-1,0111
11358 MM 11394 11065 AUTO BNF *+36,Y ,01
11370 L3 11065 00000 AUTO CF Y ,,0
11382 KO 11065 1106N AUTO TF Y ,Y ,0111
11394 J1 11093 000-5 AUTO AM START-1,5,010
11406 KO 11070 1109L AUTO TF YMAX ,START-1,0111
11418 MM 11454 11070 AUTO BNF *+36,YMAX ,01
11430 L3 11070 00000 AUTO CF YMAX ,,0
11442 KO 11070 1107- AUTO TF YMAX ,YMAX ,0111
11454 J1 11093 000-5 AUTO AM START-1,5,010
11466 KO 11075 1109L AUTO TF YMIN ,START-1,0111
11478 MM 11514 11075 AUTO BNF *+36,YMIN ,01
11490 L3 11075 00000 AUTO CF YMIN ,,0
11502 KO 11075 1107N AUTO TF YMIN ,YMIN ,0111
11514 J1 11093 000-5 AUTO AM START-1,5,010
11526 KO 11080 1109L AUTO TF NY ,START-1,0111
11538 MM 11574 11080 AUTO BNF *+36,NY ,01
11550 L3 11080 00000 AUTO CF NY ,,0
11562 KO 11080 1108- AUTO TF NY ,NY ,0111
11574 J1 11093 000-5 AUTO AM START-1,5,010
11586 KO 11085 1109L AUTO TF N ,START-1,0111
11598 MM 11634 11085 AUTO BNF *+36,N ,01
11610 L3 11085 00000 AUTO CF N ,,0
```

11622	KU	11085	11085	AUTO	TF	N	,N	,0111
11634	J1	11093	000-2	AUTO	AN	START-1,2,010		
11646	M9	12074	00000	AUTO	B	AROUND,,0		
11654				AUTO	DORG	*-3		
				*		SYMBOL TABLE AND CONSTANTS HERE		
00010		00000		3/17	FF	DS	,10	
00005		00000		3/17	KK	DS	,5	
00485		00000		3/17	FAC	DS	,485	
01402		00000		4/23	TRACE	DS	,1402	
01652		00000		4/23	FXD	DS	,1652	
02336		00000		4/23	RSGN	DS	,2336	
02112		00000		4/23	FLOAT	DS	,2112	
01772		00000		4/23	FIX	DS	,1772	
01358		00000		4/23	FIXFAC	DS	,1358	
03002		00000		4/23	FSB	DS	,3002	
02430		00000		4/23	FAD	DS	,2430	
02336		00000		5012	FSBR	DS	,2336	
03046		00000		4/23	FMP	DS	,3046	
03218		00000		4/23	FDV	DS	,3218	
03434		00000		4/23	FDVR	DS	,3434	
01314		00000		4/23	TOFAC	DS	,1314	
05028		00000		4/23	WATY	DS	,5028	
06066		00000		4/23	SMC	DS	,6066	
05650		00000		4/23	COMPLT	DS	,5650	
11658		00005		4/23	FMT	DC	5,5794	
11661		00003		4/23		DC	3,002	
11666		00005		4/23		DC	5,7850	
11669		00003		4/23		DC	3,010	
11671		00002		4/23		DC	2,03	
11676		00005		4/23		DC	5,5794	
11679		00003		4/23		DC	3,002	
11684		00005		4/24	REDO	DC	5,6280	
11687		00012			DUMOUT	DAC	12,	
11711		00102		3/13	OUTPUT	DAS	102	
11915		00001				DAC	1,'	
11920		00005		3/22	I	DS	KK	
11924		00004		3/22	JJ	DS	4	
11927		00003		3/25	II2	DS	3	
11932		00005		3/22	I2	DS	5	
11937		00005		3/22	IF	DS	5	
11947		00010		3/22	ZI	DS	FF	
11957		00010		3/22	UP	DS	FF	
11967		00010		3/22	DOWN	DS	FF	
11972		00005		3/22	COUNT	DS	5	
11980		00008		3/22	HALF1	DC	8,50000000	
11982		00002		3/22	HALF	DC	2,00	
11992		00010		3/22	YYMAX	DS	FF	
12002		00010		3/22	DY	DS	FF	
12012		00010		3/22	DX	DS	FF	
12022		00010		3/22	YLABEL	DS	FF	
12023		00001		3/22	IND	DS	1	
12028		00005		3/22	NDX	DS	KK	
12036		00008		*****	XNO1	DC	8,10000000	
12038		00002		*****	XNO	DC	2,03	
12046		00008		3/22	ONE1	DC	8,10000000	
12048		00002		3/22	ONE	DC	2,01	
12056		00008		3/22	F1	DC	8,50000000	
12058		00002		3/22	FIFTY	DC	2,02	
12063		00005		3/22	KY	DS	KK	
12073		00010		3/25	XXMIN	DS	FF	
12074	1P	01314	J2038	3/22	AROUND	BTM	TOFAC,XNO,17	

12086	1P	02430	J2048	3/22		BTM	FAD,ONE,17
12098	17	01772	000-0	3/22		BTM	FIX,0,10
12110	K6	12028	00485	3/22		TF	NOX,FAC,0
12122	2P	01314	11070	3/22		BT	TOFAC,YMAX,1
12134	1P	01402	J2022	3/24		BTM	TRACE,YLABEL,17
12146	2P	03002	11075	3/22		BT	FSB,YMIN,1
12158	1P	03218	J2058	3/22		BTM	FDV,FIFTY,17
12170	1P	01402	J2002	3/24		BTM	TRACE,DY,17
				3/22	*		
12182	2P	01314	11050	3/22		BT	TOFAC,XMAX,1
12194	2P	03002	11055	3/22		BT	FSB,XMIN,1
12206	1P	03218	J2038	3/22		BTM	FDV,XNO,17
12218	1P	01402	J2012	3/24		BTM	TRACE,DX,17
				3/22	*		
12230	1P	03046	J1982	3/25		BTM	FMP,HALF,17
12242	2P	02336	11055	3/25		BT	FSBR,XMIN,1
12254	1P	01402	J2073	3/25		BTM	TRACE,XXMIN,17
12266	1P	01314	J2002	3/25		BTM	TOFAC,DY,17
12278	1P	03046	J1982	3/22		BTM	FMP,HALF,17
12290	2P	02430	11070	3/22		BT	FAD,YMAX,1
12302	1P	01402	J1992	3/24		BTM	TRACE,YYMAX,17
				3/22	*		
12314	J6	12063	0-000	3/22		TFM	KY,0,08
12326	J6	11920	00-01	3/22		TFM	I,1,09
12338	J6	11927	00-01	3/25	RTN1	TFM	II2,1,09
12350	J0	11972	J1711	3/22	RTN111	TFM	COUNT,OUTPUT,017
12362	KJ	11972	11927	3/22		A	COUNT,II2,01
12374	KJ	11972	11927	3/22		A	COUNT,II2,01
12386	J6	1197K	000-0	3/22		TFM	COUNT,0,0610
12398	J1	11927	000-1	3/22		AM	II2,1,010
12410	KM	12028	11927	3/24		C	NOX,II2,01
12422	M6	12350	01300	3/24		BNN	RTN111,,0
12434	J5	12023	00000	3/22		TDM	IND,0,0
				3/22	*	GRID.....	
12446	J3	12063	-005-	3/22		MM	KY,50,0711
12458	32	00095	00000	3/23		SF	99-KK+1
12470	26	00485	00099	3/22		TF	FAC,99
12482	2P	01652	11080	3/22		BT	FXD,NY,1
12494	12	00485	000-1	3/22		SM	FAC,1,10
12506	2J	00485	11920	3/22		A	FAC,I,1
12518	14	00485	0-000	3/22		CM	FAC,0,8
12530	M7	12662	01200	3/22		BNE	I211,,0
				3/22	*		
12542	J6	11924	00-01	3/25	I222	TFM	JJ,1,09
12554	J0	11972	J1711	3/22	RTN332	TFM	COUNT,OUTPUT,017
12566	KJ	11972	11924	3/22		A	COUNT,JJ,01
12578	KJ	11972	11924	3/22		A	COUNT,JJ,01
12590	J6	1197K	000K0	3/22		TFM	COUNT,20,0610
12602	J1	11924	000-1	3/22		AM	JJ,1,010
12614	KM	12028	11924	3/24		C	NOX,JJ,01
12626	M6	12554	01300	3/24		BNN	RTN332,,0
				3/22	*		
12638	J5	12023	00001	3/22		TDM	IND,1,0
12650	J1	12063	000-1	3/22		AM	KY,1,010
				3/22	*		
12662	J6	11924	00-00	3/22	I211	TFM	JJ,0,09
12674	20	00485	12028	3/22	RTN444	TF	FAC,NOX,1
12686	12	00485	000-1	3/25		SM	FAC,1,10

12698	2L	00485	11924	3/22		M	FAC,JJ,1
12710	32	00095	00000	3/23		SF	99-KK+1
12722	26	00485	00099	3/25		TF	FAC,99
12734	2P	01652	11060	3/22		BT	FXD,NX,1
12746	JO	11932	J1711	3/22		TFM	I2,OUTPUT,017
12758	K1	11932	00485	3/22		A	I2,FAC,0
12770	K1	11932	00485	3/25		A	I2,FAC,0
12782	J1	11932	000-2	3/22		AM	I2,2,010
12794	J6	1193K	000P1	3/22		TFM	I2,71,0610
12806	J1	11924	000-1	3/22		AM	JJ,1,010
12818	KM	1106-	11924	3/24		C	NX,JJ,016
12830	M6	12674	01300	3/24		BNN	RTN444,,0
				3/22	*		
12842	20	00485	11920	3/22		TF	FAC,I,1
12854	33	00483	00000	3/24		CF	FAC-2
12866	32	00481	00000	3/24		SF	FAC-KK+1
12878	17	02112	000-0	3/22		BTM	FLOAT,0,10
12890	1P	01402	J1947	3/22		BTM	TRACE,ZI,17
12902	1P	03002	J2048	3/22		BTM	FSB,ONE,17
12914	1P	03046	J2002	3/22		BTM	FMP,DY,17
12926	1P	02336	J1992	3/22		BTM	FSBR,YYMAX,17
12938	1P	01402	J1957	3/24		BTM	TRACE,UP,17
12950	1P	03002	J2002	3/22		BTM	FSB,DY,17
12962	1P	01402	J1967	3/24		BTM	TRACE,DOWN,17
				3/22	*		
12974	J6	11937	00-01	3/22		TFM	IF,1,09
12986	J3	11937	000J0	3/22	RTN121	MM	IF,FF,010
12998	2J	00099	11065	3/22		A	99,Y,1
13010	27	01314	00099	3/22		BT	TOFAC,99
13022	1P	03002	J1957	3/22		BTM	FSB,UP,17
13034	M4	13246	00483	3/25		BNF	I1121,FAC-2,0
13046	J3	11937	000J0	3/22	I2221	MM	IF,FF,010
13058	2J	00099	11065	3/22		A	99,Y,1
13070	27	01314	00099	3/22		BT	TOFAC,99
13082	1P	03002	J1967	3/22		BTM	FSB,DOWN,17
13094	M4	13114	00483	3/22		BNF	I3331,FAC-2,0
13106	M9	13246	00000	3/22		B	I1121,,0
13114				3/22		DORG	*-3
13114	J3	11937	000J0	3/22	I3331	MM	IF,FF,010
13126	2J	00099	11045	3/22		A	99,X,1
13138	27	01314	00099	3/22		BT	TOFAC,99
13150	1P	03002	J2073	3/25		BTM	FSB,XXMIN,17
13162	1P	03218	J2012	3/22		BTM	FDV,DX,17
13174	17	01772	000-0	3/22		BTM	FIX,0,10
13186	11	00485	000-1	3/22		AM	FAC,1,10
13198	13	00485	-0002	3/22		MM	FAC,2,7
13210	32	00095	00000	3/22		SF	95
13222	1J	00099	J1711	3/22		AM	99,OUTPUT,17
13234	16	0009R	000J4	3/22		TFM	99,14,610
13246	J1	11937	000-1	3/22	I1121	AM	IF,1,010
13258	KM	1108N	11937	3/24		C	N,IF,016
13270	M6	12986	01300	3/24		BNN	RTN121,,0
13282	ML	13318	12023	3/22		BD	I11,IND,01
13294	L9	11687	00900	4/2	I10	WA	OUTPUT-24,00900,0
13306	M9	13366	00000	3/13		B	I12,,0
13318	L9	11687	00901	4/2	I11	WA	OUTPUT-24,00901,0
				3/13	*		ADD YLABEL OUTPUT .....
13330	1P	05028	J1653	4/24		BTM	WATY,FMT-5,17
13342	1P	06066	J2022	4/2		BTM	SWC,YLABEL,17
13354	17	05650	000-0	4/2		BTM	COMPLT,0,10
13366	1P	01314	J2022	3/13	I12	BTM	TOFAC,YLABEL,17

13378	1P	03002	J2002	3/13
13390	1P	01402	J2022	3/13
13402	J1	11920	000-1	3/13
13414	J4	11920	000N1	3/13
13426	M7	12338	01100	3/13
13438	M9	1109L	00000	3/13
00000				3/13

BTM	FSB,DY,17
BTM	TRACE,YLABEL,17
AM	I,1,010
CM	I,51,010
BNP	RTN1,,0
B	START-1,,06
DEND	









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C   AUTOMATIC LINKAGE GENERATOR FOR SPS SUBS. FOR FN II.
C
C   L. HOFFMAN, GUGGENHEIM LABS.
C
*1205
C   SPS - FN II LINKAGE AND CONSTANT AUTOMATIC GENERATOR.....
    DIMENSION VAR(20)
27  READ 1,N
    1  FORMAT(I5)
    READ 2,(VAR(I),I=1,N)
    2  FORMAT(A6)
    4  TYPE 3
    3  FORMAT(22H TYPE SUBROUTINE NAME. )
    ACCEPT 2,SNAME
    IF(SNAME)5,4,5
    5  PUNCH 6
    PRINT 6
    6  FORMAT(5HAUTO ,6X,9HDORG11036)
    PUNCH 7,SNAME
    PRINT 7,SNAME
    7  FORMAT(5HAUTO ,A6,5HDS 5)
    DO 8 I=1,N
    PUNCH 9,VAR(I)
    PRINT 9,VAR(I)
    9  FORMAT(5HAUTO ,A6,5HDS 5)
    8  CONTINUE
    PUNCH 10
    PRINT 10
10  FORMAT(5HAUTO ,6X,5HDS 7)
    PUNCH 30
    PRINT 30
    PUNCH 31
    PRINT 31
    PUNCH 32
    PRINT 32
    PUNCH 45
    PRINT 45
    PUNCH 38
    PRINT 38
    PUNCH 33
    PRINT 33
    PUNCH 34
    PRINT 34
    PUNCH 37
    PRINT 37
    PUNCH 36
    PRINT 36
    PUNCH 35
    PRINT 35
    PUNCH 41
    PRINT 41
    PUNCH 40
    PRINT 40
    PUNCH 39
    PRINT 39
    PUNCH 42
    PRINT 42
    PUNCH 43
    PRINT 43
    PUNCH 44
    PRINT 44

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```

PUNCH 46
PRINT 46
PUNCH 49
PRINT 49
PUNCH 48
PRINT 48
PUNCH 47
PRINT 47
30 FORMAT(5HAUTO ,6HFF ,4HDS ,6H,10 )
31 FORMAT(5HAUTO ,6HKK ,4HDS ,6H,5 )
32 FORMAT(5HAUTO ,6HFAC ,4HDS ,6H,485 )
45 FORMAT(5HAUTO ,6HTOFAC ,4HDS ,6H,1314 )
38 FORMAT(5HAUTO ,6HFMFAC ,4HDS ,6H,1358 )
33 FORMAT(5HAUTO ,6HTRACE ,4HDS ,6H,1402 )
34 FORMAT(5HAUTO ,6HFXD ,4HDS ,6H,1652 )
37 FORMAT(5HAUTO ,6HFIX ,4HDS ,6H,1772 )
36 FORMAT(5HAUTO ,6HFLOAT ,4HDS ,6H,2112 )
35 FORMAT(5HAUTO ,6HRSGN ,4HDS ,6H,2336 )
41 FORMAT(5HAUTO ,6HFSBR ,4HDS ,6H,2336 )
40 FORMAT(5HAUTO ,6HFAD ,4HDS ,6H,2430 )
39 FORMAT(5HAUTO ,6HFSB ,4HDS ,6H,3002 )
42 FORMAT(5HAUTO ,6HFMP ,4HDS ,6H,3046 )
43 FORMAT(5HAUTO ,6HFDV ,4HDS ,6H,3218 )
44 FORMAT(5HAUTO ,6HFDVR ,4HDS ,6H,3434 )
46 FORMAT(5HAUTO ,6HWATY ,4HDS ,6H,4920 )
49 FORMAT(5HAUTO ,6HPRA ,4HDS ,6H,5028 )
48 FORMAT(5HAUTO ,6HCOMPLT ,4HDS ,6H,5650 )
47 FORMAT(5HAUTO ,6HSWC ,4HDS ,6H,6066 )
PUNCH 11
PRINT 11
11 FORMAT(5X,1H*,12X,25HEND OF ARGUMENT ADDRESSES/5X,1H*,14X,14HSTART
111 LINKAGE. )
START=.6263415963
BLANK=0.
DO 12 I=1,N
IF(I-1)13,14,13
14 PUNCH 15,START
PRINT 15,START
15 FORMAT(5HAUTO ,A6,4HAM ,13HSTART-1,5,010)
GO TO 16
13 PUNCH 15,BLANK
PRINT 15,BLANK
16 PUNCH 17,VAR(I)
PRINT 17,VAR(I)
17 FORMAT(5HAUTO ,6X,4HTF ,A6,13H,START-1,0111)
PUNCH 18,VAR(I)
PRINT 18,VAR(I)
18 FORMAT(5HAUTO ,6X,4HBNF ,5H*+36,,A6,3H,01)
PUNCH 19,VAR(I)
PRINT 19,VAR(I)
19 FORMAT(5HAUTO ,6X,4HCF ,A6,3H,,0)
PUNCH 20,VAR(I),VAR(I)
PRINT 20,VAR(I),VAR(I)
20 FORMAT(5HAUTO ,6X,4HTF ,A6,1H,,A6,5H,0111)
12 CONTINUE
M=N/2
M=M*2
IF(M-N)21,22,21
22 PUNCH 23
PRINT 23
23 FORMAT(5HAUTO ,6X,4HAM ,13HSTART-1,1,010)

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```
GO TO 24
21 PUNCH 25
PRINT 25
25 FORMAT(5HAUTO ,6X,4HAM ,13HSTART-1,2,010)
24 PUNCH 26
PRINT 26
26 FORMAT(5HAUTO ,6X,4HB ,9HAROUND,,0/5HAUTO ,6X,4HDORG,3H*-3/
261 5X,1H*,12X,36HSYMBOL TABLE AND CONSTANTS HERE ... )
GO TO 27
END
```

DATA FOR AUTO-LINK.

```
9
X
XMAX
XMIN
NX
Y
YMAX
YMIN
NY
N.
```

## A SURVEY OF THE BEGINNING PROGRAMMING COURSE

Clarence B. Germain  
College of St. Thomas  
February 20, 1964

Last Fall, a questionnaire was sent to the 280 schools which are members of the USERS Group. 175 schools responded. The results are tabulated on the following pages.

1. No allowance has been made for non-respondents. This does bias the results.
2. Since the survey covers only schools having 1620's, the figures for the end of 1964 do not reflect the influence of schools which will acquire their first 1620 during the year.
3. A surprising number of respondents gave inconsistent answers; e.g., they indicated floating-point hardware, but not divide hardware, or they indicated that 35% of their students run their own SPS programs, while they taught SPS only to 20% of their students.
4. Figures for index registers, binary capabilities, and the 1627 plotter may not be indicative since the questionnaire was circulated too soon after announcement of these features.
5. Average enrollment in the beginning programming courses is 170 students per school per year.
6. Many of the Model II 1620's will supplement existing Model I's, not replace them.
7. Relatively few schools indicated any plans to obtain the 1443 printer.
8. The disk units will more than double in popularity during 1964 with 1/3 of all schools having at least one disk unit by the end of the year.
9. While 3% of the schools offered no course involving Fortran, 35% of the students were taught more than one version of Fortran.
10. At the end of 1963, 51% of the schools had the hardware necessary to run Fortran II; by the end of 1964, this figure will rise to 59%.
11. 85% of the students get "hands on" experience in running their own programs on the computer. This percentage is about the same regardless of what programming systems (SPS, GOTRAN, etc.) are taught.
12. Jim Moore's Multi-Trace, 1.4.003, was the most commonly mentioned trace program taught to students. However, 85% of the schools indicated that they used no trace program in their courses.
13. The figures for textbooks are for use in at least one course. Many schools use more than one text in a course. 31% of the schools use only IBM publications as texts. While a wide variety of texts, many unrelated to either Fortran or the 1620, are in use, only four commercial texts and a half-dozen IBM publications are used with any frequency. Of the non-programming type texts, numerical analysis books, particularly Stanton's, were most often mentioned.
14. The textbook percentages in no way indicate sales of books; these figures are quite different from the percentages shown here and were not a part of this study.

RESPONSES OF 175 SCHOOLS TO A SEPTEMBER 1963 QUESTIONNAIRE

Results are given as a percentage of the number of schools replying to the questionnaire. Probable errors do not exceed  $\pm 3\%$  except for items marked with an asterisk (\*) where the probable error is less than  $\pm 8\%$ . Results are given for the end of 1963 and for the end of 1964. Changes for 1964 are only for equipment now on order. Slight discrepancies in the percentages are due to rounding.

1620 Model:	1963	1964	Number of 1620's in the school:
I	98%	89%	One 95%
II	2	11	Two 5
Special Features, Model I			Special Features, Model II (1964)
AFP, Div, IDA, Edit	31	31	Automatic Floating-Point 65*
AFP, Div, IDA	3	3	Index Registers 0*
AFP, Div, Edit	0	0	Binary Capabilities 5*
AFP, Div	1	1	Installations with Printer (1964)
Div, IDA, Edit	31	31	No disk 23*
Div, IDA	14	14	1 disk 15*
Div, Edit	1	1	2 disks 54*
Div	3	3	3 disks 0*
IDA, Edit	1	1	4 disks 8*
IDA	3	3	
Edit	1	1	Type of Courses Offered:
No special features	13	13	Both credit and non-credit 5
Summary:			Non-credit courses only 36
Automatic Floating-Point	34	35	Credit courses only 13
Automatic Divide	82	82	No answer or no courses 47
Indirect Addressing	82	82	Departments which offer courses:
Additional (Edit) Instructions	64	64	Engineering 40
Storage:			Education 1
20K core, no disk	48	38	Mathematics 45
40K core, no disk	21	18	Business 31
60K core, no disk	17	13	Other 40
20K core, disk	5	12	Subjects Taught:
40K core, disk	4	9	Machine Language 32
60K core, disk	5	9	Operation of the Computer 66
Input-Output:			SPS 29
Paper Tape only	4	4	GOTRAN 17
Paper Tape and Cards	10	10	FORTTRAN with FORMAT 47
Cards only	86	86	FORTTRAN II or II-D 33
Magnetic Tape	4	4	FORGO, etc. 35
Paper Tape	13	14	Use of some library trace 13
Cards, 1622-1	83	81	Block Diagramming 63
Cards, 1622-2	13	16	Monitor I 9
Cards, RPQ to read 800 cpm	3	3	
1443 Printer		8	
Disk, one or more	14	31	
1627 Plotter	4	4	
1710	2	3	

Disks:		
No disk	86	68
1 disk	8	20
2 disks	5	11
3 disks	0	0
4 disks	1	1
Hardware necessary to run:		
Fortran II only	37	29
Fortran II and II-D	9	19
Fortran II-D only	5	11

Students are expected to write and run their own programs using:

SPS II	25
GOTRAN	15
FORTRAN with FORMAT	43
FORTRAN Pre-Compiler	28
FORTRAN II	27

Required or recommended texts:

IBM Publications

1620 Reference Manual	74
1710 Reference Manual	4
SPS Reference Manual	49
GOTRAN Reference Manual	22
1620 FORTRAN Reference Manual	61
1620 FORTRAN II Bulletin	38
FORTRAN General Information Manual	23
1620 Program Writing and Testing Bulletin	12
Introduction to IBM Data Processing Systems	15
Programming and Block Diagramming Techniques	12

Commercial Publications

Germain—Programming the IBM 1620	27
Leeson-Dimitry—Basic Programming Concepts and the IBM 1620 Computer	39
Gruenberger-McCracken—Introduction to Electronic Computers	6
McCracken—A Guide to FORTRAN Programming	38
Organick—A FORTRAN Primer	38
Colman-Smallwood—Computer Language	6
Smith-Johnson—FORTRAN Autotester	3

Utah State University

Logan, Utah

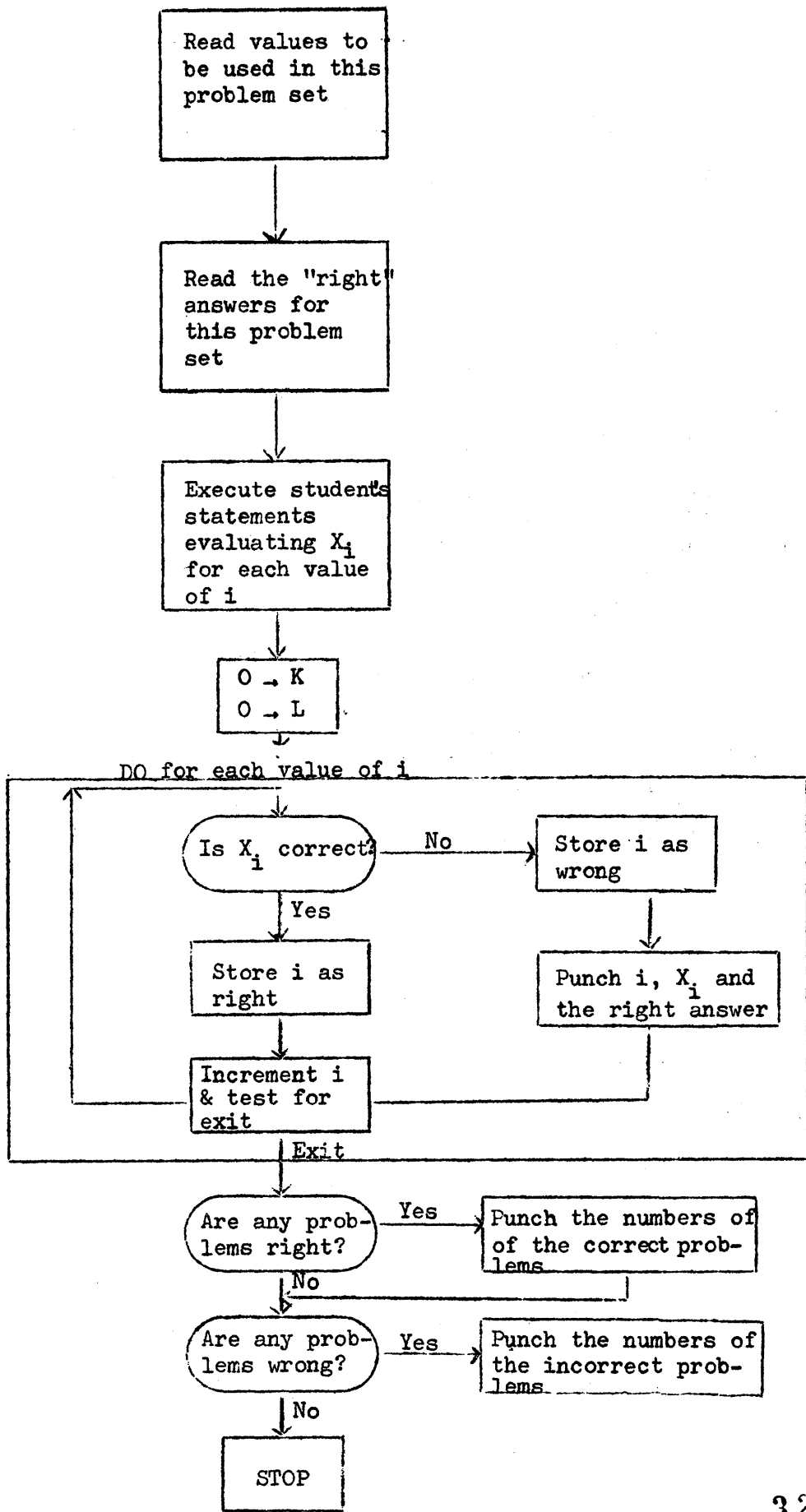
FORTRAN "TEACH" PROBLEMS

by

Wendell L. Pope

These problems are designed to be of assistance in introducing the neophyte to FORTRAN. Problem sets and programs to check them are provided for arithmetic statements, subscripted variables, fixed and floating point variables, functions and control statements, loops and input-output. The problems do not require that a student be able to write a complete program. They provide a means of acquainting him with the characteristics of FORTRAN in easy stages and help to bridge the gap between the introduction to computing and the writing of a complete program. The student's statements are checked for correctness by imbedding them in the appropriate checking program. They are checked for compilation errors by the FORGO processor, and for accuracy by the checking program itself. This is done by comparing the values computed by the student's statements to a predetermined set of "correct" values. For wrong answers, the number of the problem and the value computed are output, for right answers only the number of the problem is output.

FLOW CHART - TEACH Problem Checking Program





```

C TEACH PROBLEM NO 1 PLACE STUDENT HEADER CARD IN FRONT OF THIS CARD
  DIMENSION NWRITE(10),NWRNG(10),RIGHT(10),X(10)
  READ,A,B,C,D,E,F,G,H,(RIGHT(I),I=1,10)
C INSERT STUDENT STATEMENTS BEHIND THIS CARD
  K=0
  L=0
  DO 9950 I=1,10
  IF(X(I)-RIGHT(I)) 9912,9914,9912
9912 L=L+1
  NWRNG(L)=I
  PUNCH 9920,I,X(I),RIGHT(I)
  GO TO 9950
9914 K=K+1
  NWRITE(K)=I
9950 CONTINUE
  IF(K) 9922,9924,9922
9922 PUNCH 9923,(NWRITE(I),I=1,K)
9923 FORMAT(6H RIGHT,10I5)
9924 IF(L) 9925,9926,9925
9925 PUNCH 9927,(NWRNG(I),I=1,L)
9926 STOP
9920 FORMAT(9H PROB NO. I3,4X,9HYOUR ANS=E16.8,4X,10HRIGHT ANS= E16.8)
9927 FORMAT(6H WRONG, 10I5)
  END

```

```

10.341296 10.345599 8.6867569 -40.683394 15.683097 .0034784067
.00034329602 1.1234567
.10664629E+02 0.41901625 0.84447317E+02 0.12039480E+02
.65303219E+09 0.10259388E+01 0.14580430F+01 0.87982136F+01
.46184027E+02 0.11729093E+03

```

C TEACH PROBLEM NO 2 PLACE STUDENT HEADER CARD IN FRONT OF THIS CARD  
 DIMENSION NWRITE( 5),NWRNG( 5),RIGHT( 5),X( 5),B(4),A(4,4)  
 READ,K,L,((A(I,J),J=1,4),I=1,4),(B(I),I=1,4),(RIGHT(I),I=1,5)  
 C INSERT STUDENT STATEMENTS BEHIND THIS CARD

```

K=0
L=0
DO 9919 I=1,5
IF(X(I)-RIGHT(I)) 9912,9914,9912
9912 L=L+1
NWRNG(L)=I
PUNCH 9920,I,X(I),RIGHT(I)
GO TO 9919
9914 K=K+1
NWRITE(K)=I
9919 CONTINUE
IF(K) 9922,9924,9922
9922 PUNCH 9923,(NWRITE(I),I=1,K)
9923 FORMAT(5HRIGHT,10I5)
9924 IF(L) 9925,9926,9925
9925 PUNCH 9927,(NWRNG(I),I=1,L)
9926 STOP
9927 FORMAT(5HWRONG,10I5)
9920 FORMAT(8HPROB NO. 13,4X,9HYOUR ANS= E16.8,4X,10HRIGHT ANS= E16.8)
END
  
```

2	3						
2.3964587	3.6241346	4.1357653	5.3422587	3.3524569	4.3687946	5.3124	
6.0247685	4.1024678	5.3751468	6.0347312	7.3107386	5.3420769	6.0467	
7.3214680	8.3469201	9.3704368	10.437695	-1.3579430	-2.5347962		
5.4673001	41.125807	4913189.2	7701.4270	.86225934			

C TEACH PROBLEM NO 3 PLACE STUDENT HEADER CARD IN FRONT OF THIS CARD  
DIMENSION NWRITE(10),NWRNG(10),RIGHT(10),X(10),B(4)

READ,K,L,(B(I),I=1,4),(RIGHT(I),I=1,5)

C INSERT STUDENT STATEMENTS BEHIND THIS CARD

K=0

L=0

DO 9919 I=1,5

IF(X(I)-RIGHT(I)) 9912,9914,9912

9912 L=L+1

NWRNG(L)=I

PUNCH 9920,I,X(I),RIGHT(I)

GO TO 9919

9914 K=K+1

NWRITE(K)=I

9919 CONTINUE

IF(K) 9922,9924,9922

9922 PUNCH 9923,(NWRITE(I),I=1,K)

9923 FORMAT(5HRIGHT,10I5)

9924 IF(L) 9925,9926,9925

9925 PUNCH 9927,(NWRNG(I),I=1,L)

9926 STOP

9927 FORMAT(5HWRONG,10I5)

9920 FORMAT(8HPROB NO. I3,4X,9HYOUR ANS= E16.8,4X,10HRIGHT ANS= E16.8)

END

2 3

9.3704368 10.437695 -1.3579430 -2.5347962

157.09495 -1.3579430 15.915392 205.01975 11.0

```

C TEACH PROBLEM NO 4 PLACE STUDENT HEADER CARD IN FRONT OF THIS CARD
C TEACH PROBLEM NO 4 PLACE STUDENT HEADER CARD IN FRONT OF THIS C
  DIMENSION NWRITE(5),NWRNG(5),RIGHT (5),X(5)
  READ,A,B,C,D,E,F,G,H,K,(RIGHT(I),I=1,5)
C INSERT STUDENT STATEMENTS BEHIND THIS CARD
  K=0
  L=0
  DO 9919 I=1,5
  IF (X(I)-RIGHT(I))9912,9914,9912
9912 L=L+1
  NWRNG(L)=I
  PUNCH 9920,I,X(I),RIGHT(I)
  GO TO 9919
9914 K=K+1
  NWRITE(K)=I
9919 CONTINUE
  IF(K)9922,9924,9922
9922 PUNCH 9923,(NWRITE(I),I=1,K)
9923 FORMAT(5HRIGHT,10I5)
9924 IF(L)9925,9926,9925
9925 PUNCH 9927,(NWRNG(I),I=1,L)
9926 STOP
9927 FORMAT(5HWRONG,10I5)
9920 FORMAT(8HPROB NO. I3,4X,9HYOUR ANS= E16.8,4X,10HRIGHT ANS= E16.8:
  END
1.2457369 2.3580123 3.8609756 4.7602541 5.3025768 6.2047536
7.0367524 8.3205689 2
1.1161260 1.3001178 .78878076 2.4429843 5.6480088

```

C TEACH PROBLEM NO 5 PLACE STUDENT HEADER CARD IN FRONT OF THIS CARD

DIMENSION NWRITE(5),NWRNG(5),RIGHT(5),X(5)  
READ,A,B,C,D,E,(RIGHT(I),I=1,5)

C INSERT STUDENT STATEMENTS BEHIND THIS CARD

K=0

L=0

DO 9919 I=1,5

IF (X(I)-RIGHT(I))9912,9914,9912

9912 L=L+1

NWRNG(L)=I

PUNCH 9920,I,X(I),RIGHT(I)

GO TO 9919

9914 K=K+1

NWRITE(K)=I

9919 CONTINUE

IF(K)9922,9924,9922

9922 PUNCH 9923,(NWRITE(I),I=1,K)

9923 FORMAT(5HRIGHT,10I5)

9924 IF(L)9925,9926,9925

9925 PUNCH 9927,(NWRNG(I),I=1,L)

9926 STOP

9927 FORMAT(5HWRONG,10I5)

9920 FORMAT(8HPROB NO. I3,4X,9HYOUR ANS= F16.8,4X,10HRIGHT ANS= E16.8

END

2.3964587 3.6241346 4.1357653 5.3422587 3.3524569 112761.0 23922.7  
11018.845 7478.2325 6160.7034

```

C TEACH PROBLEM NO 6 PLACE STUDENT HEADER CARD IN FRONT OF THIS CARD
  DIMENSION NWRITE(5),NWRNG(5),RIGHT(5),X(5)
  READ,(RIGHT(I),I=1,5)
C INSERT STUDENT STATEMENTS BEHIND THIS CARD
  K=0
  L=0
  DO 9919 I=1,5
    IF (X(I)-RIGHT(I))9912,9914,9912
9912 L=L+1
    NWRNG(L)=I
    PUNCH 9920,I,X(I),RIGHT(I)
    GO TO 9919
9914 K=K+1
    NWRITE(K)=I
9919 CONTINUE
    IF(K)9922,9924,9922
9922 PUNCH 9923,(NWRITE(I),I=1,K)
9923 FORMAT(5HRIGHT,10I5)
9924 IF(L)9925,9926,9925
9925 PUNCH 9927,(NWRNG(I),I=1,L)
9926 STOP
9927 FORMAT(5HWRONG,10I5)
9920 FORMAT(8HPROB NO. I3,4X,9HYOUR ANS= F16.8,4X,10HRIGHT ANS= E16.8
  END
    ,97569023E+04  0.11699576E+04  0.15283722E+05
    .34042183E+03  0.23871681E+04
    14.369025  15.6753869  1.0367521  6.9851203
    8.4357205
    3.9857423
    9.8530247
    8.5432586

```

TEACH PROBLEM SET 1  
Arithmetic Expressions

Assume that values of A, B, C, D, E, F, G and H are in storage. Write and keypunch correct FORTRAN statements to evaluate each of the following expressions.

$$1. X(1) = A + \frac{-B}{C + D}$$

$$2. X(2) = \frac{A + B}{C - D}$$

$$3. X(3) = AB + C \frac{D}{E-F} - G$$

$$4. X(4) = A + \frac{B}{C + \frac{D}{E + \frac{F}{G-H}}}$$

$$5. X(5) = A + B^C - \frac{D}{E}$$

$$6. X(6) = \left( A + B \frac{C-D}{E} \right)^F$$

$$7. X(7) = A^{(B^2+C)^D} + \frac{\pi}{\left( E - \frac{F}{G} \right)^H} \quad (\pi = 3.141592653\dots)$$

$$8. X(8) = \frac{A/B}{C} + \frac{A}{B/C}$$

9. If  $E > B > A > H > C > F > G > 0$ , compute

$$X(9) = A + B + C + E + F + G + H$$

to obtain the most accuracy.

$$10. X(10) = A + B^2 + \frac{3}{4+D}$$

TEACH PROBLEM SET 2  
Subscripted Variables

Assume the arrays  $A = \begin{pmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \\ a_{41} & a_{42} & a_{43} & a_{44} \end{pmatrix}$  and  $B = \begin{pmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \end{pmatrix}$

are in storage. Write and keypunch correct FORTRAN statements to evaluate each of the following expressions.

1.  $X(1) = \frac{a_{12}}{a_{11}} + \frac{a_{13}}{a_{11}} + \frac{a_{14}}{a_{11}}$

2.  $X(2) = \sum_{j=1}^4 a_{1j} b_j$

3.  $X(3) = \sum_{j=1}^4 a_{2j} b_j$

4.  $X(4) = a_{31}X_1 + a_{32}X_1^2 + a_{33}X_1^3 + a_{34}X_1^4$

5.  $X(5) = \frac{\{b_1\}}{b_2}$  where  $\{b_1\}$  denotes the integer portion of  $b_1$ .

TEACH PROBLEM SET 3

Fixed point, Floating point, and subscripted variables

Assume that values of  $k$  and  $L$  and  $B = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{pmatrix}$  are in storage.

Write and keypunch correct FORTRAN statements to evaluate each of the following expressions.

1.  $X_1 = \sum_{i=1}^4 b_i^i$

2.  $X_2 = b_{k^3-5}$

3.  $X_3 = \sum_{i=1}^4 b_{ik-i}$

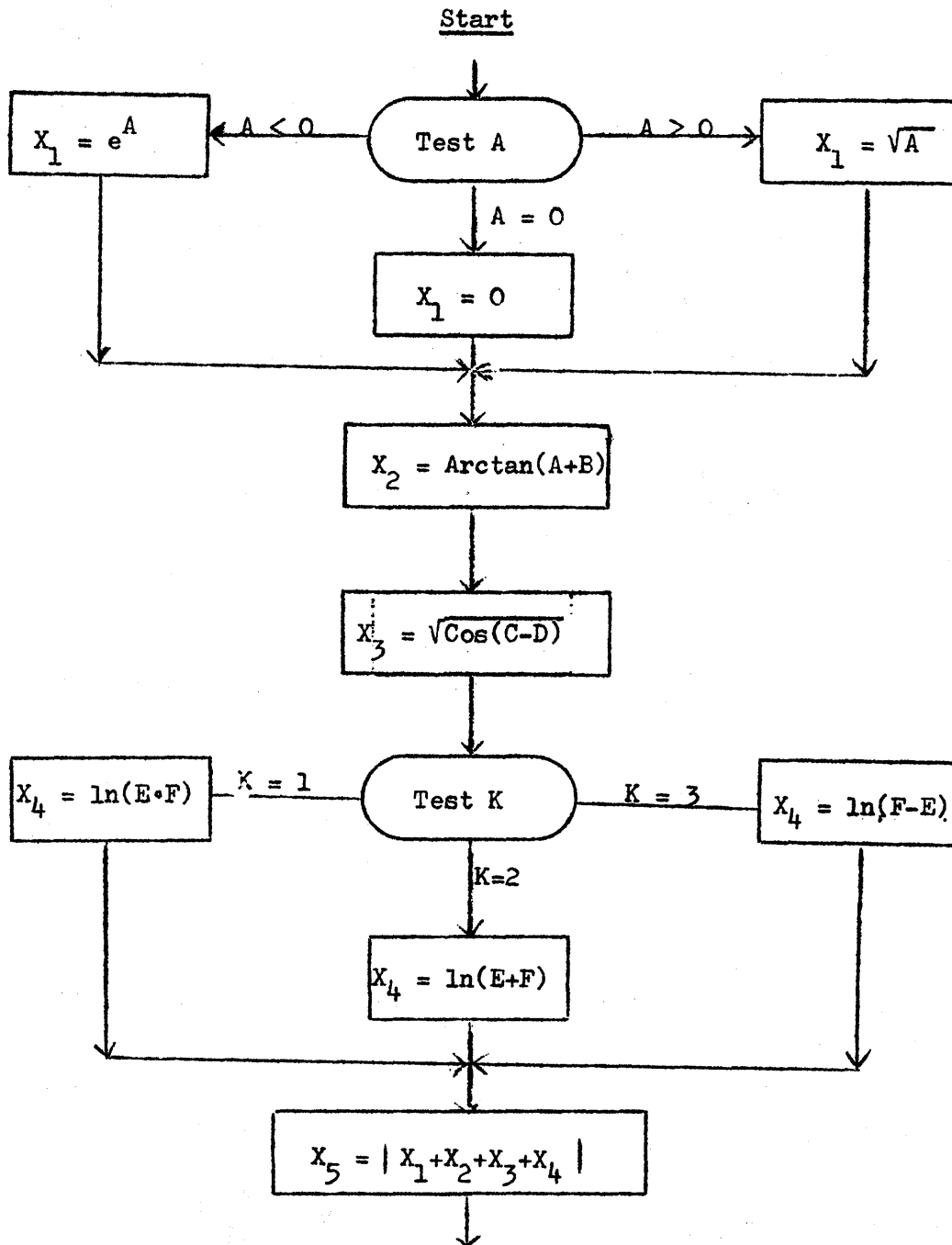
4.  $X_4 = \sum_{i=1}^4 b_i^k$

5.  $X_5 = \frac{L}{k} + k^L + k^3 - 2$



TEACH PROBLEM SET 4  
 Functions and Control Statements

Write (and keypunch) statements to evaluate  $X_1$ ,  $X_2$ ,  $X_4$  and  $X_5$  according to the instructions in the flow chart below. Assume A, B, C, D, E, F and K to be defined.



TEACH PROBLEM SET 5

Loops

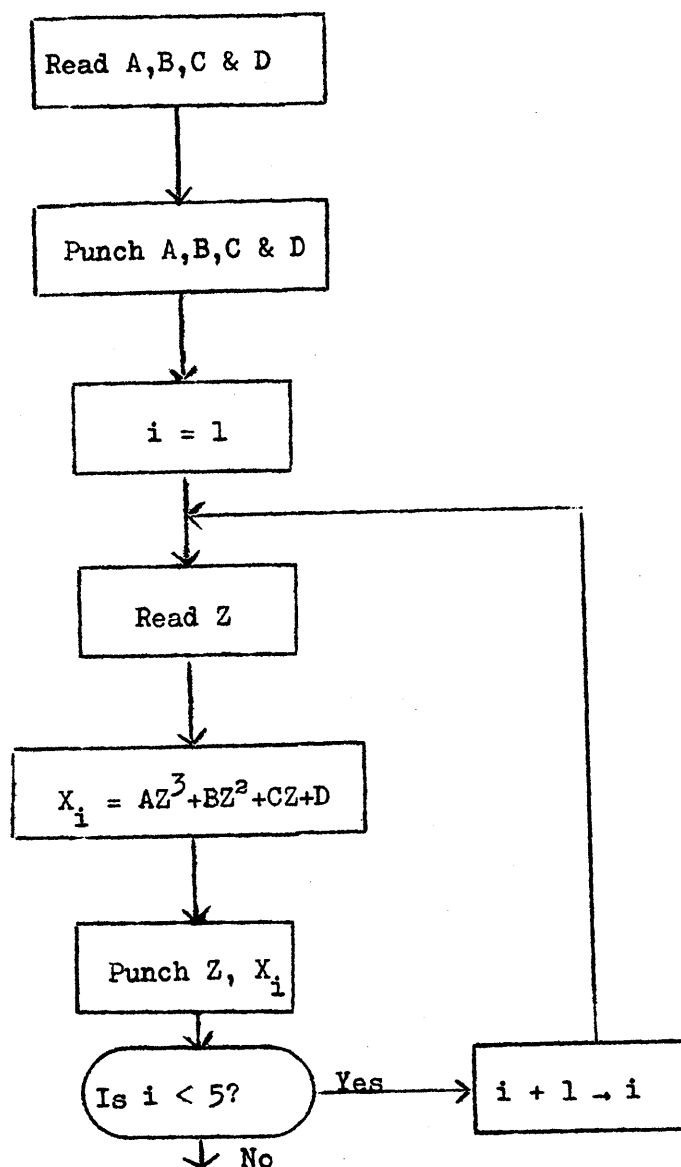
Write (and keypunch) statements to evaluate  $X_1, X_2, X_3, X_4$  and  $X_5$  in the exercises below.

1.  $X_1 = 3^2 + 6^2 + 9^2 + 12^2 + \dots + 99^2.$

2.  $X_i = \frac{A + B^C}{D} \sqrt{E \cdot X_{i-1}}, \quad i = 2, 3, 4, 5.$

TEACH PROBLEM SET 6

Input - Output (without formats)



LOAD-AND-GO SPS WITH MONITOR CONTROL

Kenneth M. Lochner and Glenn R. Ingram

## MSC ASSEMBLY SYSTEM:

### A LOAD-AND-GO SPS WITH MONITOR CONTROL

*Kenneth M. Lochner and Glenn R. Ingram*

#### I.) INTRODUCTION

This paper will discuss a monitored, load-and-go type assembler developed at Montana State College and the conditions that prompted its development. It is a report on work done by Ken Lochner, formerly of the MSC Computing Center, and soon to assume duties as chief programmer at the Dartmouth Computation Center, in the actual writing of the processor.

To suggest some of the background reasons for this processor, it is well to admit that I am a relative newcomer to 1620 ranks in completing my second academic year with a 1620, after leaving a 709 installation. Anyone who follows this path finds himself wondering why in the world he did, and then develops a feeling somewhat akin to the fellow who had a job with a circus. This particular job consisted of following behind the animals during a parade, and cleaning the street with his little shovel. After an especially trying day, he complained so bitterly that his wife asked, "If it's so bad, why don't you quit and get another job?" The man replied, "What! And get out of show business?"

If the analogy isn't exact, it may be suggestive that some things could be cleaner in the 1620 tent.

#### II.) BACKGROUND AND MOTIVATION

To indicate some of the motivation for the work discussed here, Montana State College has, since 1958, offered courses in symbolic programming. This programming is not viewed as an end in itself, but as a basis for the presentation of somewhat more sophisticated topics. This was a necessity while the computer was a 650, and the practice has continued with the Model I 1620, and more recently, the Model II. The clear advantages of FORTRAN have caused a change in the philosophy of the introductory course and in the development of library routines.

The ease of teaching FORTRAN, plus its relative machine independence and the ability of the casual user to obtain answers to real problems rather quickly, make it the logical choice for an introductory language. The newcomer is somewhat surprised to discover that FORTRAN will also produce a faster object program than SPS--despite the standard folklore--but is inclined to believe that a symbolic programming system still has a place in the academic world.

We feel that a symbolic system should be taught to students who have mastered the compiler language, and are interested in more depth in the computer. However, even the most innocent SPS devotee must admit that the pre-disk system was, at best, miserable from a teaching standpoint. Without entering into a discussion of whether SPS is a model of the mold in which symbolic systems should be cast, suffice it to say that there are rather obvious drawbacks to its use by a class of students.

Having mentioned the pre-disk system, let it be stressed that the system to be discussed, despite the "monitor" in the title, does not require disks, and has no relation to the disk monitor. Our 1620 has 60K storage and the standard features of the Model II, but the system will operate on a 40K Model I with automatic divide, indirect addressing and card input-output.

### III.) CRITERIA FOR A GOOD ASSEMBLY SYSTEM FROM A TEACHING STANDPOINT

Returning to the motivation for this system, let it be noted that anyone who has taught a symbolic system to beginning programmers is aware that syntax and logical errors abound in the programs they produce. One can visualize the standard scene in a 1620 installation: a group of students loading the assembler, loading and unloading the punch hopper, entering the object deck, watching the typewriter anxiously, and then staring in increasing bewilderment at a machine which has halted, cleared or is in an infinite loop.

This scene, repeated many times, has its effect on the patience and morale of both students and teacher.

The elimination of this scene provides a good starting point in considering what is desirable in a symbolic system. We would submit that the essential points are:

- 1.) The ability to process a number of programs in a short time, which requires
  - a.) Load and Go operation
  - b.) Batch assembly
- and 2.) Good diagnostics
  - a.) During assembly
  - b.) During program execution.

These considerations suggest that the system must always be in core, and provide the third condition:

- 3.) A "student-proof" processor; i.e., one that virtually cannot be erased.

### IV.) FEATURES OF THE MSC ASSEMBLY SYSTEM

To show how these considerations dictated the type of processor developed, and to reinforce the assertion that it was superior to the regular SPS for teaching (and debugging), we list some features of this system beside those of the standard one.

<u>Feature</u>	<u>MSC</u>	<u>Regular SPS</u>
Language	Modified SPS-09	SPS-09 or -20
Operation	Load-and-go, batching	Two passes of source plus load object program
Assembly Diagnostics	Cards	Typewriter
Basic Floating Point Operations	Built-in (rewritten)	Additional deck to load
Execution Error Diagnostics	22 detected	None
Processor	Always in core	In and out

This table gives the essential comparison of pertinent points. To continue, under the MSC system, the program can be traced during execution (at the

operator's option, by a console switch), execution can be halted by turning on a console switch, and a maximum number of instructions that a program can execute may be preset: if the program exceeds this number, the monitor will so indicate, and proceed to the next job. Hence, the only output will be the error messages, results of a trace, or answers computed by the program. Thus, the time consuming card handling of processor and object deck is eliminated, as well as the annoyance of punched output after errors have been detected. Also, the MSC processor has been made as nearly "student-proof" as possible. Once in two quarters the system was wiped out of core in a rare situation associated with a divide command, and we feel that our students are as inventive as anyone's.

#### V.) OPERATION OF THE SYSTEM

The operation of this system can be divided into three phases: a rather trivial "Job Card" search, the assembly of the source program, and the execution of the program under monitor control.

Each of the stacked programs must be preceded by a "Job Card," typically containing the student's name and program description. This card is identified by an asterisk in column 80, and the processor reads cards until such a card has been found. The job card is reproduced and assembly of the source program begins. The purpose of this is, of course, to allow this title card to be differentiated from data cards which may be stacked behind the previous program.

The assembly of the source program proceeds as in the SPS-09 assembler, except that

- 1.) During the equivalent of pass 1, the source statements are stored in core,
- 2.) During the equivalent of pass 2, the source statements are assembled,
- 3.) Diagnostic errors are punched,
- and 4.) If syntax errors were discovered during assembly, an "End of Job" card is punched, and the processor returns to the job card search phase.

If no syntax errors were detected during assembly, monitor control of program execution is initiated.

The primary purpose of the monitor part of the processor is to provide an indication of execution errors while protecting the processor. After a program has been assembled, the monitor successively examines each instruction for possible execution errors, finds equivalent direct addresses, and if there are no errors, executes the instruction. The examination precludes executing an instruction that would, for example, carry out arithmetic on a field containing a record mark, operate on a field or record that hadn't been flagged or marked, etc.

This process continues until

- 1.) Program execution is complete,
- 2.) An error is detected,
- 3.) The number of instructions executed exceeds the number allowed,
- or 4.) The operator terminates program execution by a console switch.

Any one of these four conditions causes an "End of Job" card to be punched, and the processor returns to the job card search phase. Any of the last three conditions causes an error message to be punched before the "End of Job" card.

## VI.) SUMMARY OF CLASS USE

In evaluating the effectiveness of such a system, one is tempted to compare the results of student use with those of students using another system on a different computer; e.g., SOAP on the 650 or FAP, say, on the 709 or 90. However, such a comparison is probably more a reflection of the particular computer than of the programming system, because of the differences in size and complexity of the command structure, and the storage available for a processor.

Perhaps it is more meaningful to indicate that a class of 25 students was able to make effective use of this system to learn symbolic programming for the 1620. Each student was able to have several trials with his program in a two hour period. The programs written varied from a beginning one that simply read cards, rearranged fields, and punched cards, to rather complex ones. Other simple programs involved the various standard arithmetic operations and the floating point routines, and finding roots of polynomials. Other programs included random number generators, subsequent tests for randomness, Monte Carlo evaluation of an integral, square root subroutines, simulation of index registers, and some basic "building block" routines which will be used in the simulation of a symbolic system for a fixed word length computer.

## VII.) CONCLUSION

Summing up, it may be gathered that we have been pleased with this processor. It has been useful in debugging programs other than those written by students. Its effectiveness has been purchased at the cost of the time used in the detailed scrutiny of the monitor. Balancing that cost against the time that would be used in the extensive card handling of the regular systems, the gain in convenience and the ability to handle many programs in one period, has made this processor a real contribution to our classes.

With the oncoming surge of disk storage, our interest in this system may sound comparable to the enthusiasm of the folks who invested in buggy whips just before the car was invented. If so, someone should have developed such a system three years earlier.

We feel this may have a place, even with disks, so it has been submitted to the program library, in four forms. Certainly it has been well-nigh foolproof and indestructible in core, thus providing the rapid on-off form desired. It may be obsolete, but we are too old to be innocent, and we have learned to be skeptical of Greeks bearing gifts and manufacturers bearing software.

EXAMPLES OF 1620 USE IN COLLEGE ADMINISTRATION

Noel T. Smith



**INDIANA STATE COLLEGE**

**TEST SCORE ANALYSIS**

**Wayne E. Hoover  
Computer Center  
Indiana State College  
Terre Haute, Indiana  
May, 1964**

DECK KEY

Deck 1	Source Program
Deck 2	Object Program
Deck 3	Sample Data
Deck 4	Sample Output (sorted alphabetically)

Indiana State College

Test Score Analysis

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## ABSTRACT

Title: Indiana State College Test Score Analysis

Author: Wayne E. Hoover  
Computer Center  
Indiana State College  
Terre Haute, Indiana

Purpose/Description: Suppose a student tells you that he received a score of 36 on a particular examination. This alone does not indicate his achievement in relation to the rest of the class. Before one can interpret such a number, more information is needed. It becomes evident that the degree to which the results of measurement are useful is proportional to the accuracy and thoroughness with which these results are analyzed.

In order to effectively evaluate the attainment of his students accurately and fairly, a professor must have adequate measurement techniques at his disposal. He must know how to use them properly, and how to interpret results obtained by their use. However, an unfortunate characteristic of evaluative techniques is that they are very time consuming.

It is the purpose of this project to write a program for the IBM 1620 Computer which will relieve the professors and their assistants of the burden of manually calculating from a given set of a maximum of 500 test scores, the arithmetic mean, standard deviation, alpha-three, alpha-four, reliability of the test, the rank, percentile rank, z-score, and t-score for each student.

### EQUATIONS USED

Alpha-three:

$$\text{Alpha-3} = \frac{1}{n\sigma^3} \sum_{i=1}^n (X_i - X)^3$$

Alpha-four:

$$\text{Alpha-4} = \frac{1}{n\sigma^4} \sum_{i=1}^n (X_i - X)^4$$

Arithmetic Mean:

$$X = \frac{1}{n} \sum_{i=1}^n X_i$$

Percentile Rank:

$$\text{PR} = \frac{100 (n - \text{rank})}{n}$$

Reliability of Test:

$$R_t = \frac{I^2 - X(I - X)}{\sigma^2 (I - 1)}$$

Standard Deviation:

$$= \sqrt{\frac{1}{N} \sum_{i=1}^n (X_i - X)^2}$$

T-Score (Sigma Score):

$$\text{T-Score} = 50 + 10 (\text{Z score})$$

Z-Score (Standard Score):

$$\text{Z-Score} = \frac{X_i - X}{\sigma}$$

Legend:

I = Total Possible Points

n = Number taking the test

N = n if  $n \geq 30$

n-1 if  $n < 30$

$R_t$  = Reliability of the test

$\sigma$  = Standard Deviation

X = Arithmetic Mean

$X_i$  = Individual Score

REFERENCES

Alpha-three:	p. 102 <sup>1</sup>
Alpha-four:	p. 104 <sup>1</sup>
Arithmetic Mean:	p. 45 <sup>1</sup>
Percentile Rank:	p. 37 <sup>2</sup>
Reliability of Test:	p. 152 <sup>2</sup>
Standard Deviation	p. 85 <sup>1</sup>
T-Score	p. 44 <sup>2</sup>
Z-Score	p. 44 <sup>2</sup>

<sup>1</sup> John E. Freund, Modern Elementary Statistics (Englewood Cliffs, 1963), p. (indicated above).

<sup>2</sup> Victor H. Noll, Introduction to Educational Measurement (Boston, 1957), p. (indicated above).

## COMMENTS

Alpha-three: A widely used measure of skewness and symmetry is alpha-three. When alpha-three is greater than zero, the tail of the distribution is skewed to the right, and the distribution is said to be positively skewed; it is negatively skewed if the tail is at the left and alpha-three is less than zero. The distribution is perfectly symmetrical when alpha-three is exactly equal to zero.

Alpha-four: Peakedness or kurtosis is commonly described by alpha-four. When alpha-four is exactly equal to three, the distribution is said to have the familiar bell shape of the normal distribution. When alpha-four exceeds three, the distribution is very peaked and has a relatively wide tail; the distribution is flat in the middle and has a relatively thin tail when the value of alpha-four is less than three.

Arithmetic Mean: The arithmetic mean is commonly known as the average. It is the sum of all the scores divided by the number of persons taking the test.

Percentile Rank: The percentile rank is a derived score stated in terms of the percentage of examinees in a specified group who fall below a given score point. In other words, it describes a person's relative standing within a particular group. Thus, a percentile rank of 80 means that a person's score was equal to or higher than the scores made by 80 percent of the people in a specified group.

Rank: The student with the highest score receives a rank of 1; the next highest score receives a rank of 2, and so on. The student with the lowest score thus receives the highest rank. In case of ties in rank, we assign to each of the tied observations the mean of the ranks which they jointly occupy. Hence, if two scores are tied for ranks 4 and 5, we assign each rank  $4 \frac{1}{2}$ ; if three observations are tied for ranks 10, 11, and 12, we assign each rank 11.

Reliability of the Test: This is a measure of the consistency with which a test measures what it is intended to measure; it enables one to determine how much reliance he can place on the scores yielded by the test. This particular formula is accurate enough for classroom tests and other ordinary situations; however, it usually gives an underestimate of the true reliability.

Standard Deviation: In most distributions, approximately 34 per cent of the scores lie between the mean or average and a point or score that is one standard deviation away from the mean (in either direction). In other words, approximately two-thirds of the cases or scores will fall within one standard deviation of the mean.



**T-Score (Sigma Score):** This is a standard score which has a mean of 50.0 and a standard deviation of 10.0. It is obtained by adding 50 to 10 times the z-score. For example, the z-score -1.4 and 1.4 correspond to the t-scores 36 and 64, respectively. It has been said that t-scores are the best method that has been devised for direct comparison of individual test results.

**Z-Score (Standard Score):** The basic standard score is the z-score. It tells in simple terms the difference or distance between a stated group's mean and any specified raw-score value. The mean is 0.0 and the standard deviation is 1.0.

For example, suppose a student had a score of 49 and he is to be compared with his class. Given that the mean and standard deviation are 40.0 and 6.0, respectively, we find his z-score is  $(49 - 40)/6 = 1.5$ , which corresponds to a t-score of  $50 + 10(1.5) = 65$ . In other words, this student's score is 1.5 standard deviations above the mean. This may also be interpreted to mean that he scored better than about 93 per cent of the group.

## PREPARATION OF DATA

**Data Cards:** Each IBM data input card (except the header card) must be prepared in the following manner:

<u>Columns</u>	<u>Field</u>
1-5	Student Number
7-27	Name (see Exceptions, below)
28-33	Score (floating-point with decimal in column 32; may have one place to right of decimal in column 33)
40	0
46-47	0.
55-56	0.
64-65	0.
72-73	0.

**Header Card:** The first card of the data must contain in columns 28-33, the total possible points. Columns 40-73 are punched as stated above.

**Sorting:** The data cards are then sorted in descending order, according to columns 28-33. After sorting is completed, the header card is then placed on top of the deck; it is the first data card.

**Limitations:** The program will handle a maximum of 500 individual test scores.

**Exceptions:** The program is so designed that columns 1-27 of the data input cards may be left blank without affecting the rest of the program.

## OUTPUT

**Printed Output:** The arithmetic mean, standard deviation, alpha-three, alpha-four, and the reliability of the test are printed out first (in that order.) This information is printed regardless of the switch settings.

The remaining information may or may not be printed out, depending on the switch settings.

The heading, student number and name, score, number, rank, percentile rank, z-score, and t-score is printed out next.

Finally, each line of typed output contains the following information in the indicated order: student number, name, score, a "counter" number, rank, percentile rank, z-score, and t-score. Students receiving the same score are grouped together; the typewriter double spaces between groups.

When the last card has been processed, the word STOP is printed.

**Punched Output:** The first seven cards contain the arithmetic mean, standard deviation, alpha-three, alpha-four, the reliability of the test, and two blank cards. This information is punched in columns 1-35.

**The Heading:** Name, Scrc, No, Rak, Ptrk, Z-Scrc, and T-Scrc is punched on the eighth card.

The remaining cards contain the following information

<u>Columns</u>	<u>Field</u>
1-5	Student Number
7-27	Name
28-33	Score
38-40	A "counter" number
44-49	Rank
54-59	Percentile Rank
63-68	Z-Score
71-76	T-Score

**Sorting Punched Output:** The punched output cards may be sorted alphabetically by sorting in ascending order according to columns 1-5. (This assumes that the student numbers are assigned in sequence when the last names are in alphabetical order). The first eight cards of the output will be rejected by the sorter because columns 1-5 are blank; however, save these cards to process through the Alphabetic Interpreter.

Alphabetic Interpreting of Punched Output: The information contained in the punched output cards may be printed at the top of each card by means of the IBM 548 Alphabetic Interpreter.

The alphameric interpreter board should be wired as follows:

<u>Wire Read Brushes</u>	<u>to</u>	<u>The Print Entries</u>
7-35		1-29
38-40		30-32
44-49		34-32
54-59		41-46
63-68		48-53
71-76		55-60

MACHINE REQUIREMENTS

Equipment Specifications: IBM 1620, 1622

Storage Requirements: 20K

Source Language: Afit FORTRAN

Special Features: None

## PROGRAM EXECUTION

Step-by-Step Procedure: Assume that the object deck is compiled and at hand. Then ---

1. Clear 1622 Card Read Punch.
2. Load Afit Fortran loader, object deck, and the Afit Fortran Subroutines in the 1622.
3. Check Switches.

PARITY Check Switch	STOP
I/O Check Switch	STOP
OVERFLOW Check Switch	PROGRAM

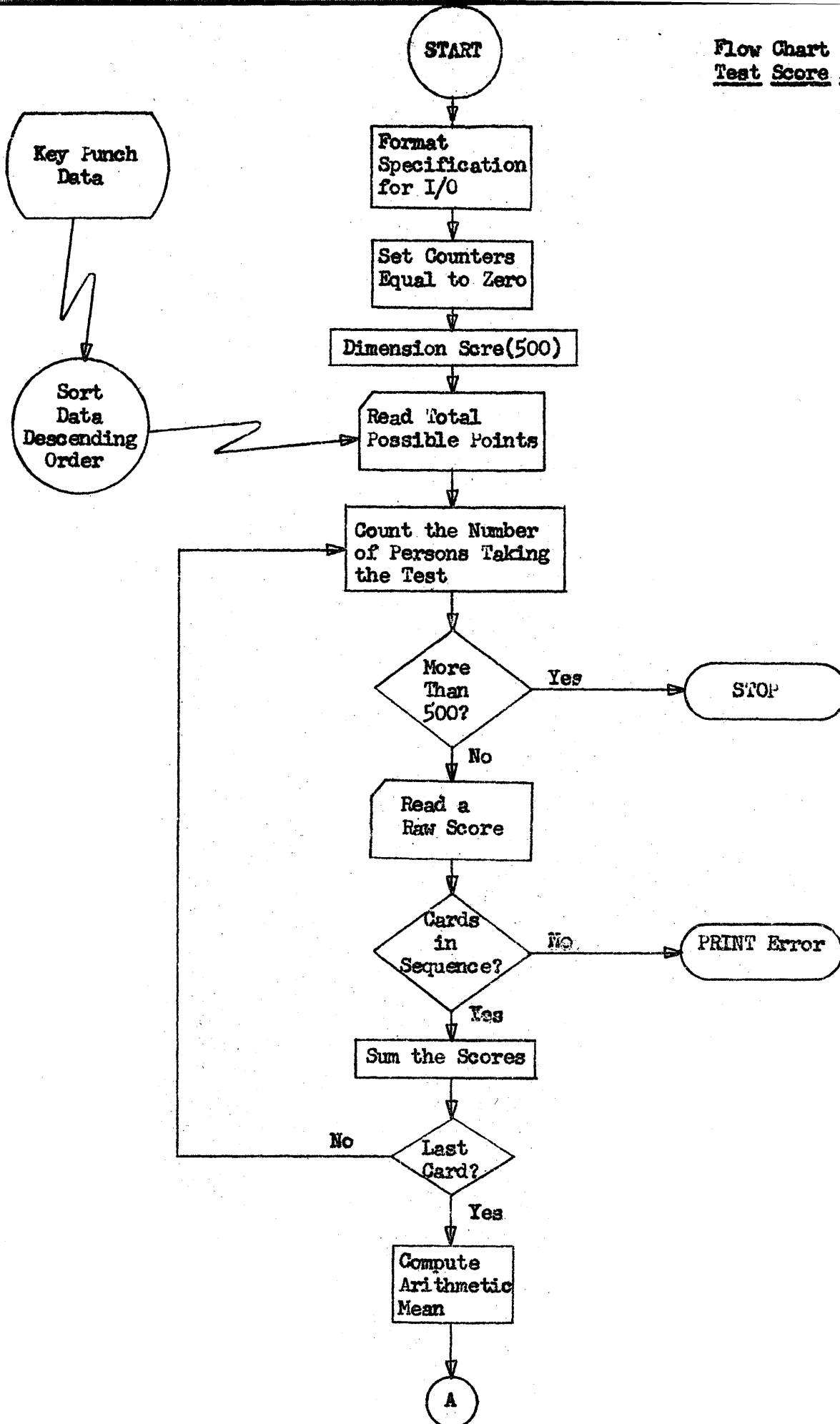
4. Press RESET and Reader LOAD.
5. Set Program Switches for desired option.  
  
Switch 2 ON for PUNCHED output.  
  
Switch 2 and 3 OFF for PRINTED output.  
  
Switch 2 and 3 ON for PRINTED and PUNCHED output.  
  
Switch 1 and 4 OFF.
6. Press START
7. Place data cards in the 1622 and press READER START.
8. Press PUNCH START if necessary.
9. Typewriter prints and/or 1622 punches.
10. Repeat Step 7. (i.e. two passes are required).
11. Typewriter prints and/or 1622 punches. When the last card has been processed, the word STOP is typed.

**Expected Stops:** If the data is not sorted in descending order, a "card out of sequence" message will be printed out and the program will branch to the STOP command. The restart procedure is then necessary.

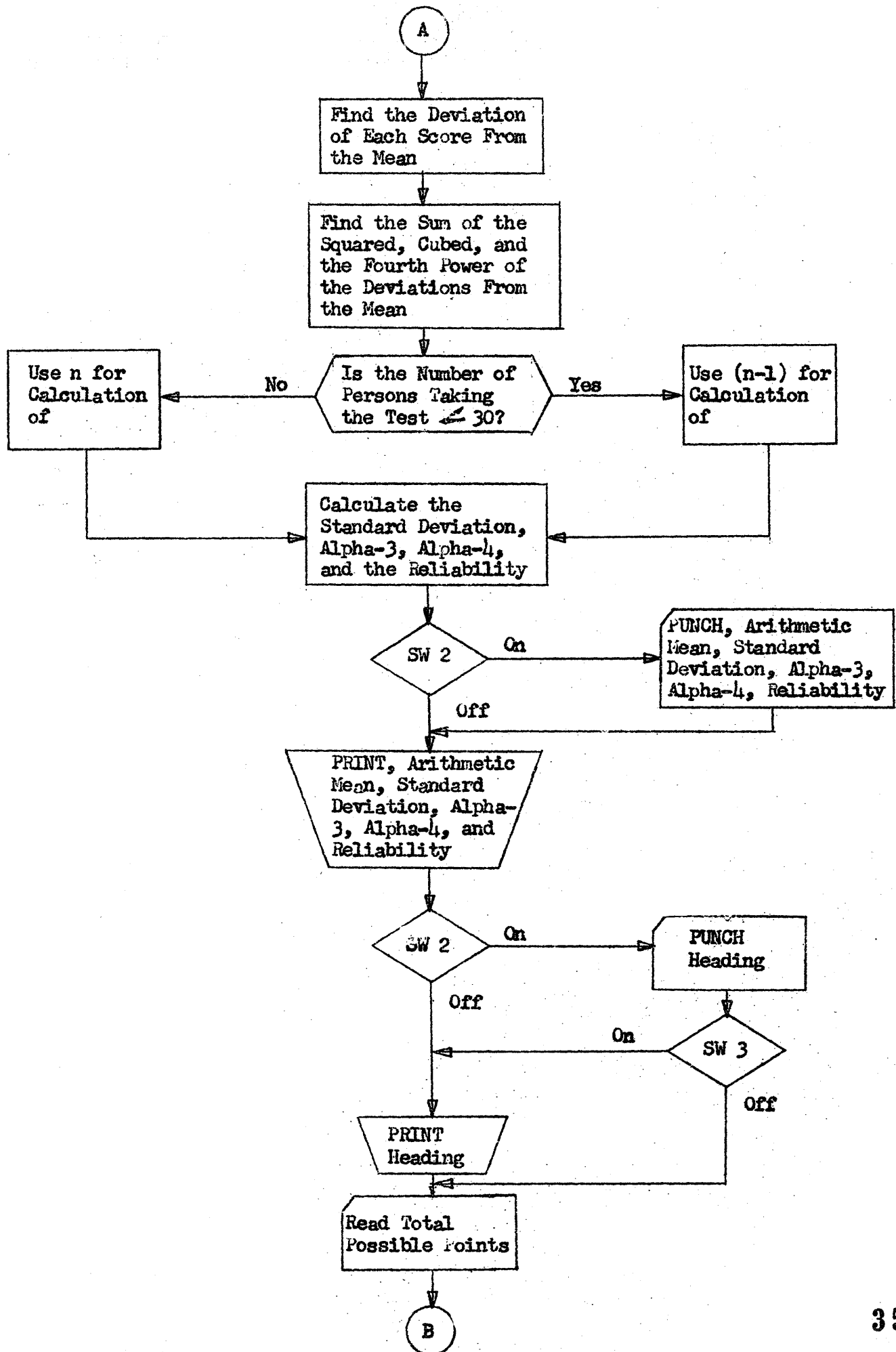
The computer will automatically stop if more than 500 scores are read in.

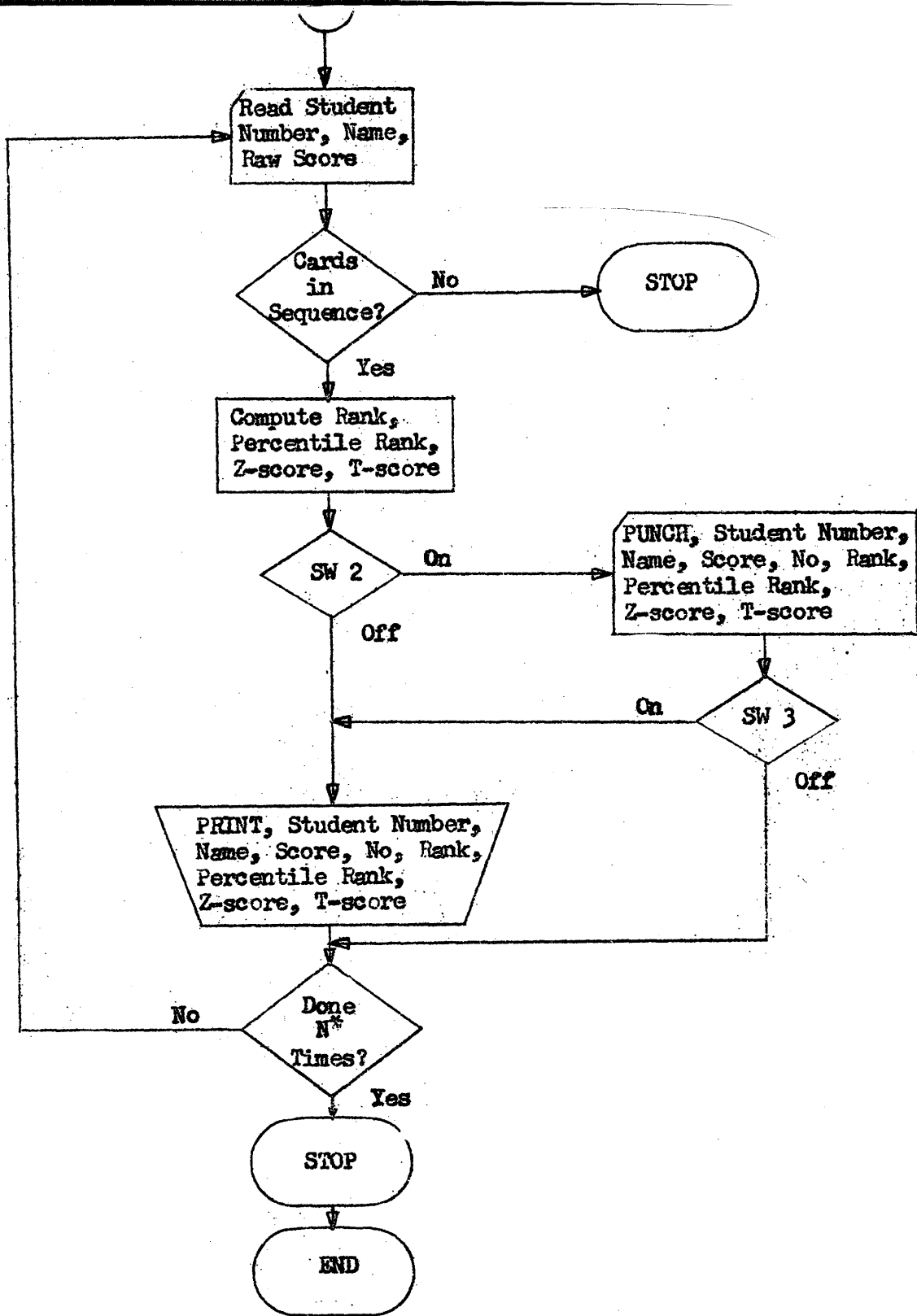
**Restart Procedure:** Press INSTANT STOP, RESET, INSERT, RELEASE, START. Repeat the step-by-step procedure for program execution.

Flow Chart for  
Test Score Analysis









\*Where N equals the number of persons taking the test.

```

08000 C TEST SCORE ANALYSIS WAYNE E HOOVER
08000 C COMPUTER CENTER INDIANA STATE COLLEGE
08000 C TERRE HAUTE, INDIANA MAY, 1964
08000 100 FORMAT(27X,F6.1,17,F9.2,F10.3,F9.3,F8.3)
08108 390 FORMAT(6X4HNAME,19X4HSCORE,5X2HNO,5X3HRNK,5X4HPTRK,6X12HZSCORE TSCORE)
08360 442 FORMAT(6X19HSTUDENT NO AND NAME,3X,5HSCORE,5X,2HNO,5X,4HRANK)
08532 443 FORMAT(11H PCNT RANK,8H Z-SCORE,8H T-SCORE)
08524 776 FORMAT(/)
08646 778 FORMAT(//)
08674 889 FORMAT(6X18HARITHMETIC MEAN IS F11.3)
08758 900 FORMAT(6X22HSTANDARD DEVIATION IS F7.3,//6X11HALPHA 3 IS F18.3)
08914 901 FORMAT(6X11HALPHA 4 IS F18.3,//6X23HRELIABILITY OF TEST IS F6.3)
09072 999 FORMAT(//20HCARD OUT OF SEQUENCE)
09146 SSCORE=0.0
09170 OLD=1000.0
09194 SAFA3=0.0
09218 SAFA4=0.0
09242 SDEV=0.0
09266 N=0
09290 DIMENSION SCORE(500)
09290 READ 100,TPTS,Z1,Z2,Z3,Z4,Z5
09374 13 N=N+1
09410 IF (N-500)14,14,99
09478 14 READ 100,SCORE(N),Z1,Z2,Z3,Z4,Z5
09586 IF (OLD-SCORE(N))16,18,18
09674 16 PRINT 999
09698 GO TO 99
09706 18 SSCORE=SSCORE+SCORE(N)
09766 OLD=SCORE(N)
09814 IF (SENSE SWITCH 9)21,13
09834 21 TNMBR=N
09858 AMEAN=SSCORE/TNMBR
09894 NMBR=N
09918 DO 30 N=1,NMBR
09930 VAR=SCORE(N)-AMEAN
09990 SDEV=SDEV+VAR**2
T0038 SAFA3=SAFA3+VAR**3
T0086 30 SAFA4=SAFA4+VAR**4
T0170 IF (N-30)37,37,35
T0238 35 TOTAL=N
T0262 GO TO 38
T0270 37 TOTAL=N-1
T0306 38 STDV2=SDEV/TOTAL
T0342 STDV=SQRT(STDV2)
T0366 ALFA3=SAFA3/(TNMBR*STDV**3)
T0426 ALFA4=SAFA4/(TNMBR*STDV**4)
T0486 REL=(TPTS*STDV2-AMEAN*(TPTS-AMEAN))/(STDV2*(TPTS-1.0))
T0654 IF (SENSE SWITCH 2)85,89
T0674 85 PUNCH 889,AMEAN
T0698 PUNCH 900,STDV,ALFA3
T0734 PUNCH 901,ALFA4,REL
T0770 89 PRINT 776

```

```

T0794 PRINT 889,AMEAN
T0818 PRINT 778
T0842 PRINT 900,STDV,ALFA3
T0878 PRINT 778
T0902 PRINT 901,ALFA4,REL
T0938 PRINT 778
T0962 IF (SENSE SWITCH 2)39,42
T0982 39 PUNCH 390
T1006 41 IF (SENSE SWITCH 3)42,44
T1026 42 PRINT 778
T1050 PRINT 442
T1074 PRINT 443
T1098 PRINT 776
T1122 44 OLD=1000.0
T1146 READ 100,TPTS,Z1,Z2,Z3,Z4,Z5
T1230 DO 83 N=i,NMBR
T1242 READ 100,SCRE(N),Z1,Z2,Z3,Z4,Z5
T1350 J=N
T1374 IF (N-NMBR)51,63,99
T1442 51 IF (SCRE(N)-SCRE(N+1))16,52,62
T1554 52 IF (J-1)99,54,53
T1622 53 IF (SCRE(N-1)-SCRE(N))16,64,54
T1734 54 SAVEN=N
T1758 55 N=N+1
T1794 IF (N-NMBR)57,58,99
T1862 57 IF (SCRE(N)-SCRE(N+1))16,55,58
T1974 58 THISN=N
T1998 ADD=(THISN-SAVEN)/2.0
T2046 SRANK=SAVEN+ADD
T2082 GO TO 64
T2090 62 IF (J-1)99,66,63
T2158 63 IF (SCRE(N-1)-SCRE(N))16,64,66
T2270 64 RANK=SRANK
T2294 GO TO 67
T2302 66 RANK=N
T2326 67 N=J
T2350 PRANK=((TNMBR-RANK)*100.0)/TNMBR
T2410 ZSCRE=(SCRE(N)-AMEAN)/STDV
T2482 TSCRE=ZSCRE*10.0+50.0
T2530 IF (SENSE SWITCH 2)73,75
T2550 73 PUNCH 100,SCRE(N),N,RANK,PRANK,ZSCRE,TSCRE
T2658 IF (SENSE SWITCH 3)75,82
T2678 75 IF(OLD-SCRE(N))16,76,78
T2766 76 PRINT 776
T2790 GO TO 79
T2798 78 PRINT 778
T2822 79 PRINT 100,SCRE(N),N,RANK,PRANK,ZSCRE,TSCRE
T2930 82 OLD=SCRE(N)
T2978 83 CONTINUE
T3014 99 STOP
T3022 END
END OF COMPILATION T354IT3810

```

Legend: Source Program

ALFA 3	=	alpha-three
ALFA 4	=	alpha-four
AMEAN	=	arithmetic mean
NMBR	=	total number of persons taking the test
PRANK	=	percentile rank
REL	=	reliability of the test
SAFA 3	=	$\sum_{i=1}^n (X_i - \bar{X})^3$
SAFA 4	=	$\sum_{i=1}^n (X_i - \bar{X})^4$
SCORE	=	individual test score (raw score)
SDEV	=	$\sum_{i=1}^n (X_i - \bar{X})$
SSCRE	=	$\sum_{i=1}^n X_i$
STDV	=	standard deviation = $\sigma$
STDV 2	=	$\sigma^2$
STDV 3	=	$\sigma^3$
STDV 4	=	$\sigma^4$
TNMBR	=	total number of persons taking the test.
TOTAL	=	TNMBR - 1.0      if TNMBR $\leq$ 30
	=	TNMBR            if TNMBR $\geq$ 30
TPTS	=	total possible points or total number of test items, which ever is greater.

TSCORE = T-SCORE  
VAR =  $(X_i - \bar{X})$   
ZSCORE = Z-SCORE  
Z1 - Z5 = Dummy Variables

ARITHMETIC MEAN IS 82.256  
 STANDARD DEVIATION IS 11.094  
 ALPHA 3 IS -.828  
 ALPHA 4 IS 3.695  
 RELIABILITY OF TEST IS .890

STUDENT NO AND NAME	SCORE	NO	RANK	PCNT RANK	Z-SCRE	T-SCRE
12550 BUSH DONALD JOSEPH	98.0	1	1.50	96.153	1.419	64.190
57795 KICKSCHL JAMES HAROLD	98.0	2	1.50	96.153	1.419	64.190
28336 FOSLER LARRY RICHARD	96.0	3	4.50	88.461	1.238	62.387
65600 PENNINGTON GAIL LEA	96.0	4	4.50	88.461	1.238	62.387
84890 TATEM JOHN DAVID	96.0	5	4.50	88.461	1.238	62.387
68090 PROPST DARRELL	96.0	6	4.50	88.461	1.238	62.387
74390 SCHAEFER JAMES MARTIN	95.0	7	7.00	82.051	1.148	61.486
42645 JENSEN JAY WAYNE	94.0	8	8.00	79.487	1.058	60.585
74565 SCHERB JOHN FREDERICK	92.0	9	9.00	76.923	.878	58.782
25890 EVANS MELVIN ERNEST	88.0	10	13.00	66.666	.517	55.177
33150 GRIFFITH JOHN H	88.0	11	13.00	66.666	.517	55.177
43336 JOHNSON STEPHEN R	88.0	12	13.00	66.666	.517	55.177
44700 KEEGAN GARY LEE	88.0	13	13.00	66.666	.517	55.177
47600 KRENKE GLEN LEE	88.0	14	13.00	66.666	.517	55.177
48006 KUYKENDALL JOHN A	88.0	15	13.00	66.666	.517	55.177
53125 MAPES REX MARSHALL	88.0	16	13.00	66.666	.517	55.177
03050 BAKER DOYLE SIMON	86.0	17	18.00	53.846	.337	53.374
70260 RICE JAMES HAROLD	86.0	18	18.00	53.846	.337	53.374
75778 SEIM KENNETH BRUCE	86.0	19	18.00	53.846	.337	53.374
54750 MC ATEE DON SCOTT	84.0	20	20.50	47.435	.157	51.571
55975 MC DOWELL MARSHA G	84.0	21	20.50	47.435	.157	51.571
39370 HOFFMAN DAVID CHARLES	82.0	22	22.50	42.307	-.023	49.768
76548 SHEEHAN JAMES EVERETT	82.0	23	22.50	42.307	-.023	49.768
60600 MOYER MICHAEL A	80.0	24	24.50	37.179	-.203	47.966
77387 SHOULDERS MICHAEL C	80.0	25	24.50	37.179	-.203	47.966
17800 CORBIN JAMES B	78.0	26	26.50	32.051	-.383	46.163
24965 ELLER LARRY ALLEN	78.0	27	26.50	32.051	-.383	46.163

24310	ECHARD WM RANDALL	76.0	28	29.50	24.358	-.563	44.360
64265	PAIGE WANDA ELAINE	76.0	29	29.50	24.358	-.563	44.360
64890	PATRICK GERRY WAYNE	76.0	30	29.50	24.358	-.563	44.360
2725	WILDER RICHARD LYNN	76.0	31	29.50	24.358	-.563	44.360
55475	MC COY CAROL RUTH	74.0	32	32.00	17.948	-.744	42.558
22977	DOUGHERTY CHARLES E	70.0	33	33.50	14.102	-1.104	38.952
25383	EMMERT ROBERT CARL	70.0	34	33.50	14.102	-1.104	38.952
02025	ARMSTEAD ROBERT LEE	69.0	35	35.00	10.256	-1.194	38.051
69130	RAUSCH MICHAEL	64.0	36	36.50	6.410	-1.645	33.544
73936	SANDERS JAMES W	64.0	37	36.50	6.410	-1.645	33.544
62097	NEWHARD NANCY FAYE	62.0	38	38.00	2.564	-1.825	31.741
28065	FOLTA JAMES VINCENT	48.0	39	39.00	.000	-3.087	19.122
	STOP						



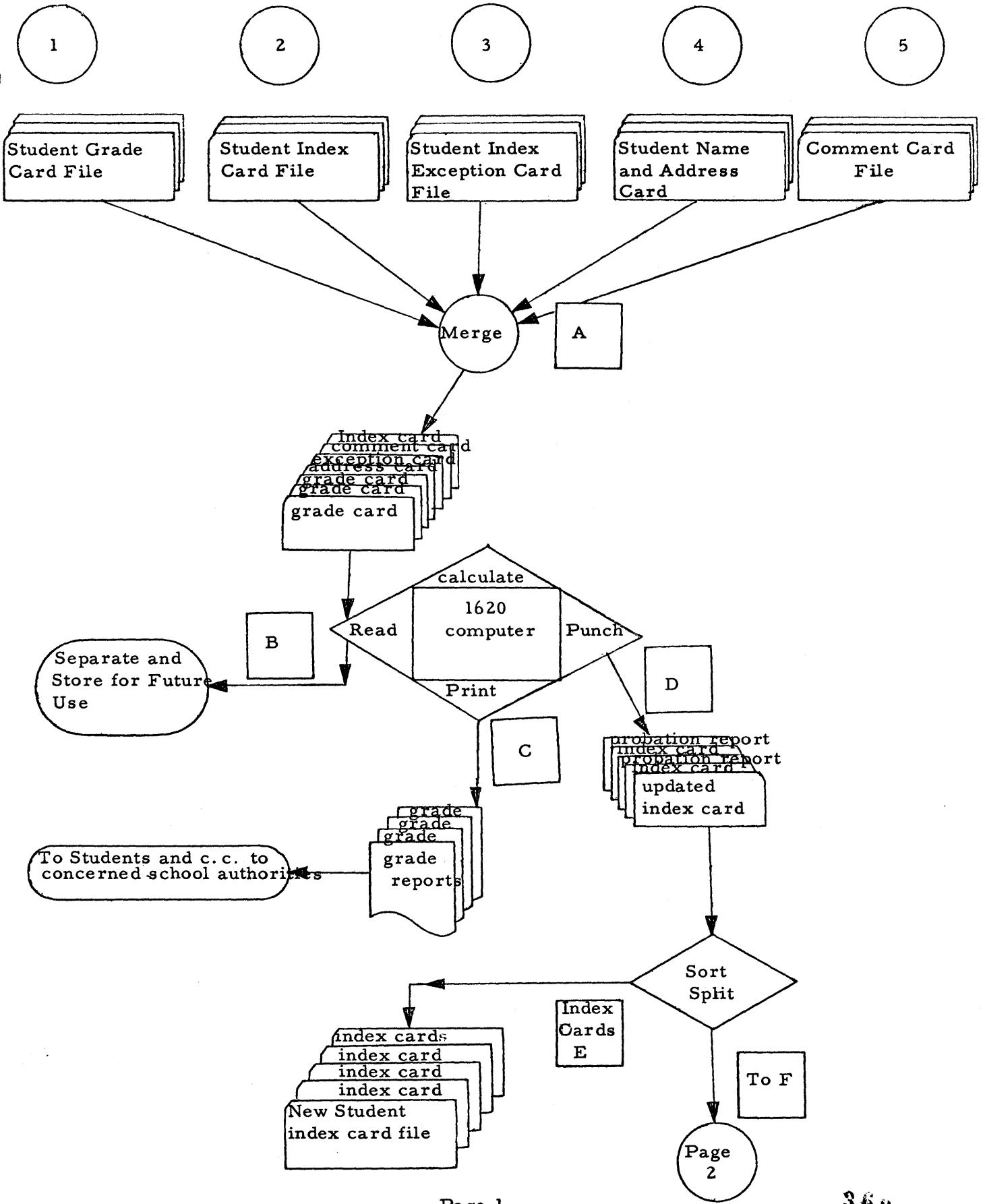
A  
PROPOSAL FOR AN  
ADDITION IN GRADE REPORTING  
PROCEDURES TO ALLOW FOR  
AUTOMATIC PROCESSING OF PROBATION STUDENTS

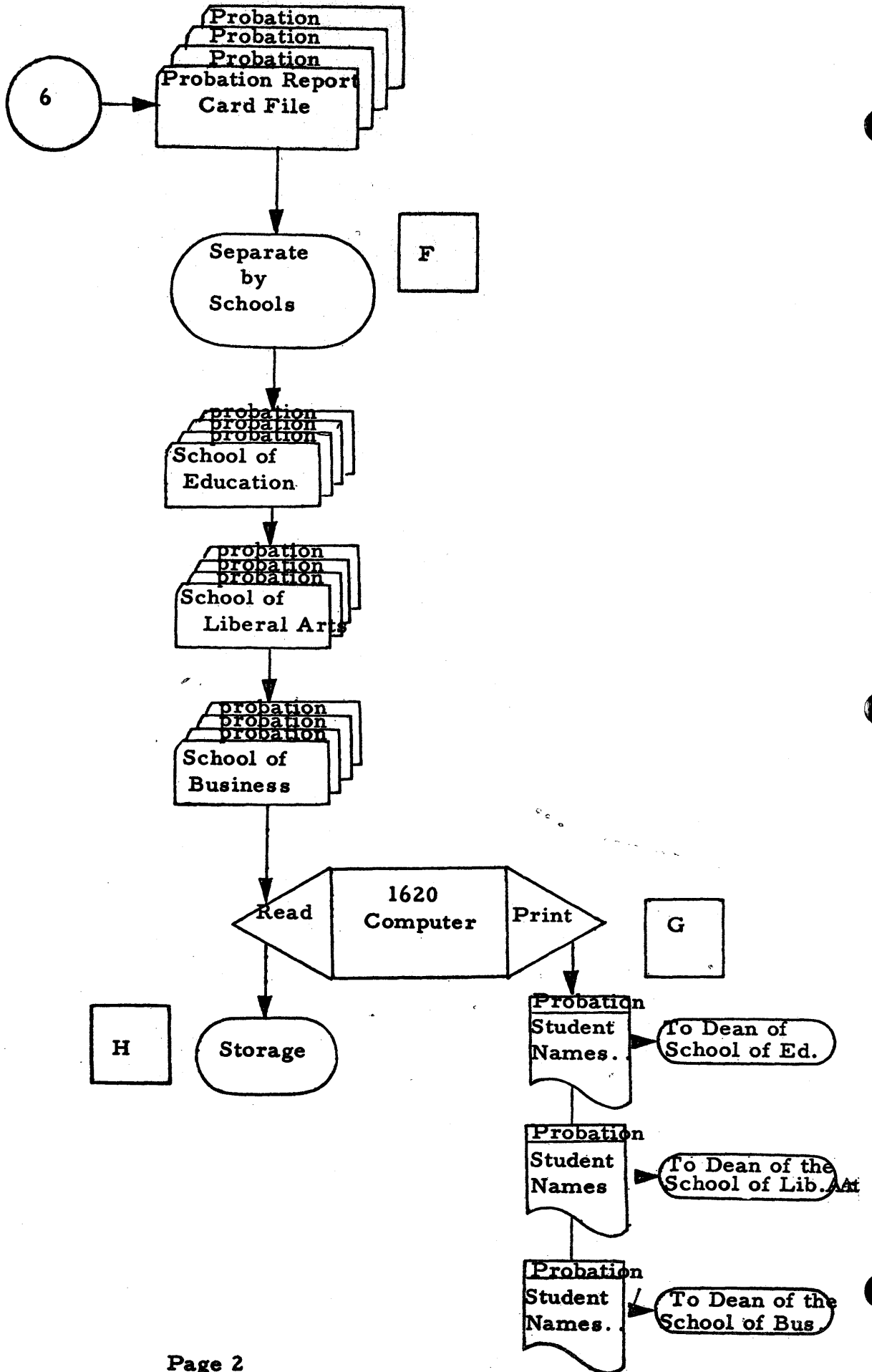
Indiana State College  
Computer Center  
May 5, 1964

**FLOW  
OF  
PROCEDURE**

### Explanation of Flow Chart

- A All card files are merged together in alphabetical sequence by student number.
- B Computer writes grade reports and updates student index file. Probation cards for failing students. Store used cards.
- C Grade reports to students and school officials.
- D Output: New student index cards and probation report writing cards.
- E New index cards go back into file for report use.
- F Probation report cards are sorted by school.
- G Probation reports are printed and sent to the deans of the schools.
- H Probation cards are stored.





CARD  
JUSTIFICATION

## Explanation of Card File

After completion of Step A (shown by the flow chart), the input to the computer consists of multiple card groups, one per student. It is the purpose of this report to show the need for including each type of card.

The speed and efficiency of any data processing procedure are largely dependant upon the volume of data to be processed. Because card volume is of such importance, it is to the users advantage to keep it at a minimum.

The five types of input data cards necessary for the student academic progress reports follow. They are:

### Card 1 - Student Class Grade Card

These cards enter the flow through the registration line. They are the yellow striped cards the student submits for each class he attempts. After registration, the cards are held until the instructors turn in their grades. Each grade, with the respective grade points, is entered on the correct grade card. The result of this activity leaves a workable file of all work completed on punched cards.

### Card 2 - Student Index Card

A continuous file of student index cards is maintained by the Computer Center. This file records the complete scholastic history and present status of each student. It is this card that records the amount of college work completed, with the grade average earned for this work - on a cumulative basis, and also on a single semester basis.

Such information as where the student lives, what social organization does he belong to, the number of credit hours transferred in from other colleges, his first two major areas of study, his minor area of study, and the sex, is all recorded in a numerical code on this card.

It is the student index card that facilitates all reports on academic progress; from a report showing the current and cumulative index of each girl living on the third floor of Reeve Hall, to a report of the numbers of hours all Education majors carried any given semester.

Because this card is so vital to our work, and because the grade report contains both cumulative and current credit hours, grade points earned, and grade point ratio, it is necessary for this card to be re-computed at each semester's end.

### Card 3 -- Student Required Index Exception Card

The purpose of this card file is to automate the detailed processing of students having scholastic problems. Because the student's academic progress is of utmost concern to the college, careful monitoring and guidance techniques are essential. The inclusion of this card greatly facilitates much of the detailed analysis work necessary.

The student index card allows the proper school authorities to carefully supervise the progress of a student. By submitting a probation form to the Computer Center, a school official can stipulate exactly what scholastic level of achievement must be met. This is done by simply stating what grade point ratio the student must earn, either on a cumulative, or semester basis. This information is then entered into the student's card group and allows the computer to analyze the student's work accordingly.

If the student fails to meet this requirement, the computer will generate a card from which a complete scholastic report can be written and sent to the appropriate official.

The card is labeled "exception" because, in the absence of such a card, the computer will use the standard required index schedule to analyze the student. See Probation Scaling.

### Card 4 -- Student Name and Address Card

The name and address card allows for automatic addressing of the grade report.

### Card 5 -- Comment Card

The comment card allows a school official a maximum of two lines (68 characters) of comment on the student grade report. Through the use of a comment code, the comment may be printed only if the student fails to meet specified grade conditions. It is also possible to print the comment under any conditions. See Comment Printing.



## Output Data Cards

Two types of cards are generated by the computer. The first, an updated student index card, replaces the input student index card (see card 2 - student index card).

The second, a probation report card (card 6), is punched for every student falling below certain minimum grade average requirements.

The purpose of this card is to allow a file of cards to be maintained on all probation students. A more complete description of this will be found under Probation Report.

## Comment Printing

The grade processing procedure utilizes two kinds of comments. The first type, those entered on the comment card, may be worded as the school official desires. However, because the probation report sent to the school authorities should explain what kind of comment was made, and who authorized the comment, the inclusion of an authority code, and of a comment classification code is necessary.

The classification of comments is as follows (without respect to wording):

<u>Code</u>	<u>Description</u>
1	Place on Probation
2	Contact the Dean of your school
3	Contact the Registrar
4	Withdrawn from school for scholastic reasons
5	Withhold Permit to Register
6	
7	
8	
9	Machine generated probation comment.

The authority codes are as follows:

<u>Code</u>	<u>Description</u>
1	Registrar
2	Dean of Students
3	Business Office
4	Dean of the School of Ed.
5	Dean of the School of Lib. Arts

6  
7  
8  
9

Computer Center

It should be noted that the presence of a comment does not necessarily mean a student is on probation.

A special punch over the comment code will suppress printing if the required index is met or exceeded. Thus, a code 1 and this special punch would read J, a 2 would read K, 3 - L, etc.

The second form of comment, machine-generated, are those printed on the grade report to notify the student he is being placed on probation, or that he is being removed from probation status. These comments will be identified by a code 9.

## Probation Scaling

As each student index card enters the computer during grade processing, the probation code found in card column 47 will be examined. A zero in this column indicates the student is not on probation this semester, any other digit means that the student was on probation last semester, i. e., a 3 would indicate a probation student for the 3 preceding semesters. If the student again fails to earn a satisfactory grade index, the probation code will be incremented by 1. If the student earns a satisfactory average, the probation code will be made zero.

As each student's grades are processed and the new cumulative hours and grade points are brought up to date, the computer will, in the absence of a required index exception card, scale the cumulative hours and find the required grade-point average. The scale is as follows:

<u>Cumulative Hours</u>	<u>Required Grade-Point Average</u>
0 - 16	1.00
17 - 32	1.25
33 - 45	1.50
46 - 60	1.80
61 +	2.00

Any student who does not meet or exceed this scale will be automatically placed on probation. The probation code for such students will be incremented and a "Probation Report Card" (card 6) will be generated.

The computer will notify such students of this condition by - in the absence of a comment card - printing on the grade report. See comment printing.

## Special Provisions

If a "Required Index Exception Card" (card 3) is present, their grade-point average will be scaled as specified by this card.

CARD  
FORMAT

Card  
1

Student Class Grade Card

Card Column

1-5  
6  
7-27  
28  
29  
30-32  
33-36  
37-40  
41-42  
43-57  
58-61  
62-65  
66-70  
71-72  
73-74  
75-77  
79  
80

Description

Student Number  
School  
Name  
Classification  
Curriculum  
Number in Class  
Department  
Course Number  
Section Number  
Course Description  
Course Number  
Time Class Meets  
Days Class Meets  
Semester Code  
Grade  
Grade Points  
Hours of Credit  
Code 1

Card  
2

Student Index Card

<u>Card Column</u>	<u>Description</u>
1-5	Student Number
6	School
7-27	Name
28-31	Cumulative Hours
32-35	Cumulative Points
36-38	Cumulative Grade Point Ratio
39-40	Semester Hours
41-43	Semester Points
44-46	Semester Grade Point Ratio
47	Probation Code
48	Grade Point Ratio
49	Hours
50-51	Housing
58-61	Total Hours Toward Graduation
62-63	Hall
64-65	House
66	Social Organization
67-69	Semester Code
70	New or Transfer Students
71	Teaching or Non-Teaching
72-73	Minor
74-75	Second Major
76-77	First Major
78	Sex
80	Code 2

Card

3

Student Required Index Exception Card

<u>Card Column</u>	<u>Description</u>
1-5	Student Number
6	School
7-27	Name
28	Authority
29-31	Required Cumulative Index to be Earned
32-34	Required Semester Index to be Earned
35-37	Semester Code
80	Code 4

Card Format

Card  
4

Student Name and Address Card

Card Column

1-5

6

7-27

28-51

52-72

80

Description

Student Number

School

Name

Street Address

City and State

Code 1



Card

5

Comment Card

Card Column

1-5  
6  
7  
8  
9-11  
12-45  
46-79  
80

Description

Student Number  
School  
Authority  
Comment Code  
Semester Code  
First Line of Comment  
Second Line of Comment  
Code 5

Card Coding For Schools

Card Column 6 Of All Cards

<u>Code</u>	<u>School</u>
1	Education
2	Liberal Arts
3	Business
4	Nursing
5	Other

Card Coding For Classification

Card Column 49 Student Index Card

Card Column 75 of Probation Report Card

<u>Code</u>	<u>Classification</u>	<u>Cumulative Number of Hours</u>
1	Freshman	0 - 27
2	Sophomore	28 - 56
3	Junior	57 - 85
4	Senior	86 - 124
5	Other	

FORMS









COST STUDY  
PROGRAM ABSTRACTS

Programs 1, 2, and 3



## PROGRAM 1

Description:

This program extends the enrollment and class hour figures in the 2 cards. The input, the 2 cards sorted by Department Number, is read from the 1622 card reader. The computer multiplies the enrollment figure times the class hour figure. The entire 2 card is reproduced for output - containing the extended amount in card columns 26-31.

When a change in Department Number occurs, the machine will generate a 3 card. Therefore, there should be one 3 card for every department entered. The 3 card will contain: The Year, Dept. No., Total Student Class Figure, and the Total FTE Staff Figure. Both totals are from the 2 cards.

Error Conditions:

As the file goes through the machine, a sequence check for equal groups is performed. If a 2 card is out of order an error message will be typed on the typewriter ("error in Dept. sequence."). Because the computer recognizes an error condition on an "Not Equal and Not High" compare break, the department to which the error card belongs may have already passed completely through. In such a case it will be necessary to:

1. Adjust the 3 card for that department.
2. Rerun the department with the card included.

Irrespective of which method is used, the run must be started from the first department for which there is no 3 card.

Switch Settings:

I/O - Stop  
Parity - Stop  
Overflow - Stop

Console Switches:

Sw. 1 Not Used  
Sw. 2 Not Used  
Sw. 3 See Operating Suggestions  
Sw. 4 Not Used

## PROGRAM 2

Description:

This program is used to pro-rate the number two cards. After Program 1 is complete, the Department Total cards (3 cards) have two total expense amounts keypunched in them. These cards are then placed back in the file of 2 cards in front of their respective department files. This program (Program 2) then accepts that file as data. The program accepts the data in sequential order (by Dept. No. - 3 card first, followed by the two cards) and pro-rates the two amounts from each 3 card into each two card.

The two amounts to be allocated are found in card columns 19-24 and 25-30 of each 3 card. The first amount is pro-rated by weighting the student class hours figure on each 2 card against the total student class hour amount on its respective 3 card. This is then applied to the first total (stored in the machine) for each 2 card's share. The second amount is pro-rated on a similar basis, i.e., FTE figures are weighted. The remainders are carried into the following dividend, allowing the last card to zero-balance the amount to be allocated and the amount allocated.

Error Conditions:

The program tests for six error conditions during processing. If an error condition is detected, the computer will type "Error N" and halt. The six error conditions follow:

Error	
1	Sequence - 3 card
2	Sequence - 2 card
3	Amount allocated on student hours doesn't zero balance
4	Amount allocated on FTE doesn't zero balance
5	Student hours from 2 cards do not equal total from 3 card
6	FTE percentages from 2 cards do not equal total from 3 card

On any of the above conditions, except 1 and 2, it will be necessary to start from the last department not completely processed.

PROGRAM 2

Switch Settings:

I/O - Stop  
Parity - Stop  
Overflow - Stop

Console Switches:

Sw. 1 Not Used  
Sw. 2 Not Used  
Sw. 3 Must be same setting as in Program 1  
Sw. 4 Not Used

## PROGRAM 3

Description:

This program allocates the pro-rated totals on the two cards to a 7 card for each course. The output from Program 2 is sorted into course within department order (card columns 5-10) with the 7 cards then being merged in. If the data is in correct order, there will be one seven card in front of every group of 2 cards.

The program accepts its data through the card reader. The computer will read all the data cards for an entire department, then punch a new 7 card with the sum of the three expense totals from the two cards (FTE Expense + Salary Expense + Student Class Hours Expense) allocated to the six different grade levels (Freshmen, Sophomore, Junior, Senior, Special, and Graduate). The total enrollment of each class is weighted against the remainder from each division is carried to the following dividend to zero-balance the last card of each group.

The 7 card generated, then, has the complete cost of the course allocated to each class of students.

Error Conditions:

The program is extremely limited as to the number of error conditions it can check. The 7 cards are presumed to have been crossfooted, and the 2 cards have been through the machine checks of Program 1 and 2. Three error conditions might arise. They are:

1. An unmatched two card
2. A sequence error
3. The 7 card allocation did not zero balance - This would happen if a 2 card was missing.

On any error condition, the typewrite will describe the error and halt. It is then necessary to start over, from the last course 7 card not punched.

Switch Settings:

I/O - Stop  
Parity - Stop  
Overflow - Stop

Console Switch Settings:

Not Used

## OPERATING SUGGESTIONS

1. Before running any program, clear the machine. This is done by pressing "Reset", then "Insert" on the computer console. Type in 260000800009, then press the R-S key. Wait about 1/8 of a second and depress "Instant Stop". Then press "Reset".
2. After clearing the machine, place one of the object decks in the read hopper and press "Load" (This is the yellow button). After the machine stops reading, press the green "Reader Start" to read in the last card. The machine will be in the manual mode at this point. Load the data in the read hopper, and blank cards in the punch hopper and press Start to begin processing.
3. If programs one and two are used with switch 3 off, there will be a duplication of output. That is, Program 1 extends the two cards with switch 3 off. This may be avoided by using the programs with switch 3 on (Programs 1 and 2). This method will allow just Program 2 to generate the 2 cards. Program 1 will then punch only the 3 cards and Program 2 will extend the 2 cards as it pro-rates them. Irrespective of the switch setting used, it must be the same for both programs. See the following flow chart for a more complete description.

Equipment Specifications  
20K 1620 Computer  
1622 Card/Punch Unit

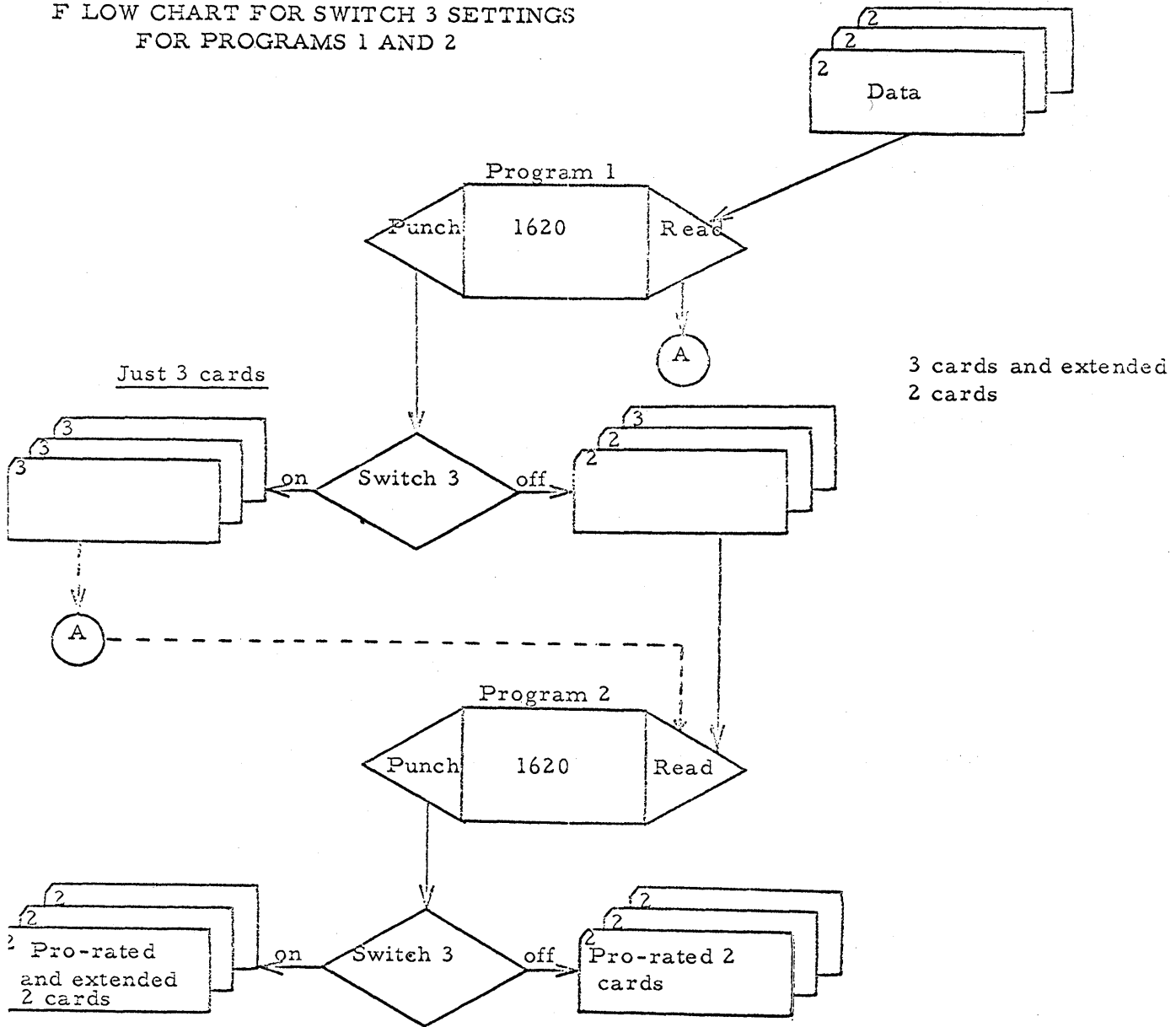
Additional Features  
TNS and TNF Instructions\*  
Indirect Addressing Not Used

\*These instructions greatly facilitate numeric conversion, but if not available, they may be simulated with several transmit digit instructions. As an example of this substitution, Program 1, line 01090 (TNS in+10, Now, Dept. No. 3-6) might be changed to:

TD	Now, in+10
TD	Now - 1, in+8
TD	Now - 2, in+6
TD	Now - 3, in+4
SF	Now - 3

This is because a number such as 0036 will be read as 70707376.

FLOW CHART FOR SWITCH 3 SETTINGS  
FOR PROGRAMS 1 AND 2



(A)

The output from Program 1 with Switch 3 ON is merged in the input file. This then becomes the input for Program 2.





A  
REPORT  
OF  
SCHOLASTIC ACHIEVEMENT  
FOR  
SPRING SEMESTER  
1963-64

June 4, 1964

395

This report contains a summary of information by department, residence halls, sororities, fraternities, non-sorority, non-fraternity and by total student body.

This report deals exclusively with the undergraduate student body.

All Students

	<u>Number of Students</u>	<u>Hours Attempted</u>	<u>Grade Points Earned</u>	<u>Grade-Point Ratio</u>
Men	3,150	40,290	89,102.0	2.21
Women	<u>2,321</u>	<u>29,697</u>	<u>74,730.0</u>	<u>2.52</u>
Total	<u>5,471</u>	<u>69,987</u>	<u>163,832.0</u>	<u>2.34</u>

Fraternities\*

	<u>Number of Students</u>	<u>Credit Hours Attempted</u>	<u>Grade Points Earned</u>	<u>Grade-Point Average</u>
Alpha Tau Omega	46	643	1,500.5	2.33
Lambda Chi Alpha	102	1,332	3,163.5	2.38
Pi Lambda Phi	31	430	988.0	2.30
Sigma Phi Epsilon	105	1,407	3,239.5	2.30
Tau Kappa Epsilon	87	1,186	2,705.0	2.28
Theta Chi	42	608	1,433.0	2.36
Total	<u>413</u>	<u>5,606</u>	<u>13,029.5</u>	<u>2.32</u>
Non-Fraternity Men	2,737	34,684	76,072.5	2.19

\*Actives Only

Sororities\*

	<u>Number of Students</u>	<u>Credit Hours Attempted</u>	<u>Grade Points Earned</u>	<u>Grade-Point Average</u>
Alpha Omicron Phi	55	774	2,072.5	2.68
Alpha Phi	39	594	1,642.5	2.77
Alpha Sigma Alpha	35	518	1,400.5	2.70
Chi Omega	50	699	1,916.0	2.74
Delta Gamma	54	796	2,231.0	2.80
Gamma Phi Beta	60	860	2,334.5	2.71
Sigma Kappa	55	778	2,109.0	2.71
Zeta Tau Alpha	53	779	2,080.0	2.67
<b>Total</b>	<b>401</b>	<b>5,798</b>	<b>15,786.0</b>	<b>2.72</b>
<b>Non-Sorority Women</b>	<b>1,920</b>	<b>23,899</b>	<b>58,944.0</b>	<b>2.47</b>

\*Actives Only

By Area Major

Major	Number of Students	Hours Attempted	Grade Points Earned	Grade-Point Average
Art	113	1,286	2,974.5	2.31
Business	794	10,073	22,351.0	2.22
English	302	3,895	10,128.0	2.60
Foreign Language	137	1,300	3,696.0	2.84
Home Economics	165	2,131	5,162.0	2.42
Industrial Education	418	5,351	11,920.0	2.23
Journalism	5	37	69.5	1.88
Liberal Arts	102	972	1,801.0	1.85
Mathematics	267	3,552	8,550.0	2.41
Music	148	2,090	4,977.5	2.38
Nursing	85	790	1,950.5	2.47
Phy. Ed. Men	421	5,511	11,252.5	2.04
Phy. Ed. Women	108	1,384	3,220.0	2.33
Science	517	6,803	16,693.5	2.45
Social Studies	587	7,740	17,654.0	2.28
Speech and Radio	115	1,538	3,757.5	2.44
Special Education	148	1,939	4,819.5	2.49
Elem. - Fresh.	342	4,729	10,175.5	2.15
Elem. - Soph.	245	3,330	7,835.0	2.35
Elem. - Jr.	280	3,964	10,269.5	2.59
Elem. - Sr.	172	1,572	4,575.0	2.91
Total	5,471	69,987	163,832.0	2.34

By Residence Halls

Female

	<u>Number of Students</u>	<u>Credit Hours Attempted</u>	<u>Grade Points Earned</u>	<u>Grade-Point Ratio</u>
Burford Hall	286	4,194	10,117.0	2.41
Erickson Hall	277	4,046	10,004.0	2.47
Pickerl Hall	290	4,334	10,669.5	2.46
Reeve Hall	326	4,724	11,392.5	2.41
Total	<u>1,179</u>	<u>17,298</u>	<u>42,183.0</u>	<u>2.44</u>

Male

	<u>Number of Students</u>	<u>Credit Hours Attempted</u>	<u>Grade Points Earned</u>	<u>Grade-Point Ratio</u>
Gillum Hall	304	4,338	9,411.5	2.17
Hulman Center	309	4,163	9,049.0	2.17
Parsons Hall	298	4,326	8,452.0	1.95
Sandison Hall	298	4,360	9,748.5	2.24
Total	<u>1,209</u>	<u>17,187</u>	<u>36,661.0</u>	<u>2.13</u>

Other Housing

	<u>Number of Students</u>	<u>Credit Hours Attempted</u>	<u>Grade Points Earned</u>	<u>Grade-Point Ratio</u>
Men	1,941	23,103	52,441.0	2.27
Women	<u>1,142</u>	<u>12,399</u>	<u>32,547.0</u>	<u>2.62</u>
Total	<u>3,083</u>	<u>35,502</u>	<u>84,988.0</u>	<u>2.39</u>

By Residence Hall

Burford Hall

<u>House</u>	<u>Number of Students</u>	<u>Credit Hours Taken</u>	<u>Grade Points Earned</u>	<u>Grade-Point Ratio</u>
2	52	743	1846.0	2.48
3	58	902	2302.0	2.55
4	60	854	1975.0	2.31
5	59	867	1937.0	2.23
6	57	828	2057.0	2.48
<b>Total</b>	<b>286</b>	<b>4194</b>	<b>10117.0</b>	<b>2.41</b>



By Residence Hall

Erickson Hall

House	Number of Students	Credit Hours Taken	Grade Points Earned	Grade-Point Ratio
2	57	861	2048.5	2.38
3	56	832	2113.5	2.54
4	54	766	1879.5	2.45
5	52	757	1827.0	2.41
6	58	830	2135.5	2.57
Total	<u>277</u>	<u>4046</u>	<u>10004.0</u>	<u>2.47</u>

By Residence Hall

Gillum Hall

Hours	Number of Students	Credit Hours Taken	Grade Points Earned	Grade-Point Ratio
2	39	561	1149.5	2.05
3	39	553	1248.0	2.26
4	38	549	1157.5	2.11
5	35	492	968.0	1.96
6	38	533	1128.5	2.12
7	40	550	1131.5	2.06
8	36	521	1184.0	2.27
9	39	579	1444.5	2.49
Total	304	4338	9411.5	2.17

By Residence Hall

Hulman Center

House	Number of Students	Credit Hours Taken	Grade Points Earned	Grade-Point Ratio
3	33	442	956.0	2.16
4	53	711	1514.0	2.13
5	59	778	1737.5	2.23
6	57	730	1555.0	2.13
7	56	805	1837.5	2.28
8	51	697	1449.0	2.08
Total	309	4163	9049.0	2.17

By Residence Hall

Parsons Hall

House	Number of Students	Credit Hours Taken	Grade Points Earned	Grade-Point Ratio
1	57	843	1786.0	2.12
2	44	643	1312.0	2.04
3	56	820	1544.0	1.88
4	36	530	1061.0	2.00
5	57	832	1471.5	1.77
6	48	658	1277.5	1.94
Total	298	4326	8452.0	1.95

By Residence Hall

Pickerl Hall

House	Number of Students	Credit Hours Taken	Grade Points Earned	Grade-Point Ratio
2	56	827	2208.0	2.67
3	57	906	2111.0	2.33
4	58	846	2070.5	2.45
5	59	865	2071.0	2.39
6	60	890	2209.0	2.48
Total	290	4334	10669.5	2.46

By Residence Hall

Reeve Hall

House	Number of Students	Credit Hours Taken	Grade Points Earned	Grade-Point Ratio
1	19	270	610.0	2.26
2	87	1235	2853.5	2.31
3	114	1707	4172.5	2.44
4	106	1512	3756.5	2.48
Total	326	4724	11392.5	2.41

By Residence Hall

Sandison Hall

House	Number of Students	Credit Hours Taken	Grade Points Earned	Grade-Point Ratio
1	13	182	301.0	1.65
2	36	545	1324.5	2.43
3	35	521	1147.5	2.20
4	37	527	1230.0	2.33
5	34	473	1073.5	2.27
6	37	546	1362.5	2.50
7	36	542	1207.5	2.23
8	34	492	1002.5	2.03
9	36	532	1099.5	2.07
Total	298	4360	9748.5	2.24

AUTOMATIC PROCESSING OF AUTOSPOT AND AUTOMAP  
PROGRAMS WITH THE 1620-1311 DISC SYSTEM

A Paper Prepared For The:

Spring Meeting  
1620 Users Group  
Brown Palace Hotel  
Denver, Colorado  
June 19, 1964

BY:

Jack T. Dunn  
AVCO Corporation  
Huntsville, Alabama

And

Ernie G. Moore  
IBM Corporation  
Huntsville, Alabama



The authors of this paper gratefully acknowledge the work of  
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I

TEXT

## INTRODUCTION

Our purpose in presenting this paper is to illustrate a technique, developed at AVCO Corporation's Huntsville Engineering Office in conjunction with IBM Corporation's Huntsville Branch Office, Systems Engineering Staff, for processing AUTOSPOT and AUTOMAP programs on the 1620 equipped with an on line 1311 Disc Drive.

This is an automatic system in that it provides a means for producing a numerical control tape deck or paper tape from a deck of source statements without going through the multi-pass card system.

This technique is not reputed to be a panacea, nor will it be advantageous for all users.

It is, however, felt to be a significant step in improving the existing means of processing numerical control programs on the 1620.

During the next few minutes we will present to you information regarding:

- (1) Rationale which led us to choose this particular method of attack.
- (2) The methods utilized by and the operation of the system.
- (3) Advantages and limitations of the system.
- (4) Savings in time and money which might reasonably be expected, and
- (5) Application to Fortran and SPS programs.

Following the discussion of these topics, Mr. Ernie Moore of IBM, Huntsville, will discuss some of the specific modifications which must be made to the machine language deck in order that this system may be utilized.

I wish to invite those of you who are not specifically interested in numerical control to visualize the decks which we will be discussing as nothing more or less than machine language object decks compiled under FORTRAN with FORMAT or SPS for 1620 Mod I.

In this way it may be possible for you to anticipate uses for this type of technique within your own organization.

#### RATIONALE

For the benefit of those who are unfamiliar with AUTOSPOT and AUTOMAP it would probably be worthwhile to take a brief glance at the nature of these two popular numerical control languages.

AUTOSPOT is a language for positioning machine tools -- basically it is for point to point operations but has limited machining capability. An AUTOSPOT source program is ordinarily written in fixed format English language statements, which are then punched into cards. These source statements are then processed through a three phase General Processor called the AUTOSPOT processor.

Following this General Processor will be a multi-phase post processor.

The post-processing required for the machine tools which we utilize at AVCO, consist of three to six additional phases. One of these phases is a two pass phase; thus a total of ten passes through the computer is required. All but one pass produces an intermediate output deck which must be subsequently loaded with the appropriate processor -- seemingly "ad infinitum".

AUTOMAP is a contouring language which requires an English language source deck and two phases of general processing. In addition, three phases of post-processing are required for our particular machine tools. Again each phase produces an intermediate output deck which must be subsequently reloaded with the appropriate processor.

The final phase of each of these produces a punch paper tape or a card deck which contains the input to the special purpose computer attached to the machine tool.

This special purpose computer actually controls the machine tool.

Upon exposure to this system, one immediately sees the need for streamlining the operation.

We chose to accomplish this streamlining by adapting AUTOSPOT and AUTOMAP to the 1620 with an on line 1311 disc.

This approach was chosen essentially because of the versatility offered by this system in other applications such as scientific/engineering computation under MONITOR/FORTRAN II-D as well as the numerical control application.

The choice of the particular technique we eventually utilized was the result of considerable thought as to what constituted the best approach. It was also the result of experience gained traveling a number of blind alleys.

The initial thinking indicated that the best approach would be to recompile each phase of AUTOSPOT and AUTOMAP under MONITOR with the read-write statements changed where necessary. This would, presumably, have involved updating the SPS source decks for all phases of the general and post processors applicable to our work; plus compiling them as if they were utility programs. It seemed at the time that this would be a simple straightforward approach which would yield results with a minimum of re-programming.

However, upon closer examination of this idea, a number of problems became evident.

The most serious of these involved the floating point subroutines which were required. AUTOSPOT and AUTOMAP originally were written and compiled to utilize the "excess fifty floating point" system. On the other hand, we came to learn that our SPS II-D (the SPS for MONITOR I) utilizes exponential numbers.

This problem in itself seemed to be insurmountable from a practical standpoint.

One other problem which we came to consider practically insurmountable is the fact that almost all phases of our general and post-processors fill 20,000 positions of core practically to the hilt. Thus no space would exist for the "supervisor" program which MONITOR requires.

These two problems alone seemed sufficient to eliminate the MONITOR approach as being simple.

One might reason that a 40 or 60K machine would eliminate the latter problem; and rightly so.

Unfortunately, we have a 20K machine -- moreover we desired, if possible, to keep AUTOSPOT and AUTOMAP in such a form that only the basic 1620 plus 1311 would be required.

There were other considerations in addition to these which led us to abandon the MONITOR approach. For example, our experience with MONITOR I led us to believe that, even if we were able to "Monitorize" AUTOSPOT and AUTOMAP, the execution speed would be agonizingly slow due to the disc relocatable subroutines, etc.

Thus after consultation with our systems engineer, we decided to investigate what we then called, a "quick and dirty" approach. This approach, we visualized, would incorporate machine language modifications to the existing compressed object decks of the various phases which we desired to implement. It was planned that these changes would link the various phases together in such a way that



each would call the next from its respective place in disc storage. In addition it was planned that the intermediate output would be loaded in an area reserved on the disc rather than being punched in cards.

It appeared, at the time, that this type of approach would result in serious problems in coding since all changes would necessarily be in machine language and since spare core storage locations were at a premium in most cases.

Still the benefits which were expected seemed to warrant proceeding along this particular avenue.

Some of the anticipated benefits were:

- (1) Relatively high speed operation
- (2) Simple operation
- (3) The AUTOSPOT and AUTOMAP decks would remain in essentially the same form as "non disc" decks so that Users Group library patches could be utilized.

After a number of troublesome but not really serious problems, our expectations began to be realized -- first with AUTOMAP -- then with the post-processor used for the Bendix controlled Milwaukee Matic Model II -- finally with AUTOSPOT and the post-processor for the G. E. controlled Milwaukee Matic Model II.

True to its name the system was quick but it was not at all "dirty" -- rather it was quite smooth and clean.

## THE SYSTEM

In essence the system as it applies to AUTOSPOT operates in the following manner.

(Slide #5A)

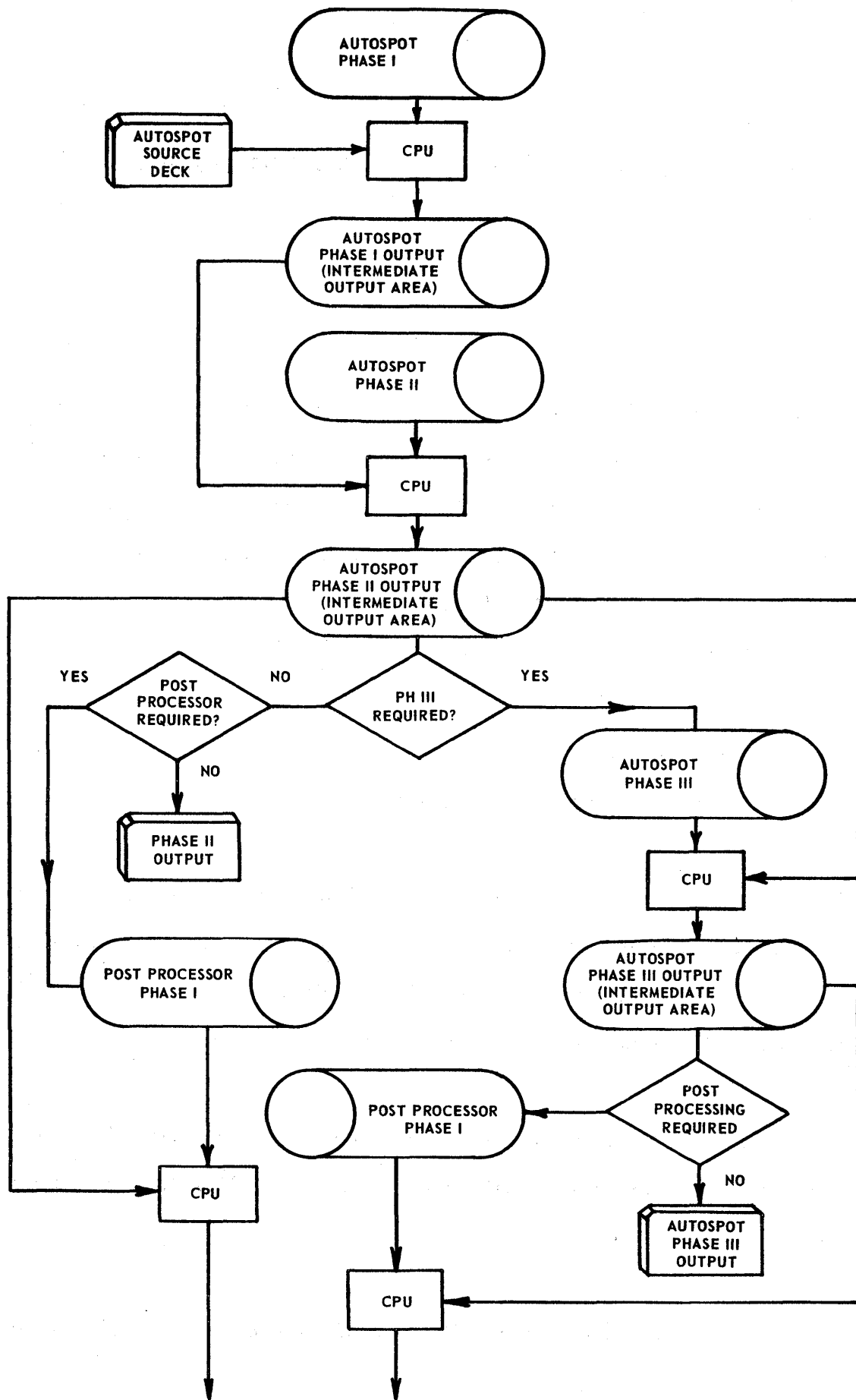
Phase I of the AUTOSPOT general processor resides in a specified position on the disc, its READ statements remaining unchanged and its PUNCH commands changed to load each output card image in an area on the disc reserved for intermediate output. This is accomplished, in general, by changing all the PUNCH commands to BRANCH to a routine which:

- (1) Writes the card image on the disc
- (2) Increments the sector address
- (3) Tests for end of job
- (4) Returns to main line program.

Thus, in order to process an AUTOSPOT program, one simply loads a control card which calls Phase I of the general processor from disc storage to memory. Phase I then reads and processes the source deck in the usual manner with the exception of punch output.

When the program encounters a FINI card (which signifies the end of job) the program branches to a routine which calls Phase II of the general processor from its disc location.

The READ commands of Phase II are modified to BRANCH the program into



a routine which reads a card image from the disc, increments the sector address and returns to the main line program. In this way the intermediate output from Phase I is read: processing proceeds in the usual manner and the resulting output is written on the disc in a fashion similar to Phase I.

When the processing of Phase II is complete and a FINI code is encountered the program tests one of the console sense switches to determine if the source program requires the use of AUTOSPOT Phase III. Phase III is not required for programs which have no tool offset or arc and slope -- that is for machining programs. If the sense switch is on (indicating contouring) the program branches to a routine which loads core with Phase III of the general processor. If, however, the switch is off (indicating no contouring) the program tests another of the sense switches to determine if the program requires post processing. If post processing is required the appropriate post processing phase is loaded into core from its location on the disc. If post processing is not required then a punch tape or a card deck containing the output of Phase II is produced.

If Phase III were required, it would be loaded into core as I mentioned previously. Processing would proceed in the usual manner with the image of each resulting card being loaded in the disc intermediate output area.

Upon completion of Phase III, indicated by a FINI code, one of the console sense switches is interrogated to determine if post processing is required; if not, a card deck containing Phase III output is produced. However, if a post

processor is required then the appropriate phase will be called from the disc into core.

The post processors applicable to our particular machine tools have been adapted to run in a manner very similar to the one just outlined.

These post processors are:

- (1) Computer Routines for The Milwaukee Matic Solid State Controlled Machining Center, User Group Library, No. 10.4.004, and
- (2) MATIC, a proprietary pseudo-language developed by Kearney and Trecker Corporation for its Bendix Dynapath controlled "Milwaukee Matics"

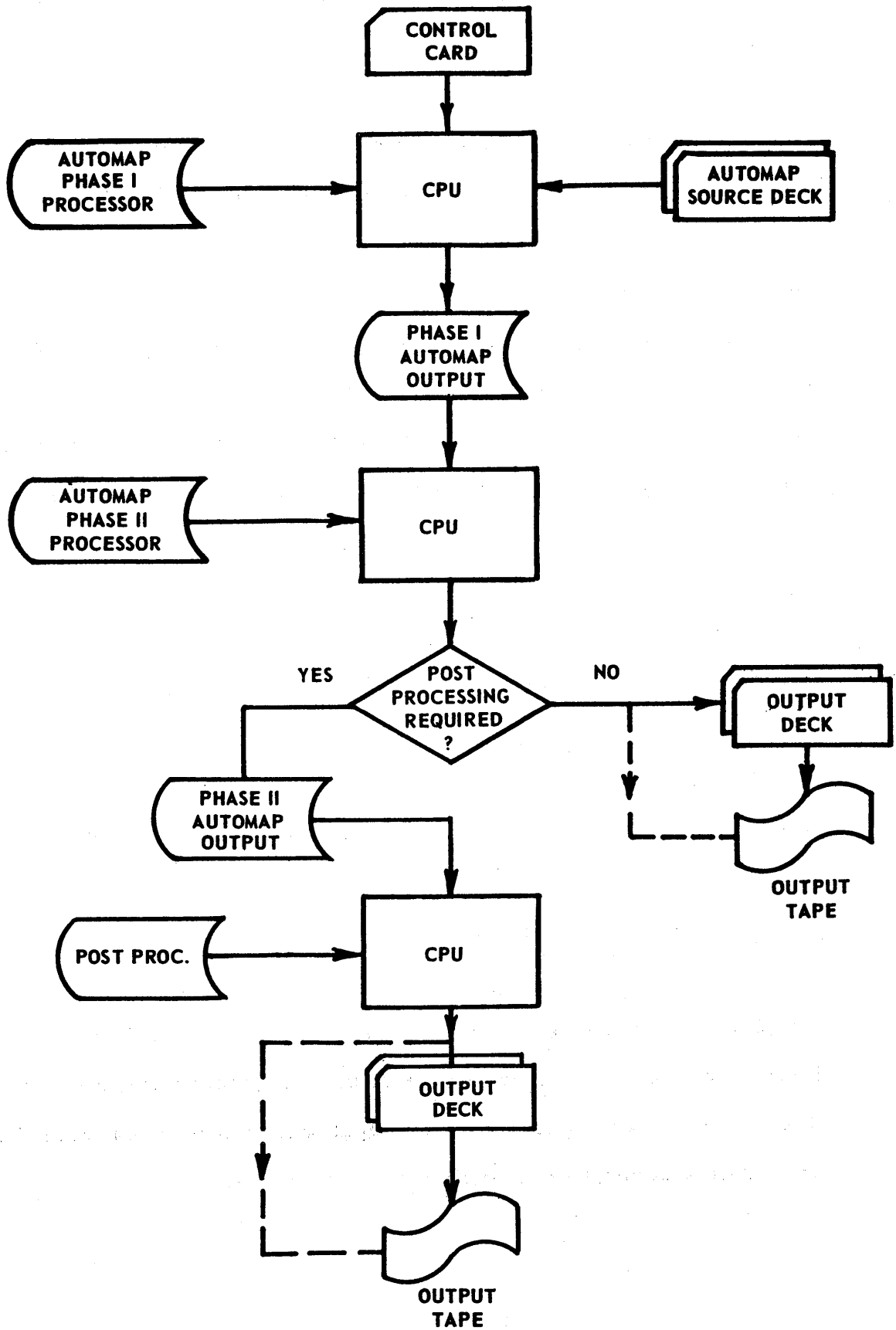
In general this is now the system operates. Let us now take a brief glance at the advantages and limitations of it.

#### ADVANTAGES AND LIMITATIONS

The principal advantages of this system as we see them are as follows:

(Slide #1A)

- (1) **Simplicity of operation:** One control card calls Phase I of AUTOMAP from disc into core -- the source deck is processed and the intermediate deck is loaded on the disc. This output will then be read from the disc after Phase II has been transferred from disc to core.



- (2) Core Clearing: Manual clearing of core to zeros or flag zeros is not required.
- (3) Card Handling: The handling of intermediate decks is eliminated thus reducing operator time and the hazard of dropped or mixed up decks.
- (4) Speed of Operation: The reading of lengthy processor decks with the 1622 is eliminated, saving computer hours and operator hours. The elimination of intermediate output decks speeds up the operation from both operator time and machine time standpoints.
- (5) Card Savings: If the company utilizing this technique does a large volume of numerical control program processing, the savings from decreased card consumption may be a significant percentage of the 1311 rental cost.
- (6) Disc READ/WRITE checks may be written into the program with as many re-reads or re-writes as desired.

At the present time we visualize that there may be two limitations or problems with this system.

- (1) Program Modifications: Program modifications which require additional core will not be possible in some phases due to the fact that the additional instructions required to implement this system cause core to be completely filled. In addition, program modifications which are desired, either as a result of program modifications published by the Users Group Library or by individual preference, must be incorporated into the deck and loaded on the disc, much in the same manner as MONITOR I patches.

- (2) Debugging Source Programs: Another problem exists if the individual users desire an intermediate deck for debugging -- let's say to find the location of a failure relative to some known point in the source program. It is often desirable to have this output to determine the last successful command, thereby locating the bug. At the present time, we utilize a debugging program in the form of a series of instructions, which will dump, from the disc, images of a number of cards before and after the location of the failure.

The seriousness of this limitation depends primarily on the method of processing utilized within the users organization. It is doubtful if it would be a serious limitation in any case.

### III SAVINGS

Having touched on the advantages of the system in general, let us return to the specific area of cost and time savings.

One might reasonably expect savings in three areas:

- (1) Savings in elapsed time from receipt, by the operator, of a source deck to production of final tape deck. This time saving is realized by the elimination of card handling between phases, core clearing, and punch clearing operations.

Our experience indicates that, normally, two hours actual processing time requires approximately three hours elapsed time. An actual test run with an average AUTOMAP program with three phases of post



processing required 60 minutes utilizing the card system exclusively but was reduced to 37.5 minutes utilizing this system.

A typical AUTOSPOT program with 6 phases of post processing ran 70.5 minutes utilizing the card system exclusively and 35.2 minutes utilizing the 1311.

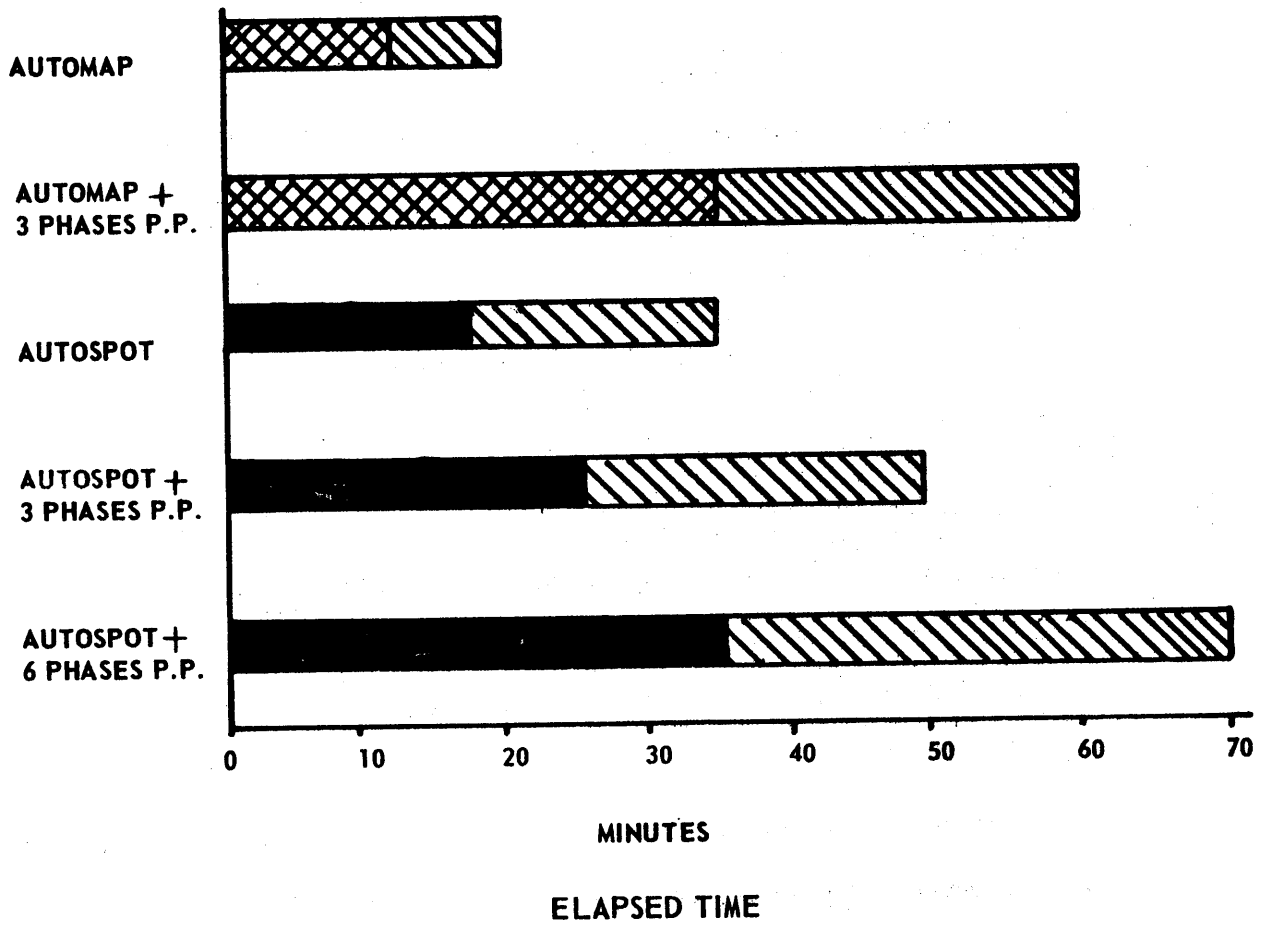
The elapsed time comparison for various AUTOSPOT and AUTOMAP plus post processor configurations are shown graphically on the next slide.

(Slide #2A)

- (2) The second area of saving is in computer meter time. This saving is realized, primarily, as a result of the elimination of the 1622 as an input/output device, in all cases except the reading of source deck and punching of tape deck. (We do not utilize the on-line tape punch).

In a test utilizing what we consider an average AUTOMAP source deck plus 3 phases of post processing the meter time utilizing card I/O was 46.8 minutes and utilizing the 1311 the time was 35.7 minutes.

A similar test with an average AUTOSPOT program with six phases of post processing indicated the running time without 1311 as 56.8 minutes and with this system as 38.8 minutes.



A bar chart showing the comparison between the two systems for various configurations of processor plus post processor is shown on the next slide.

(Slide #3A)

- (3) The third area of saving is in decreased card consumption. Obviously, this saving is realized as a result of the elimination of punched out intermediate decks.

The average AUTOMAP test deck which was utilized to determine the time comparisons previously mentioned, utilized 1600 cards with only card I/O and 450 with the help of the 1311.

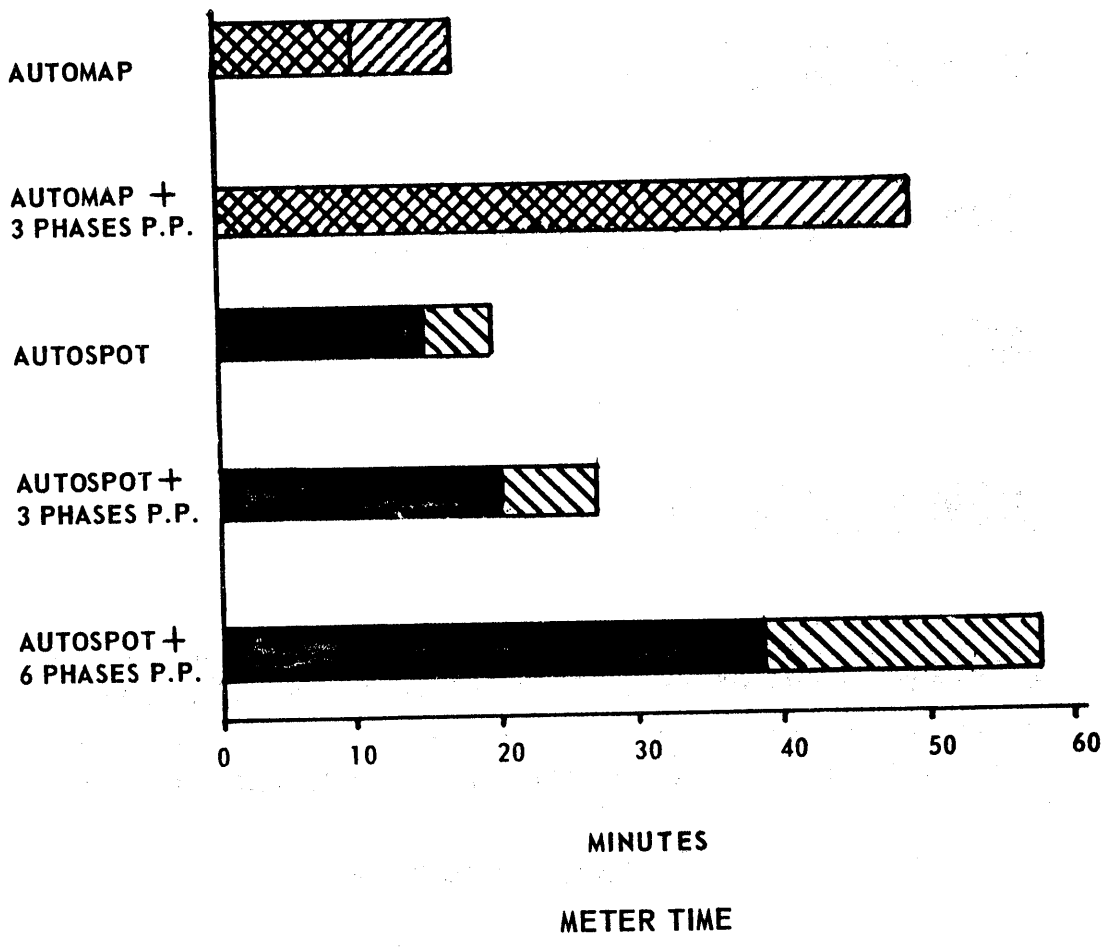
A similar test with AUTOSPOT utilized 1850 cards with card I/O and 580 cards with the 1311 as intermediate I/O medium.

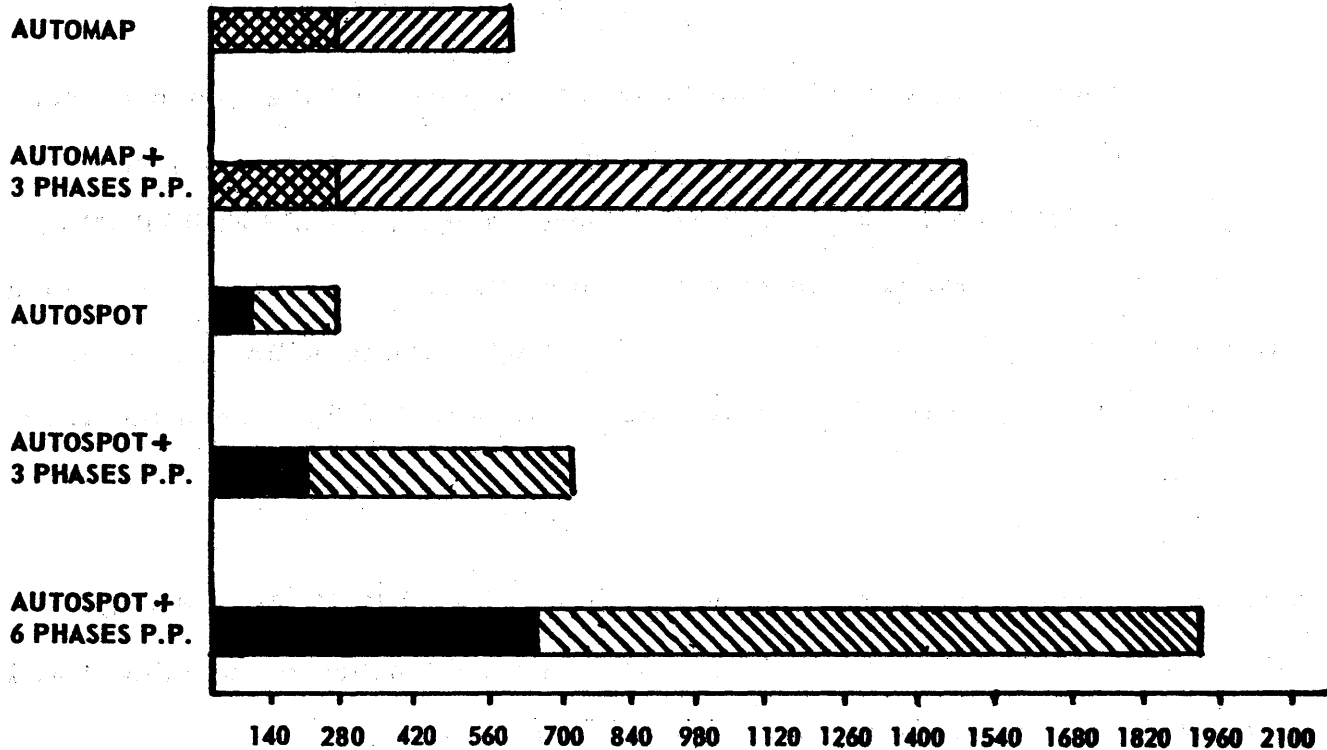
The next slide illustrates the card consumption for various configurations of processor plus post processor.

(Slide #4A)

## VI. OTHER APPLICATIONS

The technique which we have been discussing in general and which will be covered in greater detail later on, is by no means limited to numerical control program processing.





**CARD CONSUMPTION**

We have implemented it for this particular use because a considerable bulk of our processing is in the numerical control area.

The same technique has been utilized successfully for processing machine language programs compiled under "FOR TRAN with FORMAT" and SPS and even for linking two or more such programs together.

This application is somewhat troublesome from a programming standpoint but we feel that there exist areas wherein it would be useful. For example, companies which have already compiled programs under the old FOR TRAN or SPS systems might find it inconvenient or impractical to re-write and re-compile these programs under MONITOR. Thus they might wish to utilize a method similar to this to process these programs, particularly if they are of the daily run type.

The point is this: a method by which one might bypass MONITOR and still take advantage of the disc does exist, if the time and trouble for utilizing it will yield sufficient savings in the particular case in question.

We will be happy to answer questions regarding this technique at the end of our formal discussion and to discuss the details of it during the workshop which will be held this afternoon.

At this time, I would like to present Mr. Ernie Moore, Systems Engineer from Huntsville's IBM Branch Office. Mr. Moore has done a considerable amount of programming in this effort and will discuss the system in greater detail.

For those of you who already have a 1311 Disk System, I hope to show you a procedure for placing a machine language program on the disk. For those who have not yet obtained a disk system, I hope to show you how the Monitor System can be supplemented in the case of established programs which have heavy usage and are completely contained in one core image. To date, we have modified about fifteen programs to operate from the disk. Some of these programs also write data on the disk and then call the next program or phase which then uses this data written on the disk to continue the processing, as in the case of post-processors. We have modified both FORTRAN and SPS Programs. The procedure which I am going to explain is the result of our experience on these programs.

Our approach was intentionally the simplest one possible. It was to load the established program into core and then to transfer the complete core image, with minor modifications, onto the disk. When we wish to operate the program, we simply call the complete core image back into core using a single call card containing only four instructions. By this step, we have eliminated the handling of the program decks and the time to read cards into core. We are essentially using the disk as a storage medium for our card decks.

As we became more familiar with the disk, the next logical step was to have the first program in a series of programs automatically call in the second program when the first program was finished. This series of each calling the next could go on indefinitely -- the

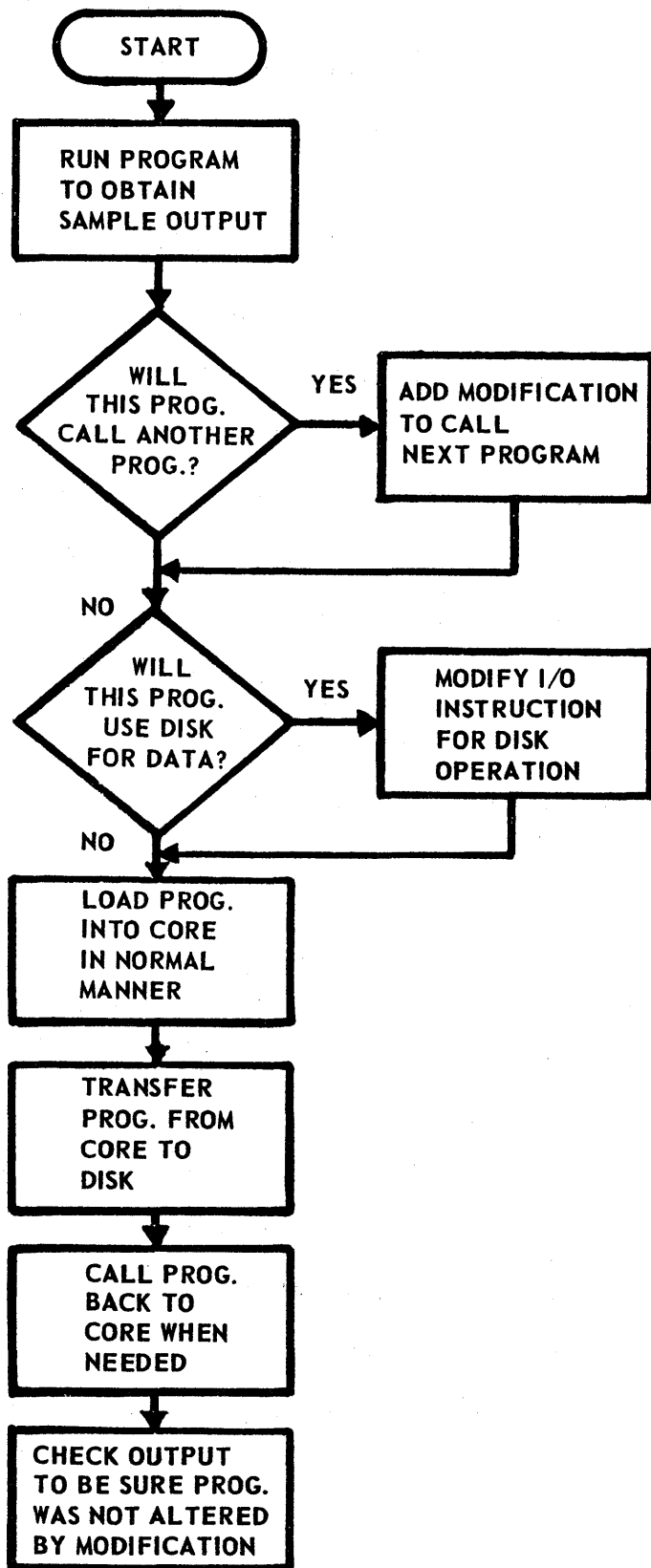
only limit being the limit of 100 core images on one disk pac. Our largest string to date is ten programs.

While we were still feeling pretty good about our first two achievements, we tackled the final step -- that of modifying the first program to place its intermediate output on the disk instead of punching cards, then modifying the second program in the string to read this intermediate data from the disk as input. The same program in turn was to write its output on the disk for the third program and so forth, on to completion. Here we have almost completely eliminated cards and their associated problems, and we have also eliminated the operator, once he has called the program in and loaded the initial data.

Slide 1 illustrates the sequence which we have just gone through. These are the steps you could take in modifying your program. You will note that before modifying the program, sample output is obtained for use later as a check to see that our modifications have not changed the logic of the original program. In the first decision block in Slide 1, we ask, "Will it be necessary for this program to call the next program in a string?" If the answer is "yes", then the program must be modified by adding a three or four instruction routine at the end. The second decision block indicates that the program must also be modified if it will use the disk instead of the cards for data storage. After the necessary changes have been made, the program is loaded into core using its own load program. The complete core image is then transferred to the disk and is available



# MODIFYING YOUR PROGRAM FOR DISK OPERATION



for call at any time it is needed. It is suggested that the output of the test run be compared with the original test data to point out any erroneous changes which have altered the logic of the program.

We will now go back and examine the individual blocks of this flowchart (Slide 1) in more detail. If it were decided that our program, when finished, should call in a second phase or program we could simply replace the final halt with the two instructions needed to call the next program from the disk. In Slide 2 we have replaced the halt at 10588 with an Op Code 34 followed by an Op Code 36. These will transfer the next program from the disk to core. This will work fine if the first instruction of the new program is located at 10612. This is because the instruction at 10612 will be the next instruction executed. From this, we can conclude that the Op Code 34 and 36 instructions to call the second program must be located just in front of the address at which the first instruction in the second program will be loaded. This is illustrated in Slide 2 where the Op Code 26 at location 10612 would be the first instruction executed in the second program.

To further illustrate this point, we will take the case of an SPS Program which has a halt at location 0 followed by a branch, in Loc 12, to the origin or beginning of the program. Here we have replaced the halt at 10588 in Slide 3 with a 49, or branch, to 19976. At location 19976, we have placed our Op Code 34 and 36. The next instruction executed after the Op Code 36 will be the instruction brought out from location 0 which is normally the first instruction in an SPS Program.

MODIFICATION TO CALL NEXT PROGRAM

CORE LOC.	OP CODE	P & Q ADDRESS	
10588	48	00000 00000	FINAL INSTRUCTION IN UNMODIFIED PROGRAM
10588	34	19962 00701	FIRST MODIFIED INSTRUCTION
10600	36	19962 00702	
10612	26		

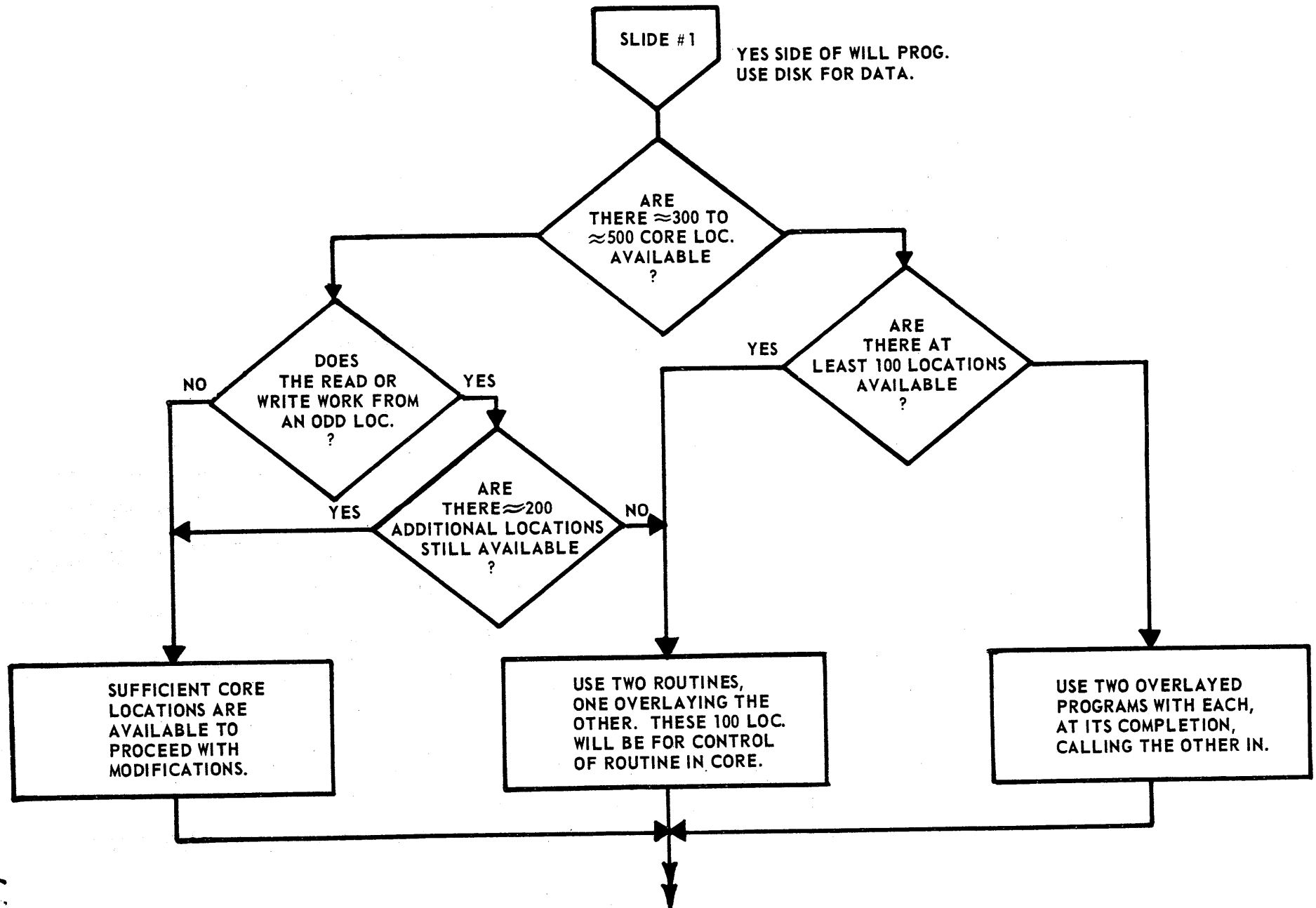
MODIFICATION TO CALL NEXT SPS PROGRAM

CORE LOC.	OP CODE	P & Q ADDRESS
10576 10588	LAST 49	INSTRUCTION IN PROG. 19976
19962 19976 19988	1000 34 36	00200 00000 19962 00701 19962 00702
00000 00012	41 49	00000 00000 01700

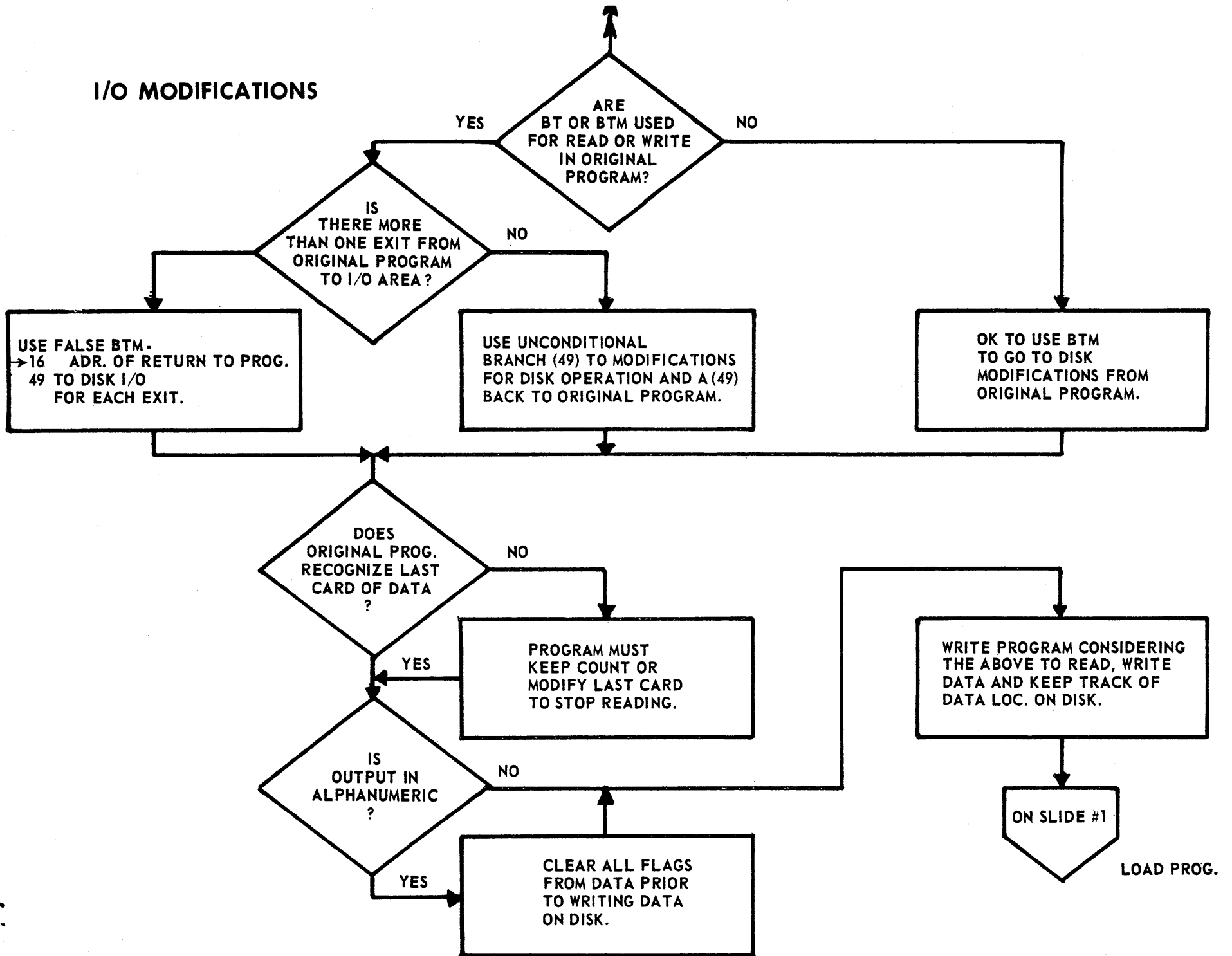
Now that the new program is in core, this is also the first instruction of our second phase or second program. This takes advantage of the wrap around feature which allows us to execute the instruction at location 0, after the instruction located at 19988 is executed. This also takes advantage of the fact that very seldom are instructions located in the upper end of core in the 19900 area. This allows us to use this area.

Let us now examine the block labeled "Modify I/O Instructions for Disk". Slide 4 is an enlargement of this block. The first decision block asks, "Are there 300 to 500 core locations available?" This figure, 300-500 is rather broad and varies depending on the number of read or write statements in the program which are to be modified. If we have this number of core positions available, then we go to the next decision block which asks, "Does the program read into or write from an area beginning with an odd address?" If this answer is "yes", then our modifications will require an additional 200 locations. These additional 200 locations are used to program the transfer of data so that the disk works out of or into an area beginning with an even address. If this is not done, the first or last digit of the data will be lost in a disk transfer. The transfer of data to core is made prior to a write disk or following a read disk instruction using a transmit record so that record marks and special characters would also be transferred; thus not altering the data in any respect. Once we find that sufficient core storage is available, we proceed to the writing of the modification program.

# I/O MODIFICATIONS



# I/O MODIFICATIONS



If, in the first block in Slide 4, we had found that sufficient core was not available, we would have gone out the no side to the next decision box which asks, "Are there at least 100 locations available?" These 100 locations will be necessary for a control program which will control a section of core storage calling either the I/O routine or the main program whenever either is needed. In other words, a portion of core will be shared by both the I/O program and the program which normally resides in this area. This way we can actually use a program which fills practically all of core and still have available the additional programming necessary to take care of reading and writing on the disk. There are very few cases where 100 locations cannot be found. In many cases, an output error message may be modified or abbreviated. The locations acquired in this manner may be used for control purposes. The locations from 0 - 80 in the product area may sometimes be used. By dumping the program out on the typewriter, there may be other areas which will become evident. In this way, sufficient area may be found to contain the instructions necessary to call the alternate program and control whichever program is in core at the present time. If we reach a situation where fewer than 100 cores are available, and if the logic of the program will allow, the best solution would be to have each of the two overlaid programs, at its completion, call its counterpart in on top of the existing program. In this way, you are alternating back and forth and each program, when executed, will automatically call the next.



The next decision block asks, "Are branch and transmit instructions used to go from the original program to its "read a card" or "punch a card" routine?" If the answer to this is "yes" we will not be able to use branch and transmit type instructions to branch our disk routines. And, in the case of this, we go out the left side of this block to the next block which asks, "Is there more than one place where the original program branches to its read or write routines?" The point here is: If branch and transmit instructions had been used by the original program to go into its routine which, after modification, we will be branching from, we may not use a branch and transmit again prior to reaching the branch back. Thus, in cases where the branch and transmit instruction is used by the original program to branch to its I/O routine, we must use a 49 type branch to branch to disk routine. But, if the original program had not used a branch and transmit instruction, we may, in turn, use branch and transmit instructions to get to and from the disk routine.

The next block says, "Does the program recognize the last card of data?" With card operation, when the last data card has been read, the card reader will stop, but the disk will continue to read sectors beyond the last data unless provisions are made to sense this last data. If the answer to this decision block is "no", then some provision must be made so that the program will not read the disk completely never knowing when it has finished the last card. This may be accomplished in two ways:

1. A count can be maintained of the number of cards stored on the disk when the data was originally written on the disk. This count can be checked as the data is read back for the last position.
2. Or, the position following the last data segment on the disk could be loaded with a special indicator which the next, or following, program will recognize as the last data area.

In the case of AUTOSPOT AND AUTOMAP, the program already made provisions for the last card by placing a fini card at the end of the data as it is written. This fini card contained a 99 and was recognized by the following program as terminating the data.

From here, we go to the next decision block which checks for input in alphanumeric form. Again, in this block we run into a uniqueness of the disk which in some cases would be an advantage, but which we must watch for. When we read the disk or write disk in alpha, flags are transferred with the data. This is contrary to reading or writing on cards. Thus, we must make provisions to remove the flags left in the data. The programs with which we are working assumed that there would be no flags in the input data and went on to set flags in the input area which were later used for data transfer. The extra flags left by the disk can cause serious errors if allowed to remain. To correct this, we used a clear flag instruction to clear all flags from the 80 positions of data prior to writing on the disk. In this way, no flags were read back from the disk into the input area.

Coming down through the flowchart (Slide 1), we have now reached the point which says load program into core in normal manner. The modifications which we have described up until now may be inserted into the original program deck in two ways:

1. We can modify the original deck prior to loading the program into core by repunching the necessary cards.
2. We can load the program in with no modifications and then write a "trailer program" which will load the modifications on top of the normal program.

Either method is satisfactory.

Now it becomes necessary to transfer the complete modified program onto the disk for recall at a later time. Again, there are a number of ways in which this can be done.

1. This can be done by placing the two necessary disk instructions into the input area. After we had loaded the program with the modifications, we would branch to these instructions which in turn would load the program on the disk. When the program is read back from the disk these two instructions (34 and 38) would still be in the input area but the assumption is that the first data read into this area would be read over these instructions and they would have no affect on the program.
2. If we had used the trailer program to load in our modifications after the initial program had been already loaded into core, we would have included these two

instructions. Again, either way is satisfactory. The main point is to get the program on the disk with the modifications.

Now all that remains is to have the program called in from the disk. Here, we may use the same philosophy which we had used when we had one program calling the following program. The main thing to remember here is that the Op Code 36 instruction must be located just in front of the first instruction to be executed in our next program. In SPS programs, we used the call routine illustrated on Slide 5. We have here an Op Code 41 followed by an Op Code 34, 16 and 49. The program is read into location 0. The Op Code 41 will do nothing. We go to the 34 which will seek the disk address which we specified in the control word located at 44. The 16 transmits immediately the 36 to location 0 and 1. We then branch back to 0 and execute this instruction which will now be a 36 or "read a disk". The new program will be read in and the next instruction executed, after the instruction located at location 0, will be the instruction located at 12 which, in the case of SPS program, will be a branch to the origin of the program.

In the case of non-SPS programs where the branch is not located in position 12, we may use the program similar to the one in the second part of Slide 5. Again, this program is read into location 0: the first instruction is a 34, "seek the disk". The second instruction, an Op Code 26, will transfer the Op Code 36, instruction, located at 0046, to a location just in front of where the next program will start after it is read in. The third instruction Op Code 49 will branch to and execute

TWO CALL PROGRAMS TO CALL AND  
START PROGRAM ON DISK

CORE LOC.	OP CODE	CONTENTS OF P & Q ADDRESS
00000	41	00044      00702
00012	34	00044      00701
00024	16	00001      00036
00036	49	00000
00044		105000 200 00000 (DISK CONTROL WORD)
00000	34	00032      00701
00012	26	START - 1   00057
00024	49	START - 12
00032		105000 200 00000 (DISK CONTROL WORD)
00046	36	00032      00702
00058		

the Op Code 36 instruction. The Op Code 36 instruction will be executed reading in the new program and the following instruction which will be the first instruction in the new program will be the next one executed.

The final block on Slide 1 points out the advisability of checking the finished program by comparing its output with the output from the program prior to any modifications. In this way, we can be relatively sure we've not altered the main philosophy of the program in any way.

I have included in the appendix a typical set of modifications for your reference. I might add one precaution -- DO NOT let any of your modified programs get on the Monitor Disk. Probably Monitor would have to be reloaded and your program most likely would not run any way.

In concluding, let me say that I hope I have brought to your attention an area of disk operation which has received very little publicity in the past. You must realize that there are very definite limitations to the use of the disk with programs modified in this manner. The program must be in a complete core image and if programs are linked together with each calling the next, the sequence is restricted and there can be little deviation without rewriting the modifications. But, for programs which will run in the same sequence, or for a single program that is run very often, a considerable savings can result. The resulting program is fast, economical and easy to operate.

I realize we have covered some rather technical material here in a rather short time. Therefore, I invite your questions either now or this afternoon during our workshop when we hope to sit down with you

and help you modify your post-processors or any other programs which you have to modify. Please bring a copy of your program listing and find out the last location that your program uses in core. Anytime in the near future that I may be of assistance, please feel free to contact me through the Huntsville Branch Office. I hope to see many of you in our workshop this afternoon.

Now, are there any questions?

II

Appendix 1

OP Code Reference Table

and

Disc Word Explanation



CODE	MNEMONIC	TYPE OPERATION
11	AM	Add Immediate
12	SM	Subtract Immediate
13	MM	Multiply Immediate
14	CM	Compare Immediate
15	TDM	Transmit Digit Immediate
16	TFM	Transmit Field Immediate
17	BTM	Branch and Transmit Immediate
21	A	Add
22	S	Subtract
23	M	Multiply
24	C	Compare
25	TD	Transmit Digit
26	TF	Transmit Field
27	BT	Branch and Transmit
31	TR	Transmit Record
32	SF	Set Flag
33	CF	Clear Flag
34	SK	Seek (Q = x07x1)
34	K	Control
35	DN	Dump Numerically
* 36	RN	Read Numerically
37	RA	Read Alphamerically
* 38	WN	Write Numerically
39	WA	Write Alphamerically

CODE	MNEMONIC	TYPE OPERATION
41	NOP	No Operation
42	BB	Branch Back
43	BD	Branch Digit
44	BNF	Branch No Flag
45	BNR	Branch No Record Mark
46	BI	Branch Indicator
47	BNI	Branch No Indicator
48	H	Halt
49	B	Branch
55	BNG	Branch No Group Mark

\*Read-Write disk modifiers on next page.

## DISK CONTROL FIELD

In order to read from or write on the disk there are four things that must be known. These are;

- (1) The disk drive number if more than one drive is attached to the system.
- (2) The five position disk sector address.
- (3) The number of sectors to be written or read.
- (4) The starting core location.

The disk control field incorporates all four of the above items into a 14 position field. Thus:

F $\emptyset$	F <sub>1</sub> , F <sub>2</sub> , F <sub>3</sub> , F <sub>4</sub> , F <sub>5</sub>	S <sub>6</sub> , S <sub>7</sub> , S <sub>8</sub>	M <sub>9</sub> , M <sub>10</sub> , M <sub>11</sub> , M <sub>12</sub> , M <sub>13</sub>
---------------	--	--	--

The disk drive number is located in F  $\emptyset$ . This drive code number varies with the number of drives attached to the system. For drive  $\emptyset$  a 1 is used. For drive 1 a 3 is used.

A sector on the disk is equal to 100 positions of core storage. There are 20,000 sectors on each disk. These sectors are numbered sequentially from 00000 - 19999. The disk control field F1 - F5 contains the sector address. This sector address determines where, on the disk, the write or read will start.

Next is the number of sectors to be read or written. This is located in S6 - S8. The maximum number of sectors that can be read or written is 200 and the minimum number is 001. The method for reading or writing fewer than 100 core locations is explained on the next page in "Read-Write Disk Modifiers".

M9 - M13 contains the core location of the leftmost position of the data transferred to or from the disk. This core location must be an even number.

In a seek, read, or write disk instruction the "P" address is the core location of the leftmost position of the disk control field. This leftmost position must be in an even location.

The "Q" address of the disk instructions contains  $\emptyset 7$  in Q8 and Q9 and a modifier in Q11. The modifier in all seek instructions is a 1. The modifier in read-write instructions is explained in "Read-Write Disk Modifiers".

## READ-WRITE DISK MODIFIERS

All read-write disk instructions must have a "Q" address of "xØ7xM" where M is the modifier. The modifier determines whether or not a group mark ( $\text{⌘}$ ) will have any effect on the data being transferred.

The write disk instruction (38) with a modifier of Ø will be determined after the first group mark encountered in core has been transferred to the disk. If no group mark is encountered the instruction will be terminated when the sector count has been decremented to ØØØ.

The read disk instruction (36) with a modifier of Ø will be terminated after the first group mark encountered on the disk has been transferred into core or, if no group mark is encountered, when the sector has been decremented to ØØØ.

The read or write disk instruction with a modifier of 2 will treat the group mark as data and transfer data until the sector count has been decremented to ØØØ.

III

Appendix 2

Machine Language Modifications  
to AUTOSPOT, AUTOMAP and  
Milwaukee Matic Post- Processor

Note:

The post-processor included here is Users  
Group Library number 10.4.004 - the  
"Computer Routines for the Milwaukee  
Matic Solid State Controlled Machining Centers".

## AUTOMAP PHASE I

Statement number 1 is a five position field for the indirect address which shows from where to start the transmission of the record (statement #2). This is done because there are two write statements in the main program and each writes from a different location. As each of the two locations are odd numbered, they must be moved to an even location and, since the only locations left are 00000 - 00080, these will be used. Statements #2 and #3 move the data from the odd numbered program output area into even numbered locations. Statement #4 writes the output data on the disk.

Statements 5, 6, and 7 check indicators; address check, wrong-length record/read back check and write check respectively and, if either the address check or write check indicator is on, a branch to a "seek" instruction (statement #10) and then a branch (statement #11) to the write instruction is made. If the WLR/RBC indicator is on a branch to the next instruction is made simply to turn off the WLR/RBC console light. This is done due to the fact that this indicator is turned on each time a record with length unequal to 100 character multiples is read or written.

Statement 8 adds one (1) to the sector address. Return to the main program is accomplished by a branch back (statement #9).

Statement #12 is the write output data disk word.

Statements 14 - 19 type the message "FINI" to indicate the end of phase I and to set up the calling of the next program.

Statements 20 - 25 are changes to the main program. 20 and 22 clear the disk output area. 21 and 23 branch to the write disk routine and transmit the starting address into the area reserved by statement #1. Statement 25 insures a group mark after the 80th position to terminate the write disk instruction.

Statements 26 - 31 dump the program on the disk in a core image.

Statements 32 - 39 load the modifications, read the first program loader card and branch to continue loading the main program.



AUTOMAP PH1  
WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
1	19802	00	0000	
2	19808	31	00000	19807
3	19820	26	00079	15390
4	19832	38	19914	00700
5	19844	46	19894	03600
6	19856	46	19868	03700
7	19868	46	19894	00700
8	19880	11	19919	00001
9	19892	42		
10	19894	34	19914	00701
11	19906	49	19832	0
12	19914	10	00000	01000
13	19926	00		
14	19928	10	42002	00000
15	19940	00	46495	5490 <sup>7</sup>
16	19952	39	19943	00100
17	19964	34	19928	00701
18	19976	48	00000	00000
19	19988	36	19928	00702

AUTOMAP PHI  
 LOADER AND CHANGES TO MAIN PROGRAM

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
20	02504	31	00000	15982
21	02516	17	19808	16063
22	11086	31	00000	15982
23	11098	17	19808	16065
24	11182	49	19952	00000
25	16062	≠		
26	15402	34	15440	00701
27	15414	16	00004	41000
28	15426	38	15440	00702
29	15438	48		
30	15440	10	40002	00000
31	15452	00		
32	00000	36	00080	00500
33	00012	36	15402	00500
34	00024	36	19802	00500
35	00036	36	19882	00500
36	00048	36	19962	00500
37	00060	49	00080	
38	00080	36	00000	00500
39	00092	49	00000	

PHASE I SPECIFICATION SECTION OF  
AUTOSPOT - MODIFICATIONS FOR DISK OPERATION

LOAD PROGRAM

00000	36	19522	00500	
00012	36	19680	00500	
00024	36	19840	00500	Load the modifications into core
00036	36	19600	00500	
00048	36	19760	00500	
00060	36	19920	00500	
00072	49	19626	9	
19522	31	19648	05819	Save data for next phase
19534	49	19976	05914	Branch to "call next program"
19546	105000	20000000		Disk control for this phase

FALSE BRANCH & TRANSMIT

19560	105200	14904948		Disk control next phase
19574	15	19969	00009	Change 42 to a 49
19586	16	19821	10701	Transfer "write from" address
19598	49	19822	0	Go to entry of program
19606	15	19969	00002	Change 49 to 42
19618	49	10660	0	Return to program

MODIFICATIONS TO PROGRAM

19526	36	00000	00500	Reset location 0 to 80
19538	26	10654	19785	Modify unit instruction to branch to unit on disk routine
19550	26	07755	19759	
19562	26	09039	19771	
19574	26	06368	19778	Modify end of program
19586	15	05911	0000 <del>7</del>	Used to save data for machining section
19698	32	07751	00000	
19710	32	09035	00000	
19722	34	19546	00701	Write program on disk
19734	38	19546	00702	
19746	48			
19748	27	19822	02753	
19760	17	19822	10877	Modified instructions to be inserted in program
19772	49	19976		
19779	49	19574		
19786	11	19821	00001	
19798	11	19816	00001	Increase address and transmit digit
19810	25	19626	00000	
19822	14	19816	19785	
19834	47	19878	01200	Has all data been moved?

PHASE I SPECIFICATION SECTION OF  
AUTOSPOT - MODIFICATIONS FOR DISK OPERATION

19846	16	19840	19878	Prepares this section for return pass
19858	16	19816	19626	
19870	49	19912	0	
19878	16	19840	19786	1st pass set-up
19890	49	19810	0	
19898	10000000	19626		
19912	38	19898	00702	Write data on disk
19924	47	19956	03600	
19936	34	19898	00701	Rewrite if we had an address check
19948	49	19912	0	
19956	11	19903	00001	Increase sector address by one
19968	42	19606	0	Return to Program
19976	34	19560	00701	
19988	36	19560	00702	Call next program
00000				

## PHASE I MACHINE SECTION

Same as Phase I specification with these exceptions

19522	41	00000	00000	
19534	49	19976	04948	Branch to end of program
19546		10520014919746		Disk address for this phase
19560		10540020000000		Disk address for next phase
19626	41	00000	00500	
19638	26	15929	19711	Modify write instruction
19650	26	07595	19719	Modify end of program
19662	32	15913	00000	
19674	34	19546	00701	
19686	38	19546	00702	Load this program on disk
19698	48			
19700	49	19574	00000	
19712	49	19534		
19720				
19746	31	05819	19648	To transmit record left by specification section of Phase I
19758	49	04984	0	

## MODIFICATION TO PHASE 2 OF AUTOSPOT TO

### RUN FROM DISK

#### PATCH PROGRAM

19466	26	13225	12033	
19478	11	13225	00010	
19490	26	13225	12055	
19502	11	13225	00004	Program patch that was in the way moved to here
19514	49	11226	0	
19522	34	19560	00701	End of this phase - seek for next phase
19534	49	19976	06850	
19546	105400	20000000		Disk control word this section
19560	1053500	1001900		Disk control word next section
19574	bbbbbbbb			

#### MODIFICATION TO ORIGINAL PROGRAM

19584	36	00000	00500	Read in last cards of modification
19596	36	19680	00500	
19608	26	01762	19712	Modify "read a card" instruction
19620	26	03219	19705	Modify error routine
19632	41	00000	00000	No Op
19644	26	11780	19720	Modify "write a card" instruction
19656	26	12024	19727	Modify end to call next program
19668	34	19546	00701	
19680	38	19546	00702	Write these modifications all on the disk
19692	48			
19694	46	19466	01200	
19706	49	197420		
19714	49	19912	4919522	

#### READ DISK DATA

19728	10000000	119600		
19742	36	19728	00702	Read a card from disk & seek if necessary then go back and read again
19754	47	19786	03600	
19766	34	19728	00701	
19778	49	19742	0	
19786	11	19733	00001	Increase sector address by one

#### TRANSFER FIELD JUST READ TO EVEN LOCATION

19798	25	13056	19600	Transfer data to area beginning with an even address
19810	14	19809	19679	
19822	46	19866	01200	Check for last transfer
19834	11	19809	00001	Increase disc location by one
19846	11	19804	00001	
19858	49	19798	0	Repeat

MODIFICATIONS TO PHASE 2 of AUTOSPOT

TO RUN FROM DISK

TRANSFER FIELD JUST READ TO EVEN LOCATION

19866	16	19809	19600	Housekeep because all data is transferred now
19878	16	19804	13056	
19890	49	01768	0	Return to program

WRITE DATA ON DISK

19898	10200000	113136		
19912	38	19898	00702	Write data on disk
19924	47	19956	03600	If address check - seek first then write data
19936	34	19898	00701	
19948	49	19912	0	
19956	11	19903	00001	Increase sector address by one
19968	49	11786	0	Return to program
19976	36	19560	00702	End of this phase - read in control program for next phase
19988	49	02318	00000	

Load program to load in modifications into core

00000	36	19466	00500	Load modifications into core and branch to the first modification
00012	36	19546	00500	
00024	36	19626	00500	
00036	36	19760	00500	
00048	36	19840	00500	
00060	36	19920	00500	
00072	49	19584	0	

## MODIFICATIONS TO AUTOSPOT PHASE 3 FOR DISK OPERATION

### READ DATA FROM THE DISK

18312	36	01984	00702	Read data from disk
18324	47	18356	03600	
18336	34	01984	00701	If not correct cylinder seek and go back to read
18348	49	18312	0	
18356	11	01989	00001	Add one to sector address
18368	25	13411	18626	Transfer data to data read in area which start with an odd address. Use transmit digit 80 times.
18380	14	18379	18705	
18392	46	18436	01200	
18404	11	18374	00001	
18416	11	18379	00001	
18428	49	18368	0	
18436	16	18374	13411	Housekeep transfer data routine and return to program
18448	16	18379	18626	
18460	49	13334	00000	

### CALL OVERLAYED ROUTINE "FLOAT" & CONTROL WORDS

18472	16	19964	10292	If program branches to FSIN or FCOS set at to call FLOAT and go to 18786
18484	49	18786	00000	
18496	16	19964	10376	
18508	49	18786	00000	
18520		00000000000000		
18534	41	00000	01700	Origin of program 01700
18546		10560020000000		Disk address this program
18560		1053500100190000		Disk address of next program

### ROUTINE TO MODIFY ORIGINAL PHASE 3 PROGRAM BEFORE IT IS PLACED ON DISK ALSO 18626 to 18705 ARE USED FOR A TRANSFER OF DATA AREA

18576	36	00000	00500	Read in a card and branch to it
18588	49	00000	0	
18596	26	13801	00063	Modify write data instruction
18608	26	13329	00071	Modify read data instruction
18620	36	13298	00500	Read in another card
18632	36	13620	00500	Read Modification into 13620
18644	16	10322	19966	
18656	16	10406	02590	Modification to original program to call float
18668	16	12882	02570	
18680	36	13272	00500	Read modifications into 13272
18692	36	00000	00500	Read original information into 00000
18704	34	18546	00701	Write complete program disk



MODIFICATIONS TO AUTOSPOT PHASE 3 FOR DISK OPERATION

ROUTINE TO MODIFY ORIGINAL PHASE 3 PROGRAM BEFORE IT IS PLACED ON DISK  
ALSO 18626 TO 18705 ARE USED FOR A TRANSFER OF DATA AREA

18716 38 18546 00702  
18728 34 13280 00701 Write "read/write data on disk" routine on disk  
18740 38 13280 00702  
18752 48 20 - 0's

IF PROGRAM NEEDS FATN CALL IN FLOAT ROUTINE

18774 16 19964 12852  
18786 34 00000 00701 Seek float and go to 19946 which will read it in  
18798 49 19946 0

WRITE OUTPUT DATA ON DISK FOR NEXT PROGRAM

18806 16 18841 13847 Housekeep transmit digit  
18818 16 18836 18626  
18830 25 18626 13847 Transmit digit 80 times  
18842 14 18836 18705 Have we transmitted digit 80 times?  
18854 46 18912 01200  
18862 11 18836 00001 Increase count or transmit digit and write data  
18878 11 18841 00001  
18890 49 18830 0  
18902 0000000000000000  
18912 38 19986 00702 Write data on disk  
18924 47 18956 03600  
18936 34 19986 00701 If at wrong cylinder seek and rewrite  
18948 49 18912 0  
18956 11 19991 00001 Increase sector count by one and return to program  
18968 49 13806 0  
18980

INSTRUCTION FOR LOCATION 0 TO 80

00000 49 01700 000 Branch to origin of program  
00010 44 18312 18966  
00022 36 13280 00702 If the read/write disk routine is in core go to R/W if not call it in  
00034 49 18312 0  
00042 44 18806 18966  
00054 36 13280 00702  
00066 49 18806 00042M3  
00080

INSTRUCTION FOR LOCATION 02528 TO 02610

02528 39 02559 00100  
02540 26 17043 16793 Error message - no tool card: "NO TC"  
02552 49 02634 0  
02560 55 56634 30#

## MODIFICATIONS TO AUTOSPOT PHASE 3 FOR DISK OPERATION

02570	44	18774	18966	
02582	49	18668	0	Check to see is "float" routine is in core - if not call it in
02590	44	18496	18966	
02602	49	17676	0	
02610				

### INSTRUCTIONS FOR LOCATION 19946 to 00000

19946	36	13628	00602	Call in float and branch to proper location
19958	49	00000	0	
19966	44	18472	18966	Check to see that float is in core if not prepare to call it in
19978	49	17708	0	
19986	10	200000	118626	Control word for read data from disk
00000				

# MODIFICATIONS TO AUTOSPOT PHASE 3 FOR DISK

## OPERATION

### Load cards for phase 3 modifications

00000	34	00056	00701	Load "float" on to disk
00012	38	00056	00702	
00024	36	18552	00500	Load in two modification cards
00036	36	18632	00500	
00048	49	18576	0	Call in next card
00056	10536700718312			Control word for float
00000	36	18312	00500	Load in six modification cards
00012	36	18392	00500	
00024	36	18472	00500	
00036	36	18712	00500	
00048	36	18792	00500	
00060	36	18872	00500	
00072	49	18576	0	Call in next card
00080				
00000	36	18930	00500	
00012	36	19892	00500	Load in 4 modification cards
00024	36	02530	00500	
00036	36	01936	00500	
00048	49	18596	0	Return to Modification program
00056	49	00042	0	Modification for write a card
00064	49	00010	0	Modification for read a card
00072				

## TRAILER CONTROL PROGRAM FOR AUTOSPOT

### DISK CONTROL WORDS

01900 X $\bar{X}$  XXXX2 0000000 Control word to call post processor  
01914 10560020000000 Disk control word to call phase 3  
01928 10200000110000 Disk control word to punch out put  
01942 10000000000000 Disk control word for last data written by last phase  
01956 10200010010000 Disk control word for write - X per data  
01970 10000010010000 Disk control word for read trace per data

### TRANSFER DATA FROM SECTOR 100000 TO 102000 SO POST PROCESSOR WILL FIND DATA

01984 26 01947 19991  
01996 34 01970 00701  
02008 36 01970 00702 Transfer data, 10,000 location at a time  
02020 34 01956 00701  
02032 38 01956 00702  
02044 11 01975 00100 Increase sector address for transfer by 100  
02056 11 01961 00100  
02068 24 01975 01947 If more data still - so back & transfer again  
02080 47 01996 01100

### CHECK FOR CARD OUTPUT

02092 34 00000 00102  
02104 39 02375 00100 Type out instructions  
02116 34 00000 00102  
02128 48 00000 00000  
02140 47 02224 00100 Check switch

### PUNCH OUTPUT ON CARDS FROM DISK

02152 34 01928 00701 Read from the disk one card at a time & punch it  
02164 36 01928 00702  
02176 38 10000 00400  
02188 11 01933 00001  
02200 24 01933 01947 If not finished, get next card  
02212 47 02152 01100

### CALL IN NEXT PROGRAM

02224 26 01927 01913 Enter here from phase 3  
02236 34 00000 00102 Enter here from phase 2  
02248 39 02441 00100 Set up switches for next program  
02260 48 00000 00000  
02272 34 01914 00701  
02284 26 19999 02317 Call in next program from disk and branch to its start  
02296 49 19988 0  
02304 36 01914 00702  
02316 00

TRAILER CONTROL PROGRAM FOR AUTOSPOT

CHECK SWITCH 3

02318	34	00000	00102	Check switch 3 for either phase 3 call or card output
02330	39	02645	00100	
02342	48	00000	00000	
02354	46	02236	00300	
02366	49	02092	0	

PRINT AREA

02374	62	66007	10056	TYPE OUT AREA:	SW 1 0
02386	55	00465	65900	NbFORb	
02398	43	41594	40056	CARDbO	
02410	64	63215	94562	UT/RES	
02422	45	63006	26341	ETbSTA	
02434	59	63076	24563	RT7SET	
02446	64	57006	26600	UPbSWb	
02458	46	56590	05545	FORbNE	
02470	67	63005	75956	XTbPRO	
02482	47	20594	56245	G-RESE	
02494	63	00626	34159	TbSTAR	
02506	63	07		T7	
02540	70	00			

LOAD CORE & DISK

02544	41	02220	00500	
02556	36	02300	00500	Call the remainder of the program into core and load the whole program disk with a correct halt at 02628
02568	36	02380	00500	
02580	36	02460	00500	
02592	36	02620	00500	
02604	34	02630	00701	
02616	36	02630	00702	
02628	48			
02630		10535001001900		
02644	62	66007	30056	SWb3bO
02656	55	00465	65900	NbFORb
02668	57	48007	32159	PHb31R
02680	45	62456	30062	ESETbS
02692	63	41596	307	TART

LOAD ROUTINE - TO LOAD PROGRAM INTO CORE

00000	36	01900	00500
00012	36	01980	00500
00024	36	02000	00500
00036	36	02140	00500
00048	36	02220	00500

TRAILER CONTROL PROGRAM FOR AUTOSPOT

LOAD ROUTINE - TO LOAD PROGRAM INTO CORE

00060 36 02540 00500

00072 49 02544 0

## AUTOMAP PHASE II

This phase is loaded on the disk in two sections. This is done because all core locations are taken and the "read in" area is defined as "DC" rather than "DS". Statements 1 - 14 load the first section and 15 - 23 load the second section.

Statements 24 - 31 are changes to the main program. Statement 24 branches to the read disk routine and statement 25 adds one (1) to the read sector address upon returning to the main program. Statement 26 transmits the field of numerical blanks to location  $\emptyset\emptyset\emptyset\emptyset$  rather than to the output area. Number 27 branches to the write disk routine. Number 28 changes the halt after the "END" typeout to a "no op". Statements 29 - 31 change record marks to group marks to insure termination of the write disk instruction after the transfer of  $8\emptyset$  characters.

Statements 32 - 42 are the read disk routine. The program branches to the read instruction (#33) and if an address check or write check occurs a branch to the seek instruction (#32) is made. Upon completion of the read operation, statement 37 checks for a "FINI" code. Upon finding a "FINI", statement 39 sets up a branch to end of program routine. Statement  $4\emptyset$  branches to the main program.

Statements 43 - 56 are the write disk instructions. The program branches to statement 44, which gives the option of either; (1) putting the output on the disk or (2) punching it in cards. Statements 45 and 46 move the output from the odd numbered core location to an even location. Statement 47 writes on the

disk. Statements 48, 49, and 50 check indicators and seeks (statement 43) if necessary. Statement 51 adds one (1) to the sector address and statement 52 returns to the main program. Statement 53 in the punch statement and statement 54 returns to the main program.

The "End of Job" message is contained in statements 57, 58, and 59. Statements 62, 63, 64 and 65 type "End of Job" and call the next program.



AUTOMAP PH2  
CORE TO DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
1	00000	36	19640	00500
2	00012	36	19720	00500
3	00024	36	00080	00500
4	00036	36	19900	00500
5	00048	49	00080	
6	00080	36	00000	00500
7	00092	49	00000	
R-W DISK				
CARDS				
8	19900	16	00004	41000
9	19912	16	00009	00000
10	19924	34	19950	00701
11	19936	38	19950	00702
12	19948	48		
13	19950	10	42001	98000
14	19962	00		
15	00000	36	15000	00500
16	00012	36	19800	00500
17	00024	36	19840	00500

AUTOMAP PH2  
CORE TO DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
18	00036	36	19920	00500
19		49	15000	
20	15000	34	15026	00701
21	15012	38	15026	00702
22	15024	48		
23	15026	10	43980	02198
		00		
FROM R-W				
DISK CARDS				
24	01714	49	19652	00000
25	01726	11	19749	00001
26	01854	31	00000	11372
27	01866	49	19770	00000
28	07872	41		
29	11290	<del>49</del>		
30	11371	<del>49</del>		
31	04603	<del>49</del>		
32	19640	34	19744	00701
33	19652	36	19744	00700
34	19664	46	19640	03600
35	19676	46	19688	03700

## AUTOMAP PH2

## CORE TO DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
36	19688	46	19640	00600
37	19700	14	11209	00004
38	19712	47	19736	01200
39	19724	16	01720	19952
40	19736	49	01726	0
41	19744	10	00000	01112
42	19756	10		
43	19758	34	19894	00700
44	19770	46	19874	00400
45	19782	31	00000	11291
46	19794	26	00079	11370
47	19806	38	19894	00700
48	19818	46	19758	03600
49	19830	46	19842	03700
50	19842	46	19758	00700
51	19854	11	19899	00001
52	19866	49	01878	0
53	19874	38	11291	00400
54	19886	49	01878	0
55	19894	10	20000	01000
56	19906	00		

AUTOMAP PH2  
CORE TO DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
57	19908	00	57487	20062
58	19920	63	41596	30043
59	19932	56	550 <del>7</del>	
60	19938	10	44000	72000
61	19950	00		
62	19952	39	19909	00100
63	19964	48	00000	00000
64	19976	34	19938	00701
65	19988	36	19938	00702

## POST PROCESSOR PHASE I

Statements 1 - 7 are changes to the main program. Statement 1 branches on indicator (equal) when the "FINI" card is read to initialize the starting sector address. Statement 2 branches to the read disk routine. Statement 3 branches to the write disk routine. Statement 4 branches to set up the calling of Phase 2 rather than Phase 3. Statement 5 branches to call the next program. Statement 6 changes the message from "Reload G.P. Output" to "Starting Pass Two." Statement 7 changes another message. The old message was, "Use Phase 2, Contouring", the new message is "Calling Ph 2, Contouring".

Statements 8 - 12 load the program on the disk.

Statements 13 - 21 load the changes into core.

Statements 22 - 31 are the read disk routines. Number 22 is a two position field to receive the transmission from the BTM entry. The group mark on the disk is the 81st character and this program only has 80 positions defined for the read in area so that the first character beyond the read in area must be saved. Statement 24 accomplishes this. Statement 25 then reads disk. Statement 26 turns off WLB/RBC console light, and statement 27 returns the digit moved by statement 24. Statement 29 returns to the main program.

Statements 34 - 44 are the write disk routine. Again in order to get all 80 positions on the disk the 81st character must be moved (#36) and a group mark placed in the 81st position (#37). Statement 38 writes on the disk, statement 39

turns off the WLR/RBC light and statement 40 replaces the digit moved.

Statement 41 adds one (1) to the sector address and statement 41 returns to the main program.

Statements 43 - 48 set up the program to call Phase 2 rather than Phase 3 if desired.

Statements 49 - 54 call the next program.

POST PROCESSOR PH1

READ-WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
1	11132	46	19784	01200
2	11556	17	19696	00000
3	19512	17	19806	00000
4	19596	17	19908	00000
5	19620	17	19952	00000
6	02033	STARTING PASS TWO (Alphamerically Coded)		
7	01985	CALLING PH 2 (Alphamerically Coded)		
CORE TO DISK				
8	01770	34	01808	00701
9	01782	16	00004	41000
10	01794	38	01808	00702
11	01806	48		
12	01808	10	58002	00000
		00		
LOADER				
13	00000	36	19694	00500
14	00012	36	19774	00500
15	00024	36	19854	00500

POST PROCESSOR PH1  
READ-WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
16	00036	36	19934	00500
17	00048	36	01770	00500
18	00060	36	00080	00500
19	00072	49	00080	
20	00080	36	00000	00500
21	00092	49	00000	
22	19694	00		
23	19696	34	19770	00701
24	19708	25	19703	01850
25	19720	36	19770	00700
26	19732	46	19744	03700
27	19744	25	01850	19703
28	19756	11	19775	00001
29	19768	42		
30	19770	10	20000	01017
31	19782	70		
32	19784	16	19775	02000
33	19796	49	11496	0
34	19804	00		



POST PROCESSOR PH1  
READ-WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
35	19806	34	19892	00701
36	19818	25	19837	10086
37	19830	15	10086	0000 <del>7</del>
38	19842	38	19892	00700
39	19854	46	19866	03700
40	19866	25	10086	19837
41	19878	11	19897	00001
42	19890	42		
43	19892	10	00000	01100
44	19904	06		
45	19906	00		
46	19908	16	19939	06000
47	19920	39	01985	00100
48	19932	42		
49	19934	10	62002	00000
50	19946	00	0000	
51	19952	39	01931	00100
52	19964	48	00000	00000
53	19976	34	19934	00701
54	19988	36	19934	00702

## POST PROCESSOR PHASE II

Statements 1 - 16 are the read disk routine. The main program branches to the read statement (#4), then checks indicators (statements 5, 6, and 7). If the disk address check or read check indicator is on a branch to a seek (#3) is made. Following a correct transfer from disk to core the input data is transferred to the odd input address. Statements 8 - 12 are needed for this. Upon completion of transfer of the 80th character, statements 13 and 14 initialize the transmit digit instruction (#8). Statement 15 adds one (1) to the sector address and #16 returns to the main program.

Statements 17 - 29 are the write disk routine. Since there are only 80 positions defined as an output area, the 81st position must be saved (#21) in order to set a group mark (#22) to terminate the read instruction of the next program. Following this is the write disk instruction (#23), indicator checking instructions (statements 24, 25, and 26) and a branch to a seek (#19), if necessary. After the seek the digit is transmitted to the 81st position (#20) before returning to statement 21. Upon completion of the transfer from core to disk, the 81st digit is replaced (#27), one (1) is added to the sector address (#28) and statement 29 returns to the main program.

Because Phase 2 is not always used and only two areas are defined on the disk for input-output, it is necessary to move the data output by Phase 2 so that input for programs to come will be properly oriented. Statements 34 - 60 do this.

Statements 30 - 33 move the output exchange statements to the high end of core, out of the way of incoming data. On completion of the exchange, statements 61 - 70 set up the call of the next program.

Statements 71 - 73 are changes to the main program. #71 branches to the read routine. #72 branches to the write routine and 73 branches to end of job routine.

Statements 74 - 79 load the program on the disk.

Statements 80 - 94 load the changes into core.

POST PROCESSOR PH2

READ-WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
1	10720	10̄	00000	01150
2	10732	00		
3	10734	34	10720	00701
4	10746	36	10720	00700
5	10758	46	10734	03600
6	10770	46	10782	03700
7	10782	46	10734	00600
8	10794	25	02365̄	15000̄
9	10806	11	10800	00001̄
10	10818	11	10805	00001̄
11	10830	14	10805	15080̄
12	10842	47	10794	01200
13	10854	16	10800	02365̄
14	10866	16	10805	15000̄
15	10878	11	10725	00001̄
16	10890	49	01844	0
17	10898	10̄	20000	01024
18	10910	70		
19	10912	34	10898	00701

POST PROCESSOR PH2

READ WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
20	10924	15	02550	00000
21	10936	25	10935	02550
22	10948	15	02550	0000 <del>7</del>
23	10960	38	10898	00700
24	10972	46	10912	03600
25	10984	46	10996	03700
26	10996	46	10912	00700
27	11008	25	02550	10935
28	11020	11	10903	00001
29	11032	49	03264	0
30	11040	31	19618	11084
31	11052	31	19938	11404
32	11064	39	03475	00100
33	11076	49	19626	0
34	11084 (19618)	00	0	
35	11087 (19621)	00	000	
36	11092 (19626)	26	19625	10903
37	11104 (19638)	12	19625	02100
38	11116 (19650)	11	19620	00001
39	11128 (19662)	14	19625	00000

POST PROCESSOR PH2  
READ WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
40	11140 (19674)	47	19734	01300
41	11152 (19686)	12	19625	00100
42	11164 (19698)	11	19620	00001
43	11176 (19710)	14	19625	00000
44	11188 (19722)	46	19686	01100
45	11200 (19734)	34	19888	00701
46	11212 (19746)	36	19888	00702
47	11224 (19758)	46	19734	00600
48	11236 (19770)	34	19874	00701
49	11248 (19782)	38	19874	00702
50	11260 (19794)	46	19770	00700
51	11272 (19806)	11	19879	00100
52	11284 (19818)	11	19893	00100
53	11296 (19830)	12	19620	00001
54	11308 (19842)	14	19620	00000
55	11320 (19854)	47	19734	01200
56	11332 (19866)	49	19952	0
57	11340 (19874)	10	00001	00005
58	11352 (19886)	00		
59	11354 (19888)	10	20001	00005
60	11366 (19900)	00		

POST PROCESSOR PH2  
 READ WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
61	11368 (19902)	00	59456	24563
62	11380 (19914)	00	62634	15963
63	11392 (19926)	23	00574	8730 <del>7</del>
64	11404 (19938)	10	62002	00000
65	11416 (19950)	00		
66	11418 (19952)	39	19901	00100
67	11430 (19964)	48	00000	00000
68	11442 (19976)	34	19938	00701
69	11454 (19988)	36	19938	00702
70	11466	<del>7</del>		

CHANGE TO MAIN PROGRAM

71	01832	49	10758	00000
72	03252	49	10936	00400
73	03372	49	11040	00000

CORE TO DISK

74	15000	34	15038	00701
75	15012	16	00004	41000
76	15024	38	15038	00702
77	15036	48		

POST PROCESSOR PH2  
READ WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CODE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
78	15038	10	60001	20000
79	15050	00		
80	00000	36	00080	00500
81	00012	36	00160	00500
82	00024	36	15000	00500
83	00036	36	10720	00500
84	00048	36	10800	00500
85	00060	36	10880	00500
86	00072	36	10960	00500
87	00084	36	11040	00500
88	00096	36	11120	00500
89	00108	36	11200	00500
90	00120	36	11280	00500
91	00132	36	11360	00500
92	00144	36	11440	00500
93	00156	36	00000	00500
94	00168	49	00000	



### POST PROCESSOR PHASE III

Statements 1 - 18 are the read routine. The main program branches to the read instruction (#5) and transmits the starting core location into the disk word (#3 and 4). Statements 6, 7, and 8 check indicators and branch to "seek" (#1) if necessary. Statements 9 - 14 transfer the input data to the odd input address. Statements 15 and 16 initialize statement 9. Statement 18 returns to the main program.

Statements 19 - 32 are the output routine. The main program branches to #23, checks indicator (program switch 2) and if it is on punches a card (#31) then returns to the main program (#32). If program switch 2 is off the program will write the output on the disk and return to the main program (#30).

Statements 33 - 35 are "fill in zeros".

Statements 38 - 40 are the end of job message.

Statements 41 - 45 type end of job and call the next program.

Statements 46 - 51 are changes to the main program. Statement 46 "branches and transmits" to the read routine and 47, 48, 49, and 51 branch and transmit to the write routine. Statement 50 branches to the end of job routine.

Statements 52 - 57 transfer a core image to the disk.

Statements 58 - 67 load the changes into core.

POST PROCESSOR PH3  
READ WRITE DISK

<u>SATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CORE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
1	19524	34	19544	00701
2	19536	49	19558	0
3	19544	10̄	00000	01000
4	19556	00		
5	19558	36	19544	00700
6	19570	46	19582	03700
7	19582	46	19524	03600
8	19594	46	19524	00600
9	19606	25	03653	00000
10	19618	14	19612	03732
11	19630	46	19674	01200
12	19642	11	19612	00001
13	19654	11	19617	00001
14	19666	49	19606	0
15	19674	16	19612	03653
16	19686	16	19617	00000
17	19698	11	19549	00001
18	19710	42		
19	19712	34	19732	00701
20	19724	49	19746	0

POST PROCESSOR PH3

READ WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CORE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
21	19732	10	20000	02039
22	19744	78		
23	19746	46	19832	00200
24	19758	15	04139	00007
25	19770	38	19732	00700
26	19782	46	19794	03700
27	19794	46	19712	03600
28	19806	46	19712	00700
29	19818	11	19737	00002
30	19830	42		
31	19832	39	03979	00400
32	19844	42	00000	00000
33	19856	00	00000	00000
34	19868	00	00000	00000
35	19880	00	00000	000
36	19890	10	64001	70000
37	19902	00		
38	19904	59	45624	56300
39	19916	62	63415	96323
40	19928	00	43565	56507
41	19940	39	03085	00100
42	19952	39	19901	00100

POST PROCESSOR PH3

READ WRITE DISK

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CORE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
43	19964	48	00000	00000
44	19976	34	19890	00701
45	19988	36	19890	00702
CHANGE TO MAIN PROGRAM				
46	01940	17	19558	00000
47	02536	17	19746	03978
48	11754	17	19746	03978
49	12546	17	19746	03978
50	12570	49	19940	00000
51	13894	17	19746	03978
CORE TO DISK				
52	03654	34	03692	00701
53	03666	16	00004	41000
54	03678	38	03692	00702
55	03690	48		
56	03692	10	62002	00000
57	03704	00		

POST PROCESSOR PH3

LOADER

<u>STATEMENT NUMBER</u>	<u>CORE LOCATION</u>	<u>OP CORE</u>	<u>P ADDRESS</u>	<u>Q ADDRESS</u>
58	00000	36	00080	00500
59	00012	36	03654	00500
60	00024	36	19524	00500
61	00036	36	19604	00500
62	00048	36	19684	00500
63	00060	36	19764	00500
64	00072	36	19844	00500
65	00084	36	19924	00500
66	00096	36	00000	00500
67	00108	49	00000	

IV

Appendix 3

Sample Problem

PG	LN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
	001			AUTOSPOT SEPTEMBER 30, 1962			
	002			AUTOSPOT PHASE2			
	003			PHASE2 SCANS OUTPUT OF PHASE1 RECORD BY RECORD			
	004			SW2 ON CARD INPUT/OUTPUT			
	005			SW4 OFF PRINT OUTPUT OF PHASE2			
	006			PUSH START TO CONTINUE ON ERROR STOPS			
	007			ERROR1 UNIDENTIFIED REC TYPE			
	008			ERROR2 TOOL NUMBER NOT IN TOOL TABLE			
	009			ERROR3 PATTERN TABLE EXCEEDED			
	009			ERROR4 PATTERN LENGTH EXCEEDED			
	009			ERROR5 PATTERN GENERATION EXCEEDED			
	009			ERROR6 PATTERN OPERATION ERROR			
1	126		DORG	01700	01700		
1	130	NEXT	RC2	CARDS	01700 46	01732	00200
1	131		RMP1	INPUT-47,,,TAPE INPUT	01712 36	12025	00300
1	132		RCOMP		01724 49	01840	00000
1	133		DORG	#-4	01731		
1	134	CARDS	RNF	CARDX,PFLAG-1,,CARD INPUT	01732 44	01780	13261
1	138		CF	PFLAG-1	01744 33	13261	00000
1	140	RCARD	RNCD	CARD1-79,,,READ FIRST CARD	01756 36	13056	00500
1	142		TFM	INLOC,CARD1-79	01768 16	13220	J3056
1	144	CARDX	RNF	RCARD,INLOC,11,NEXT CARD	01780 44	01756	1322-
1	146		TR	INPUT-47,INLOC,11	01792 31	12025	1322-
1	148		AM	INLOC,1,10,INDEX INPUT POINTER	01804 11	13220	000-1
1	150		RNR	#-12,INLOC,11	01816 45	01804	1322-
1	152		AM	INLOC,1,10	01828 11	13220	000-1
1	161	RCOMP	TFH	COMPT,COMPL62	01840 16	12752	J2754
1	162	SERCHI	CM	COMPT,COMPL620	01852 14	12752	J2772
1	183		BH	TRY26	01864 46	01988	01100
1	185		C	INPUT-46,COMPL,11	01876 24	12026	1279K
1	190		BE	CFIND	01888 46	01920	01200
1	200		AM	COMPL,2,10	01900 11	12752	000-2
2	010		B	SERCHI	01912 49	01852	00000
2	020		DORG	#-4	01919		
2	021	CFIND	TR	OUTPUT-47,INPUT-47	01920 31	12916	12025
2	022		RNF	#G36,CFIND	01932 44	01968	01920
2	023		ST	OUTPUT-46	01944 32	12917	00000
2	024		CF	CFIND	01956 33	01920	00000
2	025		BFM	OPUT,,7	01968 17	11484	-0000
2	026		B	NEXT	01980 49	01700	00000
2	028		DORG	#-4	01987		
2	051	TRY26	CM	INPUT-46,26,10,TEST FOR GO TO	01988 14	12026	000K6
2	052		DE	GOTO	02000 46	02036	01200
2	053		CM	INPUT-46,39,10,TEST FOR STOP	02012 14	12026	000L9
2	054		BJE	CONT	02024 47	02232	01200
2	055	GOTO	RNF	MAJCK,PFLAG11,,TEST FOR PATTERN	02036 44	02104	13273
2	056		RD	CFIND,INPUT-45	02048 43	01920	12027

49 19742

PG	LN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
2	057		BFM	SP,INPUT-47,7,STD IN PAT TABLE	02060 17	11288	J2025
2	058		CM	RTYPE,1,10	02072 14	12127	000-1
2	058		BE	NEXT,,,READ NEXT RECORD	02084 46	01700	01200
2	058		B	CFIND	02096 49	01920	00000
2	058		DORG	#-4	02103		
2	058	MAJCK	CM	INPUT-44,,10	02104 14	12028	000-0
2	058		RNE	CFIND	02116 47	01920	01200
2	058		CM	INPUT-42,11,10,TEST FOR Z COORD	02128 14	12030	000J1
2	058		BE	CFIND	02140 46	01920	01200
2	058		FA	INPUT-32,1XTY-11	02152 16	00469	-2187
2	058				02164 16	00445	J2040
2	058				02176 49	00422	J2114
2	058				02188 16	00469	-2223
2	058		FA	INPUT-22,1XTY-1	02200 16	00445	J2050
2	058				02212 49	00422	J2124
2	058				02224 49	01920	00000
2	058		B	CFIND	02231		
2	058	CONT	DORG	#-4	02231		
2	058		CM	INPUT-48,10,10,TEST FOR NEW MTRND SECT	02232 14	12026	000J0
2	069		BE	PATCK	02244 46	02752	01200
2	080		BH	CODE	02256 48	02868	01100
2	030		RNF	NEXT,PFLAG	02258 44	01700	13262
2	030		CF	PFLAG	02280 33	13262	00000
2	030		RNF	#G36,INPUT-46	02292 44	02328	12026
2	080		CF	INPUT-46	02304 33	12026	00000
2	080		SF	CFIND	02316 32	01920	00000
2	080		TF	RTYPE,INPUT-46,,SAVE REC TYPE	02328 26	12127	12026
2	081		RNF	CKODD,PFLAG11	02340 44	02436	13273
2	082		CT	PFLAG11	02352 33	13273	00000
2	083		RD	PLONG,PFLAG11,,TEST PAT LENGTH	02364 43	11172	13263
2	084		SF	PFLAG12	02376 32	13264	00000
2	085		TF	SCNCT,PFLAG14,6,RESET CNTR	02388 26	1325N	13266
2	085		TFM	PFLAG14,,0	02400 16	13266	0-000
2	088		TF	TYLOC,PLEV,6,SET PAT GENERATION INDICATOR	02412 26	1279-	13277
2	089		TFM	PLEV,,10	02424 16	13277	000-0
2	120	CKODD	TFM	INPUT-46,05,10,TEST FOR NEW PAT	02436 13	12026	000-5
2	130		B3	PATN,99	02448 43	02600	00099
2	140	EVEN	CM	INPUT-46,,10	02460 14	12026	000-0
2	144		BE	NEXT	02472 46	01700	01200
2	144		TF	TYPE,INPUT-43,,SAVE OPERATION CODE	02484 26	12130	12029
2	147		CM	INPUT-48,02,10	02496 14	12026	000-2
2	148		BE	CFIND	02508 46	01920	01200
2	149		CM	INPUT-46,08,10	02520 14	12026	000-8
2	150		RF	SIX	02532 46	02580	01200
2	153		TF	TOLNO,INPUT-39,,SAVE TOOL NUMBER	02544 26	12134	12033
2	154		CM	INPUT-46,04,10	02556 14	12026	000-4
2	155		RF	CFIND	02568 46	01920	01200
2	158	SIX	TF	TOLMOD,INPUT-36,,SAVE TOOL MOD	02580 26	12177	12036
2	159		B	CFIND	02592 49	01920	00000

PG	LIN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
2	160		DORG	*-4			02599
2	171	PAIN	SF	PFLAG11,,,PATTERN OPERATION	02600	32	13273 00000
2	171		SF	PFLAG10,,,FOR SAVE	02612	32	13272 00000
2	172		SM	INPUT-46,01,10,SET RECORD TYPE EVEN	02624	12	12026 000-1
2	173		AM	PIAH,01,10	02636	11	13260 000-1
2	177		AM	PCNT,01,10	02648	11	13275 000-1
2	178		TF	PIAR,PCNT,6,,TRANSFER PATTERN NUMBER	02660	26	1326- 13275
2	178		AM	PIAR,02,10	02672	11	13260 000-2
2	178		TF	TYLUC,PIAR	02684	26	12790 13260
2	179		AM	PIAR,03,10	02696	11	13260 000-3
2	180		TF	SCNCT,PIAR,,,ADDRESS OF SCAN COUNT FOR THIS PATTERN	02708	26	13295 13260
2	181		AM	PIAR,01,10	02720	11	13260 000-1
2	182		AM	PFLAG4,7,10	02732	11	13266 000-7
2	184		B	EVEN	02744	49	02460 00000
2	186		DORG	*-4			02751
2	186	PATCK	SF	PFLAG	02752	32	13262 00000
2	187		DNF	MINOR,PFLAG11	02764	44	02812 13273
2	188		BTH	SP,INPUT-45,7,STO IN PAT TABLE	02776	17	11288 J2027
2	193		CM	RTYPE,01,10	02788	14	12127 000-1
2	194		BF	NEXT,,,01 REC TYPE-NO INPUT	02800	46	01700 01200
2	198	MINOR	TR	TXTY-20,INPUT-43,,SAVE TX AND TY	02812	31	12105 12029
3	010		TR	OUTPUT-47,INPUT-47	02824	31	12916 12025
3	020		TR	OUTPUT-43,AUXIN,,SET CONTROL REC MARK	02836	31	12920 12104
3	022		BTH	DPUT,,,OUTPUT THIS RECORD	02848	17	11484 -0000
3	040		B	NEXT	02860	49	01700 00000
3	050		DORG	*-4			02867
10	010	CODE	TFM	COMP2,COMP262	02868	16	12777 J2779
10	012	FIND2	CM	COMP2,COMP268	02880	14	12777 J2785
10	014		BH	TRY13	02892	46	02948 01100
10	016		C	INPUT-46,COMP2,11,TEST FOR PATTERN OPERATION	02904	24	12026 1277P
10	018		BC	PATOP	02916	46	03748 01200
10	020		AM	COMP2,02,10	02928	11	12777 000-2
10	022		B	FIND2	02940	49	02880 00000
10	014		DORG	*-4			02947
3	060	TRY13	CM	INPUT-46,16,10	02948	14	12026 00016
3	070		BH	AUX,,,TEST FOR AUXILIARY SECTION	02960	46	03124 01100
3	072		RNF	TRANS,PFLAG11,,TEST FOR PATTERN	02972	44	03020 13273
3	074		BTH	SP,INPUT-47,7,STO IN PAT TABLE	02984	17	11288 J2025
3	075		CM	RTYPE,1,10,TEST FOR DEFINE STATEMENT	02996	14	12127 000-1
3	088		BE	NEXT,,,01 REC TYPE	03008	46	01700 01200
3	089	TRANS	CM	INPUT-46,11,10	03020	14	12026 000J1
3	090		BE	CFIND	03032	46	01920 01200
3	100		FA	INPUT-36,TXTY-11	03044	16	00469 -3079
3	100				03056	16	00445 J2036
3	100				03068	49	00422 J2114
3	110		FA	INPUT-26,TXTY-1	03080	16	00469 -3115
3	110				03092	16	00445 J2046
3	110				03104	49	00422 J2124

PG	LIN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
3	120		B	CFIND	03116	49	01920 00000
3	130		DORG	*-4			03123
3	140	AUX	CM	INPUT-46,49,10	03124	14	12026 000M9
3	142		RL	TRY30	03136	47	03232 01300
3	144		CM	INPUT-46,51,10,TEST FOR DAIMETER	03148	14	12026 000N1
3	146		BE	EQU91	03160	46	03264 01200
3	150		CM	INPUT-46,65,10	03172	14	12026 000O5
3	152		BL	CFIND	03184	47	01920 01300
3	154		CM	INPUT-46,91,10,TEST FOR TOOL CARD	03196	14	12026 000R1
3	156		BF	FJ091	<del>03208</del>	<del>46</del>	<del>19942</del> <del>01200</del>
3	158	ERR1	BTH	ERSUB,7100,8,WRITE ERROR MESSG	03220	17	11190 0P100
3	167		DORG	19942			19942
3	168	EQU91	TF	IAPT,INPUT-39,6,SAVE TOOL NO	19942	26	1322N 12033
3	174		AM	IAPT,10,10	19954	11	13225 000J0
3	176		TF	IAPT,INPUT-17,6,SAVE TIP ANGLE	19966	26	1322N 12055
3	178		AM	IAPT,04,10	19978	11	13225 000-4
3	180		B	ERSUB,36,,,OUTPUT TOOL CARD	19990	49	11226 00000
3	182		DORG	03232			03232
3	184	TRY30	CM	INPUT-46,30,10,TEST FOR REMARK STAT	03232	14	12026 000L0
3	186		BE	ERSUB,36	03244	46	11226 01200
3	188		B	ERR1,,,ILLEGAL CODE	03256	49	03220 00000
3	190		DORG	*-4			03263
4	040	EQU91	CM	TYPE,7,9,TEST FOR CSINK OR SPRILL	03264	14	12130 00-07
4	050		BE	ANGLE	03276	46	03312 01200
4	060		CM	TYPE,009,9	03288	14	12130 00-09
4	070		BNE	CFIND	03300	47	01920 01200
4	072	ANGLE	TFM	ADRES,FBLE-556	03312	16	13230 J2191
4	078	SEARCH2	CM	ADRES,FBLE-19,,SCAN FOR TOOL NO	03324	14	13230 J2728
4	080		RL	*G,4	03336	47	03360 01300
4	081		BTH	ERSUB,7200,8,TOOL NO MISSING	03348	17	11190 0P200
4	086		C	TOLNO,ADRES,11	03360	24	12134 1323-
4	088		BF	ANGLUC	03372	46	03404 01200
4	090		AM	ADRES,14,10	03384	11	13230 000J4
4	092		B	SEARCH2	03396	49	03324 00000
4	094		DORG	*-4			03403
4	096	ANGLUC	AM	ADRES,10,10,COMPUTE DEPTH	03404	11	13230 000J0
4	120		FM	-ADRES,CONV	03416	16	00469 -3451
4	120				03428	26	01260 1323-
4	120				03440	49	01262 J3409
4	125		TF	THETA,99,,SAVE CONVERTED ANGLE	03452	26	12154 00099
4	130		FD	THETA,100	03464	16	00469 -3499
4	130				03476	26	01260 12154
4	130				03488	49	01422 J2144
4	134		TF	THETA,00099	03500	26	12154 00099
4	140		FSIN	BUFFN,THETA	03512	16	14270 -3547
4	140				03524	16	14258 J2164
4	140				03536	49	13548 J2154
4	150		FCOS	BUFFD,THETA	03548	16	14270 -3583

46 19466 01200



PG	LN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
4	150				03560	16	14258 J2174
4	150				03572	49	13515 J2154
4	160		FD	BUFFN, BUFLD	03584	16	00469 -3619
4	160				03596	26	01260 J2164
4	160				03608	49	01422 J2174
4	170		IF	INPTA, 99	03620	26	12154 00099
4	190		FD	INPUT-36, TWD	03632	16	00469 -3667
4	190				03644	26	01260 J2036
4	190				03656	49	01422 J2144
4	200		IF	INPUT-36, 00099	03668	26	12036 00099
5	010		FD	INPUT-36, THCTA	03680	16	00469 -3715
5	010				03692	26	01260 J2036
5	010				03704	49	01422 J2154
5	020		IF	INPUT-36, 99, COMPUTED DEPTH TO BUFFER	03716	26	12036 00099
5	040		DM	INPUT-46, SET DEPTH CODE	03728	19	12026 00000
5	040		B	CTIND, INPUT DEPTH	03740	49	01920 00000
5	060		DORG	*-4	03747		
10	019	PATOP	SF	PATOP, PATTERN OPERATION	03740	32	03748 00000
10	020		SF	PFLAG	03760	32	13262 00000
10	021		AM	RIYPE, 5, 10, TEST FOR HIGHER GENERATION PAT	03772	13	12127 000-3
10	040		BD	ODD40, 00099	03784	43	03804 00099
10	050		B	PDUMP, OUTPUT THIS PATTERN	03796	49	04000 00000
10	055		DORG	*-4	03803		
10	060	DDL40	BNF	STORF, PFLAG&10	03804	44	03864 13272
10	070		CF	PFLAG&10	03816	33	13272 00000
10	090		SN	PTAB, 04, 10, INDEX POINTER BACKWARDS	03828	12	13260 000-4
10	099		IF	FYLOC, PTAB, SAVE ADD OF GENERATION INDICATOR	03840	26	12790 13260
10	100		AM	PTAB, 04, 10	03852	11	13260 000-4
10	110	STORE	BIM	SP, INPUT-41, 7, STORE HIGHER GENERATION PATTERN	03864	17	11280 J2025
10	111		BD	LEVOK, PLEV-1	03876	43	03976 13276
10	113		TF	PATNO, INPUT-42, SAVE PATTERN NO	03888	26	12890 12030
10	114		BT	PATLOC, PATNO, LOCATE THIS PAT IN PAT TABLE	03900	27	11396 12890
10	115		AM	COMP966, 2, 10, INDEX FST LEV POINTER	03912	11	11414 000-2
10	116		RD	*620, COMP966, 11, TEST PAT LEVEL	03924	43	03944 1141M
10	117		B	TFM1	03936	49	03964 00000
10	118		DORG	*-4	03943		
10	118		IFM	PLEV, 11, 10, SET PAT LEV INDICATOR	03944	16	13277 000J1
10	119		H	LEVOK	03956	49	03976 00000
10	120		DORG	*-4	03963		
10	121	TFM1	TFM	PLEV, 1, 10	03964	16	13277 000-1
10	125	LEVOK	CM	RIYPE, 1, 10, TEST FOR OUTPUT	03976	14	12127 000-1
10	150		BE	NEXT	03988	46	01700 01200
15	030	PDUMP	TF	PATNO, INPUT-42	04000	26	12890 12030
15	040		BT	PATLOC, PATNO, LOCATE THIS PATTERN	04012	27	11396 12890
15	050		AM	COMP966, 02, 10, INDEX POINTER FORWARD	04024	11	11414 000-2
15	052		BNF	*836, PATOP, SWITCH FOR PATTERN GENERATION	04036	44	04072 03748
15	054		CF	PATOP	04048	33	03748 00000
15	056		TF	GNO, COMP966, 11, SAVE GENERATION LEV INDICATOR	04060	26	13515 1141M

PG	LN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
15	060		BNF	CKRM2, PFLAG&7, TEST 3RD LEV FLAG	04072	44	04268 13269
15	070		BD	*620, COMP966, 11	04084	43	04104 1141M
15	080		B	RING1	04096	49	04344 00000
15	090		DORG	*-4	04103		
15	100		BNF	LOAD2, PFLAG&8, IS A PAT DEFINED ON A PAT	04104	44	04128 13270
15	105		BTM	ERSUB, 7500, 6, TEST FOR PAT GEN ERROR	04116	17	11190 0P300
15	120	LOAD2	SF	PFLAG&8, SEC LEV SWITCH	04128	32	13270 00000
15	130		TF	PLOC2, COMP966, LOAD SEC GEN OPERATION BUFFERS	04140	26	12795 11414
15	140		AM	PLOC2, 03, 10	04152	11	12795 000-3
15	150		TF	PCNT2, PLOC2	04164	26	12800 1279
15	160		A	PCNT2, PLOC2, 11	04176	21	12800 1279M
15	170		SN	PCNT2, 08, 10, SET LENGTH OF THIS PATTERN	04188	12	12800 000-4
15	172		AM	PLOC2, 1, 10	04200	11	12795 000-1
15	173		TR	OPER2-1, PLOC2, 11, SAVE OPER, SURF, PATNO, AND CORD TRANS	04212	31	12801 1279M
15	174		TF	PATNO, SURF262	04224	26	12890 12806
15	175		AM	PLOC2, 1, 10	04236	11	12795 000-1
15	176		BNK	*-12, PLOC2, 11, INDEX SCAN POINTER	04248	45	04236 1279M
15	330		B	PDUMP12	04260	49	04012 00000
15	340		DORG	*-4	04267		
15	350	CKRM2	BNF	CKSTM, PFLAG&8, TEST FOR 2ND GEN PAT	04268	44	04332 13270
15	360		BD	*620, COMP966, 11	04280	43	04300 1141M
15	370		B	RING1	04292	49	04344 00000
15	380		DORG	*-4	04299		
15	390		TF	PLOC3-4, PLOC2-4, SET 3RD LEV OPER BUFFERS	04300	31	12840 12791
15	430		SF	PFLAG&7, SET 3RD GEN PAT SWITCH	04312	32	13269 00000
15	440		IF	LOAD2&12, RELOAD 2ND LEV BUFFS	04324	49	04140 00000
15	450		DORG	*-4	04331		
15	460	CKSTM	BD	LXAS, COMP966, 11, TEST FOR SIMPLE PAT	04332	43	04128 1141M
15	470	RING1	AM	COMP966, 03, 10	04344	11	11414 000-3
15	480		TF	PTCNT, COMP966	04356	26	12182 11414
15	490		A	PTCNT, COMP966, 11	04368	21	12182 1141M
15	500		SN	PTCNT, 08, 10, SET LENGTH OF PAT	04380	12	12182 000-8
15	505		TF	PTCNT, PTCNT, SET PAT LIMIT FOR REVERSE OPERATION	04392	26	12187 12182
15	510		AM	COMP966, 02, 10	04404	11	11414 000-2
15	520		TF	OUTPUT-44, COMP966, 11, SURF TO OUTPUT BUFF	04416	26	12919 1141M
15	530		BNF	FLAG7, PFLAG&8	04428	44	04472 13270
15	540		BD	*620, SURF2-1	04440	43	04460 12803
15	545		B	FLAG7	04452	49	04472 00000
15	550		DORG	*-4	04459		
15	560		TF	OUTPUT-44, SURF2	04460	26	12919 12804
15	600	FLAG7	BNF	PRIOR, PFLAG&7	04472	44	04516 13269
15	620		BT	*620, SURF3-1	04484	43	04504 12852
15	625		B	PRIOR	04496	49	04516 00000
15	630		DORG	*-4	04503		
15	640		TF	OUTPUT-44, SURF3	04504	26	12919 12853
15	650	PRIOR	BD	*620, INPUT-45, TEST FOR NEW SURF	04516	43	04536 12027
15	655		B	TRA10	04528	49	04548 00000
15	660		DORG	*-4	04535		

PG	LN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
15	670		TF	OUTPUT-44,INPUT-44	04536	26	12919 12028
15	680	TRALO	TFM	OUTPUT-46,10,10	04548	16	12917 00010
15	690		TR	OUTPUT-43,AUXIN	04560	31	12920 12104
15	691		BTM	OPUT,,,OUTPUT DASH SURF	04572	17	11484 -0000
17	030	SETUP	AM	COMP966,01,10	04584	11	11414 000-1
17	032		TR	ST-20,COMP966,11,SAVE COORD TRANSFORM	04596	31	12891 1141M
17	036		AM	COMP966,21,10	04608	11	11414 000K1
17	038	XHAIN	TR	OUTPUT-47,COMP966,11,EST MACHINE POINT TO OPUT BUFF	04620	31	12918 1141M
17	039		AM	COMP966,1,10	04632	11	11414 000-1
17	040		CM	OUTPUT-46,26,10,TEST FOR GOVD OR STOP	04644	14	12917 000K6
17	042		BL	*624	04656	47	04680 01300
17	048		TR	OUTPUT-51,OUTPUT-47,,SHIFT REC IN OPUT BUFF	04668	31	12912 12916
15	900		BD	ER30,CNO,,TEST FOR 1ST GEN PAT	04680	43	04744 13515
15	901		TFM	LASTIG94,*620,,SET RETURN ADDRESS	04692	16	06206 -4712
15	902		D	EXEC	04704	49	05896 00000
15	903		DORC	*-4	04711		
15	904		SF	TRAN	04712	32	08516 00000
15	904		SF	TRIGF	04724	32	13267 00000
15	905		B	ANYMO	04736	49	04894 00000
15	906		DORC	*-4	04743		
15	907	ERJ6	CM	INPUT-46,36,10	04744	14	12026 000L6
15	908		BNE	TRYC2	04756	47	04780 01200
15	909	ERR2	BTM	EKSUB,7600,,PAT OPER ERROR	04768	17	11190 0P600
15	914	TRYC2	BD	G3,CNO-1	04780	43	04856 13514
15	915		TFM	NXT11-2,*620	04792	16	06734 -4612
15	916		B	EXEC2,,SECOND GEN PAT	04804	49	06264 00000
15	917		DORC	*-4	04811		
15	918		TFM	OPER2,,10	04812	16	12802 000-0
15	919		CF	DIRAN	04824	33	09120 00000
15	920		SF	DI	04836	32	09348 00000
15	923		D	ER36-32	04848	49	04712 00000
15	924		DORC	*-4	04855		
15	925	G3	CM	OPCR3,36,0110,3RD GEN PAT	04856	JM	12851 000L6
15	925		BE	ERR2	04868	46	04760 01200
15	925		RNF	EXEC3,G3	04880	44	05028 04856
15	926		CF	G3	04892	33	04856 00000
15	927		CM	IRAN3-31,12,10,TEST TO SAVE 1ST PT OF PAT	04904	14	12857 000J2
15	928		BL	LOC1	04916	47	04960 01300
15	931		TF	FG3X,IRAN3-21	04928	26	13481 12867
15	932		TF	FG3Y,IRAN3-11	04940	26	13491 12877
15	933		B	EXEC3	04952	49	05028 00000
15	934		DORC	*-4	04959		
15	936	LOC1	CM	IRAN2-31,12,10	04960	14	12808 000J2
15	938		PL	LOC3	04972	47	05016 01300
15	940		TF	FG3X,IRAN2-21	04984	26	13481 12818
15	941		TF	FG3Y,IRAN2-11	04996	26	13491 12828
15	942		D	EXEC3	05008	49	05028 00000
15	943		DORC	*-4	05015		

PG	LN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
15	944	LOC3	TR	FG3X-9,ST-20,,SAVE 1ST PT OF PAT	05016	31	13472 12891
15	944	EXEC4	CM	TYPE-1,2,10,TEST FOR LINE MILL	05028	14	12129 000-2
15	945		BNE	EX2C24	05040	47	05164 01200
15	966		BD	EX2,INPUT-46,,MILL INSTR-CHECK FOR INVERT	05052	43	05140 12026
15	967		BD	EX2-20,OPER3,,TEST FOR 3RD GEN INVERT	05064	43	05120 12851
15	968		CM	OPER2,36,10,,TEST FOR 2ND GEN INVERT	05076	14	12802 000L6
15	969		BHL	LOC7	05088	46	05208 01300
15	970		TFM	EX2E23,OPER2	05100	18	05163 12802
15	971		B	*640	05112	49	05152 00000
15	972		DORC	*-4	05119		
15	973		TFM	EX2E23,OPER3	05120	16	05163 12851
15	974		B	*620	05132	49	05152 00000
15	975		DORC	*-4	05139		
15	975	EX2	TFM	*E23,INPUT-46,,MILL INVERSION	05140	16	05163 12026
15	975		BTM	MILLR	05152	17	10236 -0000
15	975		CM	OPER2,34,10,TEST FOR 2ND GEN INVERSION	05164	14	12802 000L4
15	975		BH	LOC7	05176	46	05208 01100
15	975		BT	INVRT,OPER2	05188	27	07544 12802
15	976		B	LOC6	05200	49	05252 00000
15	976		DORC	*-4	05207		
15	977	LOC7	CM	OPER2,36,10,TEST FOR 2ND GEN REVERSAL	05208	14	12802 000L6
15	978		BNE	LOC6	05220	47	05252 01200
15	979		BTM	REVRS,36,10	05232	17	08184 000L6
15	980		B	LOC8	05244	49	05376 00000
15	981		DORC	*-4	05251		
15	982	LOC6	CM	IRAN2-31,11,10,TEST FOR 2ND GEN ROTAT W/TRANS	05252	14	12808 000J3
15	983		BNE	LOC9	05264	47	05320 01200
15	984		RNF	LOC10,TRIGF,,SWITCH FOR ANGLE CONVERSION	05276	44	05300 13267
15	985		BT	SICO,IRAN2-1,,COMPUTE SINE AND COSINE OF ROTAT ANGLE	05288	27	10096 12838
15	990	LOC10	BT	ROTD,SICO-1,,ROTATE THIS POINT	05300	27	07682 10095
15	991		B	LOC8	05312	49	05376 00000
15	992		DORC	*-4	05319		
15	993	LOC9	CM	IRAN2-31,11,10,TEST FOR 2ND GEN ROT W/O TRANS	05320	14	12808 000J1
15	994		BNE	LOC8	05332	47	05376 01200
15	995		RNF	LOC10,TRIGF,,HAS ANGLE BEEN CONVERTED	05344	44	05300 13267
15	996		BT	SICO,IRAN2-21	05356	27	10096 12818
15	998		B	LOC10	05368	49	05300 00000
15	999		DORC	*-4	05375		
16	100	LOC8	CM	OPER3,40,10,TEST FOR 2ND GEN DIFFERENCE TRANS	05376	14	12851 000M0
16	101		BNE	*636	05388	47	05424 01200
16	102		CM	IRAN3-31,10,10	05400	14	12857 000J0
16	103		WC	LOC11	05412	46	05596 01200
16	104		BTM	DIRAN,,10,COMPUTE 2ND GEN DIFF TRANS	05424	17	09120 000-0
16	105		CM	OPER3,34,10,TEST FOR 3RD GEN INVERSION	05436	14	12851 000L4
16	107		BH	LOC13	05448	46	05472 01100
16	109		BT	INVRT,OPER3	05460	27	07544 12851
16	115	LOC13	CM	IRAN3-31,13,10,TEST FOR 3RD GEN ROTATION	05472	14	12857 000J3
16	116		BNE	LOC14	05484	47	05540 01200

PG	LIN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
16	117		BNE	LOC19,TRIGF,,ANGLE TEST	05496	44	05520 13261
16	119		BT	SIC0,TRAN3-I	05508	27	10096 12687
16	123	LOC15	BT	ROT0,SIC0-I	05520	27	07682 10095
16	124		B	LOC11	05532	49	05596 00000
16	125		DORG	*-4	05539		
16	126	LOC14	CM	TRAN3-11,11,10,TEST FOR 3RD GEN ROT W/O TRANS	05540	14	12887 000J1
16	127		BNE	LOC11	05552	47	05596 01200
16	128		BNE	LOC15,TRIGF	05564	44	05520 13267
16	128		BT	SIC0,TRAN3-21	05576	27	10096 12667
16	130		B	LOC15	05588	49	05520 00000
16	131		DORG	*-4	05999		
16	132	LOC11	CM	INPUT-46,40,10,TEST FOR 3RD GEN DIFF TRANS	05596	14	12026 000M0
16	133		BNE	*E36	05608	47	05644 01200
16	134		CM	INPUT-40,10,10	05620	14	12032 000J0
16	135		BE	LOC18	05632	46	05816 01200
16	136		BTM	TRAN2,1,10,COMPUTE 3RD GEN DIFF TRANS	05644	17	08814 000-0
16	137		CM	INPUT-46,34,10,TEST FOR INVERSION	05656	14	12026 000L4
16	139		BNE	LOC16	05668	46	05892 01100
16	141		BT	INVRT,INPUT-46	05680	27	07544 12026
16	145	LOC16	CM	INPUT-40,13,10,TEST FOR ROT W/TRANS	05692	14	12032 000J3
16	146		BNE	LOC20	05704	47	05760 01200
16	147		BNE	LOC19,TRIGF	05716	46	05740 13267
16	148		BT	SIC0,INPUT-10	05728	27	10096 12062
16	153	LOC19	BT	ROT0,SIC0-I	05740	27	07682 10095
16	154		B	LOC18	05752	49	05816 00000
16	155		DORG	*-4	05759		
16	156	LOC20	CM	INPUT-40,11,10,TEST FOR ROT W/O TRANS	05760	14	12032 000J1
16	157		BNE	LOC18	05772	47	05816 01200
16	158		BNE	LOC19,TRIGF	05784	44	05740 13267
16	159		BT	SIC0,INPUT-30	05796	27	10096 12042
16	161		B	LOC19	05808	49	05740 00000
16	162		DORG	*-4	05815		
16	163	LOC18	BTM	TRAN,,10,FINAL TRANSFORMATION	05816	17	08816 000-0
16	165		BTM	OPUT,,,OUTPUT THIS POINT	05828	17	11484 -0000
16	164		C	COMP966,RTCNT	05840	24	11414 12187
16	165		BL	XMATN	05852	47	04620 01300
16	166		SF	D1,,,RESET SWITCHES ON LAST PT OF PAT	05864	32	09348 00000
16	167		CF	DTRAN	05876	33	09120 00000
16	168		B	ER36-32	05888	49	04712 00000
16	171		DORG	*-4	05895		
17	201	EXEC	CM	TYPE-1,2,10,TEST FOR LINE MILL	05896	14	12129 000-2
17	202		BNE	*E48	05908	47	05956 01200
17	203		CM	INPUT-46,34,10,TEST FOR MILL INVERSION	05920	14	12026 000L4
17	204		BNE	TRYRV	05932	46	06000 01100
17	204		BTM	MILLR,INPUT-46	05944	17	10236 J2026
17	204		CM	INPUT-46,34,10,INVERSION TEST	05956	14	12026 000L4
17	204		BH	*E32	05968	46	06000 01100
17	204		BT	INVRT,INPUT-46	05980	27	07544 12026

PG	LIN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
17	205		B	TRYRO	05992	49	06092 00000
17	206		DORG	*-4	05999		
17	208	TRYRV	CM	INPUT-46,36,10,TEST FOR REVERSAL	06000	14	12026 000L6
17	209		BNE	TRYRO	06012	47	06092 01200
17	210		BT	REVR5,INPUT-46,,36	06024	27	08184 12026
17	212		BTM	OPUT,,,OUTPUT THIS POINT	06036	17	11484 -0000
17	212		TFM	OUTPUT-90,,10,RESET GOTO OR STOP SWITCH	06048	16	12913 000-0
17	213		C	COMP966,RTCNT,,TEST FOR LAST POINT	06060	24	11414 12182
17	214		BH	ER36-32	06072	46	04712 01100
17	215		B	TRYRV624,,,GTY NEXT POINT	06084	49	06024 00000
10	216		DORG	*-4	06091		
17	217	TRYRO	CM	INPUT-40,13,10,TEST FOR ROT W/TRANS	06092	14	12032 000J3
17	218		BNE	TRYR1	06104	47	06208 01200
17	219		BNE	LAST1-12,TRIGF,,ANGLE CHECK	06116	44	06140 13267
17	220		BT	SIC0,INPUT-10	06128	27	10096 12062
17	225		BT	ROT0,SIC0-I	06140	27	07682 10095
17	226	LAST1	BT	TRAN,INPUT-40,,FINAL TRANS	06152	27	08816 12032
17	227		BTM	OPUT,,,OUTPUT THIS POINT	06164	17	11484 -0000
17	229		C	COMP966,RTCNT,,TEST FOR LAST PT THIS PAT	06176	24	11414 12187
17	230		BL	XMATN,,,NEXT POINT	06188	47	04620 01300
17	231		B	,,,RETURN TO MAIN ROUTINE	06200	49	00000 00000
17	232		DORG	*-4	06207		
17	233	TRYR1	CM	INPUT-40,11,10,TEST ROT W/O TRANS	06208	14	12032 000J1
17	234		BNE	LAST1	06220	47	06152 01200
17	235		BNE	LAST1-12,TRIGF	06232	44	06140 13267
17	236		BT	SIC0,INPUT-30	06244	27	10096 12042
17	238		B	LA,T1-12	06256	49	06140 00000
17	239		DORG	*-4	06263		
17	240	EXEC2	CM	TYPE-1,2,10,TEST FOR LINE MILL	06264	14	12129 000-2
17	241		BNE	E2	06276	47	06376 01200
17	242		BD	E3,INPUT-46,,TEST FOR MILL INVERSION	06288	43	06356 12026
17	243		CM	OPER2,34,10,TEST FOR 2ND GEN MILL INVERT	06300	14	12802 000L4
17	244		BH	NOW36	06312	46	06420 01100
17	245		TFM	*E23,OPER2	06324	16	06347 J2802
17	246	E4	BTM	MILLR,OPER2	06336	17	10236 J2802
17	247		B	*E28	06348	49	06376 00000
17	248		DORG	*-4	06355		
17	249	E3	TFM	E4611,INPUT-46	06356	16	06347 J2026
17	250		B	E4	06368	49	06336 00000
17	251		DORG	*-4	06375		
17	251	E2	CM	OPER2,34,10,TEST FOR 2ND GEN INVERT	06376	14	12802 000L4
17	251		BH	*E32	06388	46	06420 01100
17	251		BT	INVRT,OPER2	06400	27	07544 12802
17	251		B	NOW13	06412	49	06464 00000
17	251		DORG	*-4	06419		
17	251	NOW36	CM	OPER2,36,10,TEST INVERSION	06420	14	12802 000L6
17	251		BH	NOW13	06432	47	06464 01200
17	251		BTM	REVR5,,10	06444	17	08184 000-0

PG	LIN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
17	252		B	NOT1,12	06456	49	06524 00000
17	254		DORG	*-4	06463		
17	253	NOWL3	CM	TRAN2-31,13,10,TEST FOR 2ND GEN ROTAT W/TRANS	06464	14	12808 000J3
17	254		HNE	NOI11	06476	47	06792 01200
17	255		HNF	NOI1,TRIGF	06488	44	06512 13267
17	256		BT	SIC0,TRAN2-1	06500	27	10096 12838
17	261	NOI1	BT	ROTO,SIC0-1	06512	27	07682 10095
17	261		CM	INPUT-46,40,10,TEST FOR 2ND LEV DIFF TRAN	06524	14	12076 000M0
17	261		BNE	*E36	06536	47	06572 01200
17	261		CM	INPUT-40,10,10	06548	14	12032 000J0
17	261		BE	NOT2E12	06560	46	06680 01200
17	262		BTM	TRAN,,10	06572	17	09120 000-0
17	263		CM	INPUT-46,34,10,TEST FOR INVERT	06584	14	12026 000L4
17	264		DH	NXTI3	06596	46	06620 01100
17	265		BT	INVRT,INPUT-46	06608	27	07544 12026
17	266	NXTI3	CM	INPUT-40,13,10,TEST FOR ROT W/TRANS	06620	14	12032 000J3
17	267		HNE	NXTI1	06632	47	06736 01200
17	268		HNF	NOI2,TRIGF	06644	44	06888 13267
17	269		BT	SIC0,INPUT-10	06656	27	10096 12062
17	274	NOI2	BT	ROTO,SIC0-1	06668	27	07682 10095
17	275		BTM	TRAN,,10,FINAL TRANS	06680	17	08516 000-0
17	276		BTM	DPUT,,,OUTPUT THIS POINT	06692	17	11484 -0000
17	276		C	COMP96,RYCNT,,TEST FOR LAST PT OF PAT	06704	24	11414 12187
17	277		HL	XHAIN	06716	47	34620 01300
17	278				06728	49	06728 00000
17	279		DORG	*-4	06735		
17	280	NXTI1	CM	INPUT-40,11,10,TEST FOR ROT W/O TRANS	06736	14	12032 000J1
17	281		HNE	NOT2E12	06748	47	06680 01200
17	282		HNF	NOI2,TRIGF	06760	44	06688 13267
17	283		BT	SIC0,INPUT-30	06772	27	10096 12642
17	285		B	NOI2	06784	49	06668 00000
17	286		DORG	*-4	06791		
17	291	NOWI1	CM	TRAN2-31,11,10,TEST FOR 2ND GEN ROT W/O TRANS	06792	14	12808 000J1
17	292		HNE	NOT1E12	06804	47	06524 01200
17	293		HNF	NOI1,TRIGF	06816	44	06512 13267
17	293		BT	SIC0,TRAN2-21	06828	27	10096 12818
17	295		B	NOI1	06840	49	06512 00000
17	290		DORG	*-2	06849		
17	217	INIT	RCTY		06850	34	00000 00102
17	217		TFM	OUTPUT-50,00,10	06862	16	12913 000-0
17	217		TR	TANPT-4,TDLE1-4,,REINITIALIZE PHASE2	06874	31	13221 13279
17	217		R	NEXT	06886	49	01700 00000
17	217		DORG	*-4	06893		
15	715	ANYM0	HNF	NORFL,PFLASER,,TEST FOR 2ND GEN PAT	06894	44	07522 13270
15	720		C	PLOC2,PCNT2,,TEST FOR LAST SUBPATTERN	06906	24	12799 12800
15	730		DH	OFF2	06918	46	07278 01100
15	731		AM	PLOC2,02,10,INDEX POINTER FOR 2ND GEN PAT	06930	11	12795 000-2
15	734	TRYGO	CM	PLOC2,26,610,TEST FOR GOT0	06942	14	12799 000K6

PG	LIN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
15	734		BE	*E36	06954	46	06990 01200
15	734		CM	PLOC2,39,10,TEST FOR STOP	06966	14	12795 000L9
15	734		HNE	NOGO	06978	47	07210 01200
15	734		SM	PLOC2,01,10	06990	12	12795 000-1
15	734		TR	OUTPUT-47,PLOC2,11	07002	31	12916 12799
15	734		AM	PLOC2,1,10	07014	11	12795 000-1
15	734		FS	OUTPUT-32,FG2X	07026	16	00469 -7061
15	734				07038	16	00445 J2931
15	734				07050	49	00402 J3439
15	734		FS	OUTPUT-22,FG2Y	07062	16	00469 -7097
15	734				07074	16	00445 J2941
15	734				07086	49	00402 J3449
15	734		TR	ST-20,FG2X-9	07098	31	12891 13430
15	734		TFM	PICNT,,, THIS PAT INCLUDES A GOT0 OR A STOP	07110	16	12182 -0000
15	734		TF	COMP96,PLOC2	07122	26	11414 12795
15	734		TFM	LASTI,54,*E20,,OPERATE ON A GOT0 OR STOP	07134	16	06206 -7154
15	724		B	EXEC	07146	49	05896 00000
15	724		DORG	*-4	07153		
15	734		SF	TRIGF	07154	37	13267 00000
15	734		SF	TRAN	07166	37	08516 00000
15	734		AM	PLOC2,1,10	07178	11	12795 000-1
15	734		BNR	*-12,PLOC2,11,INDEX 2ND GEN POINTER	07190	45	07178 12799
15	734		B	TRYGO-12	07202	49	06930 00000
15	734		DORG	*-4	07209		
15	736	NOGO	SM	PLOC2,1,10	07210	12	12795 000-1
15	738		TR	OPER2-1,PLOC2,11,RELOAD 2ND GEN BUFFERS	07222	31	12801 12799
15	740		TF	PATNO,SURF2&2	07234	26	12890 12806
15	742		AM	PLOC2,1,10	07246	11	12795 000-1
15	744		BNR	*-12,PLOC2,11	07258	45	07246 12799
15	850		B	PDUMP&12,,,OUTPUT THIS PAT	07270	49	04012 00000
15	860		DORG	*-4	07277		
15	900	OFF2	CF	PFLAG0,,END OF 2ND GEN PAT	07278	33	13270 00000
15	905		SF	DTRAN,,RESET 2ND GEN CONTROL SWITCHES	07290	32	09120 00000
15	907		SF	DIG1	07302	32	09349 00000
15	908		CF	D4	07314	33	09752 00000
15	909		CF	D&E1	07326	32	09753 00000
15	908		TFM	FG2X-7,500,9	07338	16	13432 00000
15	909		TFM	FG2Y-7,500,9	07350	16	13442 00000
15	915		HNF	LEAV,,PFLAG&7,,TEST FOR 3RD GEN PAT	07362	44	07514 13269
15	920		C	PLOC3,PCNT3,,TEST FOR LAST SUBPATTERN	07374	24	12844 12849
15	930		DH	OFF3	07386	46	07466 01100
15	932		AM	PLOC3,1,10,INDEX 3RD GEN POINTER	07398	11	12844 000-1
15	935		TR	OPER3-1,PLOC3,11,RELOAD 3RD GEN OPERATION BUFFERS	07410	31	12850 12844
15	937		TF	PATNO,SURF3&2	07422	26	12890 12855
15	939		AM	PLOC3,1,10	07434	11	12844 000-1
15	941		BNR	*-12,PLOC3,11,INDEX 3RD GEN POINTER	07446	45	07434 12844
16	060		B	PDUMP&12,,,OUTPUT NEXT PATTERN	07458	49	04012 00000
16	070		DORG	*-4	07465		

PG	LN	LABEL	OP	OPCRANDS	LOCN	OP	INSTRUCTIONS
16	080	DEFF	CF	PFLAG7,,RIGHT OR GEN SWITCHES	07466	33	13269 00000
16	081		SF	G3	07478	32	04856 00000
16	084		ST	DTRAN	07490	32	09614 00000
16	086		ST	D4	07502	32	09752 00000
16	090	LEAVE	P	DEXT,,LOAD NEXT INPUT RECORD	07514	49	01700 00000
16	100		DORG	*-4	07521		
16	110	NOBFL	RNF	LEAVE,PFLAG7,,TEST FOR END OF ENTIRE PAT	07522	44	07514 13269
16	020		R	DEFT2,,OUTPUT NEXT PAT	07534	49	07290 00000
16	130		DORG	*-2	07543		
17	300			INVERSION SUBROUTINE			
17	302	INVRT	CM	OUTPUT-46,11,10,TEST FOR Z COORD	07544	14	12917 000J1
17	303		RF	PTOMX612	07596	46	07630 01200
17	310		CM	INVRT-1,32,10	07568	14	07543 000L2
17	311		BNF	FLIPY	07580	47	07632 01200
17	312		BNF	PTOMX,OUTPUT-36,,TEST FOR INVERT RIGHT OR LEFT	07592	44	07618 12927
17	320		CF	OUTPUT-36,,MINUS TO PLUS	07604	33	12927 00000
17	322		RR		07616	42	00000 00000
17	324		DORG	*-10	07617		
17	326	PTOMX	SF	OUTPUT-36,,PLUS TO MINUS	07618	32	12927 00000
17	328		RR		07630	42	00000 00000
17	330		DORG	*-10	07631		
17	332	FLIPY	BNF	PTOMY,OUTPUT-26,,TEST FOR INVERT UP OR DOWN	07632	44	07658 12937
17	334		CF	OUTPUT-26	07644	33	12937 00000
17	336		RR		07656	42	00000 00000
17	338		DORG	*-10	07657		
17	340	PTOMY	SF	OUTPUT-26	07658	32	12937 00000
17	342		RR		07670	42	00000 00000
17	400			ROTATION SUBROUTINE			
17	402	ROTO	CM	OPER2,36,10,TEST FOR REVERSAL	07682	14	12802 000L6
17	402		RNE	R4	07694	47	07776 01200
17	402		AM	PTCNT,03,10	07706	11	12182 000-3
17	402		CM	PTCNT,11,610	07718	14	1218K 000J1
17	402		BF	*632	07730	46	07762 01200
17	402		SM	PTCNT,03,10	07742	12	12182 000-3
17	402		R	R4624	07754	49	07800 00000
17	402		DORG	*-4	07761		
17	402		SM	PTCNT,03,10	07762	12	12182 000-3
17	402		RR		07774	42	00000 00000
17	402		DORG	*-10	07775		
17	402	R4	CM	OUTPUT-46,11,10,TEST FOR Z COORD ONLY	07776	14	12917 000J1
17	403		RF	REVR5-4	07788	46	08180 01200
17	414		FM	OUTPUT-36,THCOS	07800	16	00469 -7835
17	414				07812	26	01260 12927
17	414				07824	49	01262 J3419
17	416		TF	BUFFN,99,, X.COS*THETA TO BUFFER	07834	26	12164 00099
17	418		FM	OUTPUT-26,THSIN	07848	16	00469 -7883
17	418				07860	26	01260 12937
17	418				07872	49	01262 J3429

PG	LN	LABEL	OP	OPCRANDS	LOCN	OP	INSTRUCTIONS
17	420		FS	BUFFN,99	07884	16	00469 -7919
17	420				07896	16	00445 J2164
17	420				07908	49	00402 -0099
17	422		FM	OUTPUT-26,THCOS	07920	16	00469 -7955
17	422				07932	26	01260 12937
17	422				07944	49	01262 J3419
17	424		TF	BUFFD,99,, Y.COS*THETA	07956	26	12174 00099
17	426		FM	OUTPUT-36,THSIN	07968	16	00469 -8003
17	426				07980	26	01260 12927
17	426				07992	49	01262 J3429
17	428		FA	BUFFD,99	08004	16	00469 -8039
17	428				08016	16	00445 J2174
17	428				08028	49	00422 -0099
17	430		TF	OUTPUT-36,BUFFN,, X.COS*THETA-Y.SIN*THETA	08040	26	12927 12164
17	432		TF	OUTPUT-26,BUFFD,, Y.COS*THETA-EX.SIN*THETA	08052	26	12917 12174
17	434		CM	OUTPUT-46,16,10,TEST FOR A SLOPE	08064	14	12917 000J6
17	436		RNE	*662	08076	47	08130 01200
17	438		TFH	*823,OUTPUT-16,,ALGEBRAICALLY SUM ANGLES	08088	16	08111 J2947
17	440		FA	OUTPUT-16,ROTO-1	08100	16	00469 -8135
17	440				08112	16	00445 J2947
17	440				08124	49	00422 -7681
17	442		RR		08136	42	00000 00000
17	444		DORG	*-10	08137		
17	446		CM	OUTPUT-46,14,10,TEST FOR AN ARC OR A SLOPE	08138	14	12917 000J4
17	448		RL	*-14	08150	47	08136 01300
17	450		TFH	*-39,OUTPUT-6	08162	16	08123 J2957
17	452		R	*-74	08174	49	08100 00000
17	454		DORG	*-2	08183		
17	500			REVERSE SUBROUTINE			
17	504	REVR5	SM	PTCNT,1,10	08184	12	12182 000-1
17	505		BVR	*-12,PTCNT,11,INDEX POINTER BACKWARDS	08196	45	08184 1218K
17	505		AN	PTCNT,2,10	08208	11	12182 000-2
17	507		CM	PTCNT,26,610,TEST FOR A GOTO OR A STOP	08220	14	1218K 000K6
17	508		BL	ND39	08232	47	08482 01300
17	509		BM	*832,OUTPUT-50	08244	43	08276 12913
17	509		SF	NORLV	08256	32	08688 00000
17	509		R	*820	08268	49	08288 00000
17	509		DORG	*-4	08275		
17	509		SF	NORVEL	08276	32	08689 00000
17	509		SM	PTCNT,1,10	08288	12	12182 000-1
17	510		TR	OUTPUT-51,PTCNT,11,SHIFT 4 LOCATIONS LEFT	08300	31	12912 1218K
17	511		BD	DEIN2,GND,TEST FOR FIRST GEN PAT	08312	43	08468 13515
17	512	GOST	CM	OUTPUT-46,11,10,TEST FOR Z COORD	08324	14	12917 000J1
17	513		NE	DEIN2	08336	46	08468 01200
17	514		FA	OUTPUT-36,ST-11	08348	16	00469 -8383
17	514				08360	16	00445 J2927
17	514				08372	49	00422 J2900
17	515		FA	OUTPUT-26,ST-1	08384	16	00469 -8419

PG	LN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
17	515				08396	16	00445 J2937
17	515				08408	49	00422 J2910
17	516		CM	OUTPUT-50,,10	08420	14	12913 000-0
17	517		BE	DEIN2	08432	46	08468 01200
17	518		TR	48,OUTPUT-51,,SHIFT 4 LOCATIONS RIGHT	08444	31	00048 12912
17	519		TR	OUTPUT-47,48	08456	31	12916 00048
17	520	DEIN2	SH	PTCNT,2,10,INDEX POINTER	08468	12	12182 000-2
17	521		DR		08480	42	00000 00000
17	522		DORG	*-10	08481		
17	523	ND37	SM	PTCNT,1,10	08482	12	12182 000-1
17	524		TR	OUTPUT-47,PTCNT,11,THIS POINT TO OUTPUT BUFFER	08494	31	12916 1218K
17	525		R	GOST-12	08506	49	08312 00000
17	556		DORG	*-2	08515		
17	598			* SUBROUTINE TRAN PERFORMS FINAL TRANSFORMATION			
17	600	TRAN	BNF	TRAN1,TRAN,0	08516	44	08588 08516
17	601		CF	TRAN,,,FIRST TIME SWITCH	08528	33	08916 00000
17	602		CM	INPUT-40,12,10,TEST FOR NEW FIRST POINT	08540	14	12032 000J2
17	603		BL	TRA3	08552	47	08958 01300
17	612		TF	ST-11,INPUT-30,,SAVE FIRST POINT	08564	26	12900 12042
17	614		TF	ST-1,INPUT-20	08576	26	12910 12052
17	604	TRAN1	CM	OPCR2,36,10,TEST FOR REVERSAL	08588	14	12802 000L6
17	604		BNE	NOREV	08600	47	08688 01200
17	604		AM	PTCNT,03,10	08612	11	12182 000-3
17	604		CM	PTCNT,11,610	08624	14	1218K 000J1
17	604		BNE	*632	08636	47	08668 01200
17	604		SM	PTCNT,03,10	08648	12	12182 000-3
17	604		B	POINT	08660	49	08784 00000
17	604		DORG	*-4	08667		
17	604		SM	PTCNT,03,10	08668	12	12182 000-3
17	604		B	NOREV24	08680	49	08712 00000
17	604		DORG	*-4	08687		
17	605	NOREV	CM	OUTPUT-40,11,10,TEST FOR Z COORD	08688	14	12917 000J1
17	608		BE	POINT	08700	46	08784 01200
17	617		TA	OUTPUT-30,ST-11	08712	16	00469 -8747
17	617				08724	16	00445 J2927
17	617				08736	49	00422 J2900
17	618		FA	OUTPUT-26,ST-1	08748	16	00469 -8783
17	618				08760	16	00445 J2937
17	618				08772	49	00422 J2910
17	622	POINT	BNF	*E20,NOREV	08784	44	08804 08688
17	623		R	*E44	08796	49	08840 00000
17	623		DORG	*-4	08803		
17	623		CM	OUTPUT-50,,10,TEST FOR GOTO OR STOP	08804	14	12913 000-0
17	624		BE	TRAN4	08816	46	08884 01200
17	625		AM	COMP266,04,10	08828	11	11414 000-4
17	625		BNF	*E20,NOR:V61	08840	44	08860 08689
17	625		R	TRAN4	08852	49	08884 00000
17	625		DORG	*-4	08859		

PG	LN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
17	626		TR	48,OUTPUT-51	08860	31	00048 12912
17	628		TR	OUTPUT-47,48,,RIGHT SHIFT 4 LOCATIONS	08872	31	12916 00048
17	639	TRA4	AM	COMP966,1,10	08884	11	11414 000-1
17	640		BNR	*-12,COMP966,11,INDEX FIRST GEN POINTER	08896	45	08894 11414
17	641		AM	COMP966,1,10	08908	11	11414 000-1
17	630		TFM	OUTPUT-50,,10,RESET SPECIAL SWITCHES	08920	16	12913 000-0
17	632		CF	NOREV	08932	33	08688 00000
17	634		CF	NOREV1	08944	33	08689 00000
17	664		BB		08956	42	00000 00000
17	645		DORG	*-10	08957		
17	648	TRA3	BNF	TRA8,PFLAG67,,TEST FOR 3RD GEN PATTERN	08958	44	09098 13269
17	649		BNF	TRA7,TRAN2,,ARE 2ND GEN PATTERNS LINKED TOGETHER	08970	44	09078 09614
17	649	TR10	BNF	*E20,D161,,ARE 1ST GEN PATS LINKED	08982	44	09002 09349
17	650		B	TRA6	08994	49	09022 00000
17	650		DORG	*-4	09001		
17	650		TR	ST-20,FG2X-9,,SAVE 1ST POINT OF 2ND GEN PAT	09002	31	12891 13430
17	651		R	TRAN1	09014	49	08588 00000
17	652		DORG	*-4	09021		
17	653	TRA6	CM	TRAN2-31,12,10,TEST FOR NEW 1ST PT OF THIS 1ST GEN PAT	09022	14	12808 000J2
17	654		BL	TRAN1	09034	47	08588 01300
17	657		TF	ST-11,TRAN2-21,,SAVE FIRST POINT	09046	26	12900 12018
17	658		TF	ST-1,TRAN2-11	09058	26	12910 12828
17	659		B	TRAN1	09070	49	08588 00000
17	650		DORG	*-4	09077		
17	651	TRA7	TR	ST-20,FG3X-9,,SAVE FIRST PT OF 3RD GEN PAT	09078	31	12891 13472
17	663		B	TRAN1	09090	49	08588 00000
17	664		DORG	*-4	09097		
17	665	TRA8	BNF	TRAN1,PFLAG68,,TEST FOR 2ND GEN PAT	09098	44	08588 13270
17	666		B	TR10	09110	49	08982 00000
17	667		DORG	*-2	09119		
17	900			* DTRAN SUBROUTINE COMPUTES DIFF TRANSFORMATION WHICH LINKS 1ST GEN 0			
17	904	DTRAN	BNF	D1,OTRAN,0,1ST TIME SWITCH	09120	44	09348 09120
17	905		BNF	OUT1,D161	09132	44	09290 09349
17	906		CF	D161	09144	33	09349 00000
17	907		CM	TRAN2-31,12,10,TEST FOR NEW 1ST FOR 1ST GEN PAT	09156	14	12808 000J2
17	908		BNL	OUT162	09168	46	09292 01300
17	912		TR	FG2X-9,ST-20	09180	31	13430 12891
17	912	TR61	BNF	OUT1,PFLAG67,,TEST FOR 3RD GEN PAT	09192	44	09290 13269
17	912		CM	INPUT-40,10,10,TEST FOR PAT OPERATION	09204	14	12032 000J0
17	912		BNL	OUT1	09216	47	09290 01200
17	912		CM	TRAN3-31,12,10,NO NEW 1ST PT FOR ENTIRE PAT	09228	14	12857 000J2
17	912		BE	SHIFT	09240	47	09278 01300
17	912		TF	FG3X,TRAN3-21	09252	26	13481 12867
17	912		TF	FG3Y,TRAN3-17	09264	26	13491 12877
17	912		BB		09276	42	00000 00000
17	912		DORG	*-10	09277		
17	912	SHIFT	TR	FG3X-9,FG2X-9,,RESET 1ST PT FOR PAT	09278	31	13472 13430
17	913	OUT1	BB		09290	42	00000 00000

PG LIN LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
17 914	DORG	*-10	09291		
17 915	TF	FG2X,TRAN2-21,,RESET 1ST PT	09292	26	13437 12818
17 916	TF	FG2Y,TRAN2-11	09304	26	13449 12828
17 918	TF	ST-11,TRAN2-21	09316	26	12900 12818
17 918	TF	ST-11,TRAN2-11	09328	26	12910 12828
17 919	B	TRK1	09340	49	09192 00000
17 920	DORG	*-4	09347		
17 921 D1	BNE	D4,D1,01,,1ST TIME SWITCH	09348	MM	09480 09348
17 922	CF	D1	09360	33	09348 00000
17 923	CM	TRAN2-31,12,10,TEST FOR NEW 1ST POINT	09372	14	12808 000J2
17 924	BSL	HLW1	09384	46	09580 01300
17 927	TR	NXG2X-9,ST-20,,SAVE 1ST PT OF NEXT 1ST GEN PAT	09396	31	13451 12891
17 929	FS	NXG2X,FG2X	09408	16	00469 -9443
17 929			09420	16	00445 J3460
17 929			09432	49	00402 J3439
17 930	FS	NXG2Y,FG2Y	09444	16	00469 -9470
17 930			09456	16	00445 J3470
17 930			09468	49	00402 J3449
17 939 FAD	CM	OUTPUT-46,11,10,TEST FOR Z COORD	09480	14	12917 000J1
17 939	BNE	*614	09492	47	09506 01200
17 939	BB		09504	42	00000 00000
17 939	DORG	*-10	09505		
17 939	FA	OUTPUT-36,NXG2X	09506	16	00469 -9541
17 939			09518	16	00445 J2927
17 939			09530	49	00422 J3460
17 940	FA	OUTPUT-26,NXG2Y	09542	16	00469 -9577
17 940			09554	16	00445 J2937
17 940			09566	49	00422 J3470
17 941	BB		09578	42	00000 00000
17 941	DORG	*-10	09579		
17 943 NEWT	TF	NXG2X,TRAN2-21,,SAVE NEW TRANSFORM	09580	26	13460 12818
17 944	TF	NXG2Y,TRAN2-11	09592	26	13470 12828
17 945	B	FAD-72	09604	49	09408 00000
17 946	DORG	*-2	09613		
17 948		* DTRAN2 SUBROUTINE COMPUTES DIFF TRAN BETWEEN 2ND GEN PATS			
17 950 DTRAN2	BNE	D4,DTRAN2,0,,1ST TIME SWITCH	09614	MM	09752 09614
17 951	CF	DTRAN2	09626	33	09614 00000
17 952	CM	TRAN2-31,12,10	09638	14	12857 000J2
17 953	BL	D12	09650	47	09688 01300
17 956	TF	FG3X,TRAN3-21,,SAVE NEW 2ND GEN PAT 1ST POINT	09662	26	13481 12867
17 957	TF	FG3Y,TRAN3-11	09674	26	13491 12877
17 958	BB		09686	42	00000 00000
17 959	DORG	*-10	09687		
17 960 DF2	CM	TRAN2-31,12,10	09688	14	12808 000J2
17 961	BL	D13	09700	47	09738 01300
17 964	TF	FG3X,TRAN2-21,,SAVE NEW 1ST GEN START POINT	09712	26	13481 12818
17 965	TF	FG3Y,TRAN2-11	09724	26	13491 12828
17 966	BB		09736	42	00000 00000

PG LIN LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
17 967	DORG	*-10	09737		
17 968 D13	TR	FG3X-9,ST-20,,TRANSFORM EQUALS 1ST POINT OF PAT	09738	31	13472 12891
17 970	BB		09750	42	00000 00000
17 971	DORG	*-10	09751		
17 972 D4	BNE	*614,D4,0	09752	MM	09766 09752
17 974	BB		09764	42	00000 00000
17 975	DORG	*-10	09765		
17 976	BNE	D5,D4E1	09766	44	09910 09754
17 977	CM	TRAN3-31,12,10	09778	14	12857 000J2
17 978	BL	D6	09790	47	10010 01300
17 981	TF	NXG3X,TRAN3-21,,SAVE 1ST POINT OF NEXT 2ND GEN PAT	09802	26	13502 12867
17 982	TF	NXG3Y,TRAN3-11	09814	26	13512 12877
17 983 D7	FS	NXG3X,FG3X	09826	16	00469 -9861
17 983			09838	16	00445 J3502
17 983			09850	49	00402 J3481
17 984	FS	NXG3Y,FG3Y	09862	16	00469 -9897
17 984			09874	16	00445 J3512
17 984			09886	49	00402 J3491
17 985	CF	D4E1	09898	33	09753 00000
17 922 D5	CM	OUTPUT-46,11,10,TEST FOR Z COORD	09910	14	12917 000J1
17 993	BNE	*614	09922	47	09936 01200
17 994	BB		09934	42	00000 00000
17 995	DORG	*-10	09935		
17 996	FA	OUTPUT-36,NXG3X	09936	16	00469 -9971
17 996			09948	16	00445 J2927
17 996			09960	49	00422 J3502
17 997	FA	OUTPUT-26,NXG3Y	09972	16	00469 J0007
17 997			09984	16	00445 J2937
17 997			09996	49	00422 J3512
17 998	BB		10008	42	00000 00000
17 999	DORG	*-10	10009		
18 000 D6	CM	TRAN2-31,12,10	10010	14	12808 000J2
18 001	BL	D8	10022	47	10066 01300
18 004	TF	NXG3X,TRAN2-21	10034	26	13502 12818
18 005	TF	NXG3Y,TRAN2-11	10046	26	13512 12828
18 006	B	D7	10058	49	09826 00000
18 007	DORG	*-4	10065		
18 008 D8	TR	NXG3X-9,ST-20	10066	31	13493 12691
18 010	B	D7	10078	49	09826 00000
18 040	DORG	*66	10095		
18 040		* SICO SUBROUTINE - COMPUTES SINE AND COSINE OF ROT ANGLE			
18 041 SICO	FM	RCUNV,SICO-T	10096	16	00469 J0131
18 041			10108	26	01260 13409
18 041			10120	49	01262 J0095
18 041			10132	26	00009 00099
18 042	TF	9,99,,ANGLE IN RADIAN	10144	16	14270 J0179
18 043	FSIN	HSIN,9	10156	16	14258 J3479
18 043			10168	49	13548 -0009

PG	LIN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
18	044		FCOS	THCRS,9	10180	16	14270 J0215
18	044				10192	18	14258 J3419
18	044				10204	49	13516 -0009
18	045		CF	TRIGF	10216	33	13267 00000
18	046		BR		10228	42	00000 00000
18	200		DORG	*-4	10235		
18	201	MILLR	CR	* MILLR SUBROUTINE PERFORMS INVERSIONS FOR MILLING OPERATIONS			
18	202		TR	TYPE,20,9,TEST FOR MILL OR FACE MILL	10236	14	12130 00-20
18	203		RR	FACE	10248	46	10980 01100
18	204		TR	OUTPUT-47,A10-42,,NEXT POINT TO OUTPUT BUFFER	10260	31	12916 13012
18	204		SM	PICNT,1,10,INDEX POINTER BACKWARDS	10272	12	12182 000-1
18	206		BHR	*-17,PICNT,11	10284	45	10272 1218K
18	208		AM	PICNT,1,10	10296	11	12182 000-1
18	210		RNF	M1,ADRS,,TEST FOR ARC OR SLOPE	10308	44	10340 1048K
18	212		TR	A10-42,PICNT,11,THIS POINT TO AUX BUFF	10320	31	13012 1218K
18	214		B	ADRS	10332	49	10484 00000
18	216		DORG	*-4	10339		
18	218	H1	TR	OUTPUT-47,PICNT,11, 1ST POINT TO OUTPUT BUFF	10340	31	12916 1218K
18	220		CM	OUTPUT-46,26,10,TEST FOR GOTO OR STOP	10352	14	12917 000K6
18	222		BL	M2	10364	47	10434 01300
18	224		RD	*632,OUTPUT-50,,TEST FOR GOTO OR STOP	10376	43	10408 12913
18	226		SF	NOREV	10388	32	08688 00000
18	228		B	*E20	10400	49	10420 00000
18	230		DORG	*-4	10407		
18	232		SF	NOREV&1	10408	32	08689 00000
18	234		TR	OUTPUT-51,OUTPUT-47,,LEFT SHIFT 4 LOCATIONS	10420	31	12912 12916
18	236		RR		10432	42	00000 00000
18	238		DORG	*-10	10433		
18	240	M2	CM	OUTPUT-46,14,10,SLOPE OR ARC CHECK	10434	14	12917 000J4
18	242		RNL	*E26	10446	46	10472 01300
18	244		SM	PICNT,1,10	10458	12	12182 000-1
18	246		BR		10470	42	00000 00000
18	248		DORG	*-10	10471		
18	250		BNF	M6,ADRS,,WAS LAST POINT A SLOPE OR AN ARC	10472	44	10924 1048K
18	252	ADRS	TR	OUTPUT-36,A10-31,,X AND Y TO OUTPUT BUFF	10484	26	12927 13023
18	254		TR	OUTPUT-26,A10-21	10496	26	12937 13033
18	256		CM	OUTPUT-46,16,10,WAS LAST PT A SLOPE	10508	14	12917 000J6
18	258		RNF	MFADE36	10520	47	10732 01200
18	260		CM	OUTPUT-21,53360,,TEST SIZE OF ANGLE	10532	14	12942 N3360
18	262		BL	*E56	10544	47	10600 01300
18	264		FS	OUTPUT-16,D360	10556	16	00469 J0991
18	264				10568	16	00445 J2947
18	264				10580	49	00402 J3399
18	266		B	*-60	10592	49	10532 00000
18	268		DORG	*-4	10599		
18	270		TFM	MFADE35,090	10600	16	10731 J3369
18	272		TFM	MFADE30,00422	10612	16	10726 -0422
18	274		CM	OUTPUT-21,53270,,TEST FOR ANGLE GREATER THAN 270 DEGREES	10624	14	12942 N3270

PG	LIN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
18	276		BNH	M3	10636	47	10800 01100
18	278		TFM	MFADE30,00402	10648	16	10726 -0402
18	280	M4	CM	MILLR-1,32,610,TEST FOR INVERT RIGHT	10660	14	1023N 000L2
18	282		RNF	*E24	10672	47	10696 01200
18	284		TFM	MFADE35,0270	10684	16	10731 J3389
18	186	MFADE	FA	OUTPUT-16,090	10696	16	00469 J0731
18	186				10708	16	00445 J2947
18	287		CM	A10-41,14,10,TEST FOR CW ARC	10720	49	00422 J3369
18	288		BL	*E36	10732	14	13013 000J4
18	289		CM	A10-41,26,10,TEST FOR GOTO OR STOP	10744	47	10780 01300
18	290		BL	M2&24	10756	14	13013 000K6
18	290		BL	M2&24	10768	47	10458 01300
18	191		CF	ADRS,,RESET ARC OR SLOPE SWITCH	10780	33	10484 00000
18	292		B	M2&24	10792	49	10458 00000
18	293		DORG	*-4	10799		
18	294	M3	CM	OUTPUT-21,52900,7,TEST FOR ANGLE LESS THAN 90 DEG	10800	14	12942 N2900
18	295		BL	M4	10812	47	10660 01300
18	296		CM	OUTPUT-21,53180,7,TEST FOR ANGLE GREATER THAN 180 DEG	10824	14	12942 N3180
18	298		RR	M5	10836	46	10892 01100
18	300		CM	MILLR-1,32,610,TEST FOR INVERT RIGHT OR LEFT	10848	14	1023N 000L2
18	302		BE	MFADE	10860	46	10696 01200
18	304		TFM	MFADE30,00402	10872	16	10726 -0402
18	306		B	MFADE	10884	49	10696 00000
18	308		DORG	*-4	10891		
18	310	M5	CM	MILLR-1,34,610,TEST FOR INVERT UP OR DOWN	10892	14	1023N 000L4
18	312		BE	MFADE	10904	46	10696 01200
18	314		B	M5-20	10916	49	10872 00000
18	316		DORG	*-4	10923		
18	318	M6	SF	ADRS,,SET ARC OR SLOPE SWITCH	10924	32	10484 00000
18	320		TR	A10-42,OUTPUT-47,,NEXT PT TO AUX BUFF	10936	31	13012 12916
18	322		TR	OUTPUT-25,AUXIN,,SET CONTROL REC MARK	10948	31	12938 12104
18	324		TFM	OUTPUT-46,12,10,SET COORD CODE	10960	16	12917 000J2
18	326		B	M2&24	10972	49	10458 00000
18	328		DORG	*-4	10979		
18	320	FACE	RNF	*E38,FACE,,TEST FOR FACE MILL	10980	44	11018 10980
18	322		TR	OUTPUT-47,A10-42,,THIS POINT TO OUTPUT BUFF	10992	31	12916 13012
18	324		CF	FACE	11004	33	10980 00000
18	326		RR		11016	42	00000 00000
18	328		DORG	*-10	11017		
18	330		TR	A10-42,OUTPUT-47,,SHIFT PT TO AUX BUFF	11018	31	13012 12916
18	332		TF	9,COORD&6	11030	26	00009 11414
18	332		AM	9,1,10	11042	11	00009 000-1
18	333		BHR	*-12,9,11,INDEX 1ST GEN POINTER	11054	45	11047 0000R
18	334		AM	9,1,10	11066	11	00009 000-1
18	335		TR	OUTPUT-47,9,11,NEXT POINT TO OUTPUT BUFF	11078	31	12916 0000R
18	336		BHR	*E20,A10-20,,TEST FOR DEPTH	11090	45	11116 13034
18	337		B	*E56	11102	49	11158 00000
18	338		DORG	*-4	11109		



PG	LN	LABEL	OP	OPR ANDS	LOCN	OP	INSTRUCTIONS
18	339		TR	OUTPUT-25,AID-20,,SHIFT DEPTH TO 1ST PT IN PAT	11110	31	12938 13034
18	340		TR	AID-20,AUX0,,SET CONTROL REC MARK	11122	31	13034 13011
18	341		TR	AID-41,2,,,XCHANGE COORD CODES	11134	15	13013 00002
18	342		TR	OUTPUT-46,1	11146	15	12917 00003
18	343		SE	FAIL	11158	32	10980 00000
18	344		ER		11170	42	00000 00000
18	345		DRG	#-10	11171		
16	170	PLONG	TR	ERSUR,7600,8	11172	17	11190 0P400
16	171	MESSG	PAC	3,,	11135		3,3,E
16	172	ERSUR	TR	#-1,AUXIN	11190	31	11189 12104
16	173		WATY	MESSG	11202	39	11185 00100
16	174		B	,,,PUSH START TO CONTINUE	11214	48	00000 00000
16	175		TR	OUTPUT-47,INPUT-47	11226	31	12916 12025
16	176		RTM	DPUI,,,OUTPUT THIS REC	11238	17	11484 -0000
16	177		TR	INPUT,AUXIN	11250	31	12072 12104
16	178		TR	OUTPUT,AUX0	11262	31	12963 13011
16	179		B	NEXT	11274	49	01700 00000
19	010		DRG	#62	11287		
19	014			*SUBROUTINE SP-STORES VARIABLE REC IN PAT TABLE			
19	015	SP	TR	PTAP,#-1,611,STORE THIS RECORD IN PATTERN TABLE	11288	31	1326- 1128P
19	020		AM	PTAP,1,10	11300	11	13260 000-1
19	025		AM	PFLAGE4,1,10	11312	11	13266 000-1
19	030		CM	PTAB,19999,7,TEST FOR END OF PAT STORAGE	11324	14	13260 J9999
19	035		BH	PATEX	11336	46	11382 01100
19	040		BNR	#614,SP-1,11	11348	45	11362 1128P
19	045		BH		11360	42	00000 00000
19	050		DRG	#-10	11361		
19	054		AM	SP-1,1,10,INDEX LENGTH OF REC POINTER	11362	11	11287 000-1
19	060		B	SP612	11374	49	11300 00000
19	062		DRG	#-4	11381		
19	065	PATEX	RTM	ERSUR,7990,8,ALLOWABLE PAT STORAGE EXCEEDED	11382	17	11190 0P300
19	070		DRG	#62	11395		
20	000			* SUBROUTINE PATLOC-SCANS PAT TABLE TO LOCATE PAT TO BE OUTPUT			
20	030	PATLOC	TRM	COMP96,GN02,99,,RESET TO START OF PAT TABLE	11396	16	11414 J4509
20	040	COMP9	C	00000,PATLOC-1,,TEST FOR PAT NO	11408	24	00000 11395
20	050		BNI	#614	11420	47	11434 01200
20	060		BH		11432	42	00000 00000
20	062		DRG	#-10	11433		
20	070		TR	ADD9611,COMP966	11434	26	11469 11414
20	080		AM	ADD9611,0%,10	11446	11	11469 000-5
20	090	ADD9	A	COMP96,00000,,MOVE POINTER TO FIRST NEXT PAT IN PAT TABLE	11458	21	11414 00000
20	100		B	COMP9	11470	49	11408 00000
21	012		DRG	#62	11483		
21	013			*SUBROUTINE DPUI---OUTPUTS CARDS OR TAPE W/ OR W/O TYPED LISTING			
21	040	DPUI	PC2	PNCHO,,,TEST FOR CARD OUTPUT	11484	46	11586 00200
21	016		WPT	OUTPUT-47,,,PAPER TAPE OUTPUT	11496	38	12916 00200
21	017		BC4	#620	11508	46	11528 00400
21	018		B	PRINT	11520	49	11554 00000

PG	LN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
21	019		DRG	#-4	11527		
21	020	CK99	CM	OUTPUT-46,99,10,TEST FOR END OF PROGRAM	11528	14	12917 000R9
21	021		BE	FINI	11540	46	12006 01200
21	022	TR001	BR		11552	42	00000 00000
21	023		DRG	#-10	11553		
21	024	PR011	WPT	OUTPUT-47,,,OUTPUT LISTING ON TYPEWRITER	11554	38	12916 00100
21	026		INTY		11566	34	00000 00108
21	028		B	CK99	11578	49	11578 00000
21	030		DRG	#-4	11585		
21	032	PNCHO	AM	RSIZE,01,10,CARD OUTPUT	11586	11	13245 000-1
21	034		BNR	PNCHO,RSIZE,11,COUNT LENGTH OF THIS RECORD	11598	45	11586 1324N
21	036		AM	RSIZE,01,10	11610	11	13245 000-1
21	038		SM	RSIZE,OUTPUT-47	11622	12	13245 J2916
21	040		A	OUTLOC,RSIZE	11634	21	13235 13245
21	042		CH	OUTLOC,CARD0-5,,TEST FOR FULL OUTPUT BUFFER	11646	14	13235 J3210
21	044		RT	NOFIT	11658	46	11690 01100
21	046		S	OUTLOC,RSIZE	11670	22	13235 13245
21	048		B	FILL	11682	49	11914 00000
21	050		DRG	#-4	11689		
21	052	NOFIT	S	OUTLOC,RSIZE	11690	22	13235 13245
21	054	DIG0	TRM	JULOC,00300,6,RESET REMAINDER OF CARD TO ZERO	11702	15	1323N 00000
21	056		AM	OUTLOC,01,10	11714	11	13235 000-1
21	058		CM	OUTLOC,CARD0-6	11726	14	13235 J3211
21	053		DI	DIG0	11733	47	11702 01300
21	052		AM	SEQNO,01,10	11750	11	13250 000-1
21	054		TR	CARD0,SEQNO,,SET SEQUENCE NO	11762	26	13215 13250
21	066		WNGD	CARD0-79,,PUNCH ONE CARD	11774	38	13134 00400
21	067	SEQNO	TR	REC01,,,TEST FOR LISTING	11786	46	11890 00400
21	071	PR011	BNE	REC11,PRLOC,11	11798	44	11878 1324-
21	072		WNTY	PRLOC,,6,OUTPUT THIS REC ON TYPEWRITER	11810	38	1324- 00100
21	072		INTY		11822	34	00000 00108
21	073		AM	PRLOC,01,10	11834	11	13240 000-1
21	074		BNR	#-12,PRLOC,11	11846	45	11834 1324-
21	075		AM	PRLOC,01,10	11858	11	13240 000-1
21	076		B	PR011	11870	49	11798 00000
21	077		DRG	#-4	11877		
21	078	REC11	TRM	PRLOC,CARD0-79,,RESET POINTERS	11878	16	13240 J3136
21	081	REC00	TRM	OUTLOC,CARD0-79	11890	16	13235 J3136
21	082		BNE	REC11,TRIGET	11902	44	11938 13268
21	083	FILL	TR	OUTLOC,OUTPUT-47,6,THIS REC TO OUTPUT BUFFER	11914	31	1323N 12916
21	084		A	OUTLOC,RSIZE,,INDEX OUTPUT BUFF REGISTER	11926	21	13235 13245
21	085	REC11	TRM	RSIZE,OUTPUT-47	11938	16	13245 J2916
21	086		CM	OUTPUT-46,99,10,TEST FOR END OF PROGRAM	11950	14	12917 000R9
21	087		BE	FINI	11962	47	11552 01200
21	088		WPT	FINI,TRIGET	11974	44	12006 13268
21	089		CT	TRIGET	11986	33	13268 00000
21	090		B	DIG,,,PUNCH LAST CARD	11998	49	11702 00000
21	091		DRG	#-4	12005		

49-19972

PG	LIN	LABEL	OP	OPERANDS	LUCN	OP	INSTRUCTIONS
21	098	FINI	HATY	MESS1...END OF PHASE2	12006	39	13339 00100
21	100		R	INIT...INITIALIZE FOR NEXT JOB	12010	49	00050 00000
21	101		ORGR	#-4	12025		49 19522
1	010	INPUT	DS	48	12072		48
1	015		DC	1,3,INPUT	12072		1
1	016	AUXIN	DS	32	12104		32
1	017		DC	1,3,AUXIN	12104		1
1	020	EXTY	DS	21	12125		21
1	022	RTYPE	DC	2,00	12127		2 -0
1	030	TYPE	DS	3	12130		3
1	040	FOLNO	DS	4	12134		4
1	041	TWO	DC	10,5170000000	12144		10 N120000000
1	042	THETA	DS	10	12154		10
1	043	BUFFN	DS	10	12164		10
1	044	BUFFD	DS	10	12174		10
1	046	TJEMOD	DS	3	12177		3
1	046	PTCNT	DC	5,00000	12182		5 -0000
1	046	RTCNT	DC	5,00000	12187		5 -0000
1	047	TBLE	DS	560	12747		560
1	050	COMPL	DSA	#67	12752		5 J2754
1	055		DC	2,17	12754		2 J7
1	060		DC	2,18	12756		2 J8
1	070		DC	2,19	12758		2 J9
1	080		DC	2,21	12760		2 K1
1	090		DC	2,22	12762		2 K2
1	095		DC	2,27	12764		2 K7
1	099		DC	2,28	12766		2 K8
1	099		DC	2,29	12768		2 K9
1	100		DC	2,90	12770		2 R0
1	104		DC	2,99	12772		2 R9
1	110	COMP2	DSA	#67	12777		5 J2779
1	113		DC	2,32	12779		2 L2
1	116		DC	2,34	12781		2 L4
1	120		DC	2,36	12783		2 L6
1	124		DC	2,40	12785		2 M0
16	139	TYLOC	DC	5,00000	12790		5 -0000
16	141	PLOC2	DC	5,00000	12795		5 -0000
16	141	PCNT2	DC	5,00000	12800		5 -0000
16	142	OPER2	DC	2,00	12802		2 -0
16	143	SORT2	DC	2,00	12804		2 -0
16	143		DC	2,00	12806		2 -0
16	144	TRAN2	DS	33	12839		33
16	145	PLOC3	DC	5,00000	12844		5 -0000
16	146	PCNT3	DC	5,00000	12849		5 -0000
16	147	OPER3	DC	2,00	12851		2 -0
16	148	SORT3	DC	2,00	12853		2 -0
16	148		DC	2,00	12855		2 -0
16	149	TRAN3	DS	33	12888		33

PG	LIN	LABEL	OP	OPERANDS	LUCN	OP	INSTRUCTIONS
16	145	PAIRO	DC	2,00	12890		2 -0
16	146	SF	DS	21	12911		21
16	148		DC	2,00	12913		2 -0
16	149		DC	2,00	12915		2 -0
16	150	OUTPUT	DS	48	12963		48
16	151		DC	1,3,OUTPUT	12963		1
16	152	AUXD	DS	48	13011		48
16	153		DC	1,3,AUXD	13011		1
16	153	AID	DS	43	13054		43
16	153		DC	1,3	13055		1
16	154	CARD1	DS	80	13135		80
16	155	CARD0	DS	80	13215		80
16	156	IGLOC	DSA	CARD1-79	13220		5 J3056
16	157	TABPI	DSA	TBLE-556	13225		5 J2191
16	157	ADRES	DSA	TBLE-556	13230		5 J2191
16	159	OUTLOC	DSA	CARD0-79	13235		5 J3136
16	162	PRLOC	DSA	CARD0-79	13240		5 J3136
16	164	PSIZE	DSA	OUTPUT-47	13245		5 J2918
16	166	SECHO	DC	5,00000	13250		5 -0000
16	167	SORT1	DC	5,00000	13255		5 -0000
16	168	PIAG	DSA	GNOE933	13260		5 J4508
16	168	PIAG	DC	2,00	13262		2 -0
16	168	TRIG	DC	5,-00000	13267		5 -000-
16	168		DC	6,000000	13273		6 -00000
16	168	PCNT	DC	2,00	13275		2 -0
16	168	PLIV	DC	2,00	13277		2 -0
16	168		DC	1,3	13278		1
16	169	TITLE	DSA	TBLE-556,TBLE-556,CARD0-79,CARD0-79,OUTPUT-47,,GNOE933	13283		5 J2191
16	169				13293		5 J3136
16	169				13303		5 J2918
16	169				13313		5 -0000
16	169				13325		5 -000-
16	169		DC	5,-00000	13331		5 -00000
16	169		DC	6,000000	13333		6 -00000
16	169		DC	2,00	13336		2 -0
16	169		DC	3,000	13336		3 -0
16	170	MESS1	SA	11,END PHASE2	13339		11 11,END PHASE2
16	172	D30	DC	10,5290000000	13369		10 N290000000
16	174	E170	DC	10,5110000000	13379		10 N318000000
16	175	E270	DC	10,5170000000	13389		10 N327000000
16	178	E360	DC	10,5360000000	13399		10 N336000000
18	810	RC170	DC	10,5217000000	13409		10 M917453293
18	825	TRC1	DC	10	13419		10
18	829	TRC1	DC	10	13429		10
18	830	TRC1	DC	10	13439		10
18	832	TRC1	DC	10	13449		10
18	834	TRC1	DC	10	13459		10
18	836	TRC1	DC	10	13460		10
18	837	TRC1	DC	10	13460		10

PG	LN	LABEL	OP	OPERANDS	LOCN	OP	INSTRUCTIONS
18	898	NXG2Y	DS	10	13470	10	
18	899		DC	1,0	13471	1	
18	900	FGJX	DS	10	13481	10	
18	901	FGJY	DS	10	13491	10	
18	902		DC	1,0	13492	1	
10	902	NXG1X	DS	10	13502	10	
18	903	NXG1Y	DS	10	13512	10	
18	904		DC	1,0	13513	1	
18	905	GND	DC	2,00	13515	2	-0
6	110		END	NEXT	01700		

MODIFICATION TO PHASE 2 OF AUTOSPOT TO  
RUN FROM DISK

PATCH PROGRAM

19466	26	13225	12033	
19478	11	13225	00010	
19490	26	13225	12055	
19502	11	13225	00004	Program patch that was in the way moved to here
19514	49	11226	0	
19522	34	19560	00701	End of this phase - seek for next phase
19534	49	19976	06850	
19546	10540020000000			Disk control word this section
19560	10535001001900			Disk control word next section
19574	bbbbbbbbbb			

MODIFICATION TO ORIGINAL PROGRAM

19584	36	00000	00500	Read in last cards of modification
19596	36	19680	00500	
19608	26	01762	19712	Modify "read a card" instruction
19620	26	03219	19705	Modify error routine
19632	41	00000	00000	No Op
19644	26	11780	19720	Modify "write a card" instruction
19656	26	12024	19727	Modify end to call next program
19668	34	19546	00701	
19680	38	19546	00702	Write these modifications all on the disk
19692	48			
19694	46	19466	01200	
19706	49	197420		
19714	49	19912	4919522	

READ DISK DATA

19728	10000000119600			
19742	36	19728	00702	Read a card from disk & seek if necessary then go back and read again
19754	47	19786	03600	
19766	34	19728	00701	
19778	49	19742	0	
19786	11	19733	00001	Increase sector address by one

TRANSFER FIELD JUST READ TO EVEN LOCATION

19798	25	13056	19600	Transfer data to area beginning with an even address
19810	14	19809	19679	
19822	46	19866	01200	Check for last transfer
19834	11	19809	00001	Increase disc location by one
19846	11	19804	00001	
19858	49	19798	0	Repeat

# MODIFICATIONS TO PHASE 2 of AUTOSPOT

## TO RUN FROM DISK

### TRANSFER FIELD JUST READ TO EVEN LOCATION

19866	16	19809	19600	Housekeep because all data is transferred now
19878	16	19804	13056	
19890	49	01768	0	Return to program

### WRITE DATA ON DISK

19898		10200000	113136	
19912	38	19898	00702	Write data on disk
19924	47	19956	03600	If address check - seek first then write data
19936	34	19898	00701	
19948	49	19912	0	
19956	11	19903	00001	Increase sector address by one
19968	49	11786	0	Return to program
19976	36	19560	00702	End of this phase - read in control program for next phase
19988	49	02318	00000	

### Load program to load in modifications into core

00000	36	19466	00500	Load modifications into core and branch to the first modification
00012	36	19546	00500	
00024	36	19626	00500	
00036	36	19760	00500	
00048	36	19840	00500	
00060	36	19920	00500	
00072	49	19584	0	





**GOOD YEAR**  
**GOODYEAR AEROSPACE**  
CORPORATION  
ARIZONA DIVISION  
LITCHFIELD PARK, ARIZONA

GENERAL RAY TRACE PROGRAM

Presented to the 1620 Users Group  
(Western Region) Meeting at Tempe,  
Arizona, December 12, 1963.

Presented by D. H. O'Herren

AAP-18375

December 12, 1963



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GENERAL RAY TRACE PROGRAM

This program is based upon a paper by Gordon H. Spencer titled "A General Ray Tracing Procedure", IBM Research Paper RC-549. Spencer's paper applies, with a few extensions, to the problem of tracing a light ray thru surfaces which may be rotationally symmetric, cylindrical, or conic and may have arbitrary orientations with respect to a reference system. Certain adjustments and additions have been included in the program. This paper does not attempt to delve deeply into the mathematics behind the program. It is felt that a brief description of some of its advantageous points along with a sample ray trace problem would be of more interest. The sample will attempt to illustrate the value of the program as an optics system evaluation tool.

In general, surfaces are described in a local coordinate system. This local system is then positioned in the optical system thru translation and rotation of the local system in relation to the basic coordinate system. It is thereby possible to position surfaces off the optical axis and "tilted" to the desired degree.

One feature the program possesses which is not common to some other ray trace programs is the capability to trace rays thru cylinders and prisms. These types of lenses are fairly common in modern optics systems making this an important feature.

The basic surface equation suggested by Spencer which is used in the program is:

$$F(X, Y, Z) = AX^2 + BY^2 + CZ^2 + Z = 0.$$

This represents the surface obtained by revolving a conic section with vertex at the origin about the Z axis. In general, a ray intersects such surfaces at two points. By using vertex equations set up so that the first iteration point lies on the plane  $Z = 0$  thru the vertex, the iteration starts closer to the desired

intersection point than to the extraneous point.

For cylinders with conic cross-section in the XZ plane, the surface equation becomes:

$$F = AX^2 + CZ^2 + Z = 0$$

The rulings of this cylinder are parallel to the Y axis. A similar equation will describe a cylinder with rulings parallel to the X axis.

The equation:

$$\frac{X^2}{(aY)^2} + \frac{Z^2}{(bY)^2} = 1$$

describes a cone with apex at the origin with the Y axis as the principal axis.

Cross-sections parallel to the XZ plane are ellipses. This equation may be rewritten in the form:

$$F = Z + D \sqrt{EX^2 + Y^2} = 0$$

To iterate with this equation, the following rule holds:

$$Z \geq 0 \text{ for } D < 0$$

$$Z \leq 0 \text{ for } D > 0$$

Therefore it is possible to represent three types of surfaces with the single equation:

$$F = AX^2 + BY^2 + CZ^2 + Z + D \sqrt{EX^2 + Y^2} = 0$$

The constants A, B, C, D, and E are program input data.



SURFACE CARDS:

Card A:

- XO(I)            Cols. 1 - 11    X - coordinate of origin of local system, punched in same fashion as the initial index of refraction on the header card.
- YO(I)            Cols. 12 - 22   Y - coordinate of origin of local system, punched in same fashion as X - coordinate
- ZO(I)            Cols. 23- 33    Z - coordinate of origin of local system, punched in same fashion as X - coordinate.
- Alpha(I)        Cols. 34 - 44   Y axis euler angle ( $\alpha$ ) in decimal degrees, punched in same fashion as X - coordinate.
- Beta(I)          Cols. 45 - 55   X axis euler angle ( $\beta$ ) in decimal degrees, punched in same fashion as X - coordinate.
- Gamma(I)        Cols. 56 - 66   Z axis euler angle ( $\gamma$ ) in decimal degrees, punched in same fashion as X - coordinate.
- Cols. 67 - 77   Blank.
- Cols. 78 - 80   It is suggested, but not required, that the surface cards be punched 01A, 01B, 01C, 02A, 02B, etc., in Cols. 78 - 80 to ensure that they are kept in the correct order.

Card B:

The first five fields give the coefficients in the surface equation  $F = AX^2 + BY^2 + CZ^2 + Z + D \quad EX^2 + Y^2 = 0$ , as follows:

- A(I)            Cols. 1 - 11, A (punched in same fashion as X - coordinate on Card A)
- B(I)            Cols. 12 - 22B (punched in same fashion as A).
- C(I)            Cols. 23 - 33C (punched in same fashion as A).
- D(I)            Cols. 34 - 44D (punched in same fashion as A).
- E(I)            Cols. 45 - 55E (punched in same fashion as A).

Refr(I)                    Cols. 56 - 66 Index of refraction of medium following the surface, except in the case of reflection, when the negative of the index of refraction for the previous surface is used. (Field is punched in same fashion as A).

                          Cols. 67 - 77 Blank

                          Cols. 78 - 80 May be punched as suggested for Card A.

Card C:

AP1(I)                    Cols. 1 - 11 X - coordinate of center of circular-annular aperture or coordinate of center of hyperbolic aperture or X lower bound of base of rectangular-trapezoidal aperture.

AP2(I)                    Cols. 12 - 22 Y - coordinate of center of circular-annular aperture or coordinate of center of hyperbolic aperture or X upper bound of base of rectangular-trapezoidal aperture.

AP3(I)                    Cols. 23 - 33 Inner radius of circular-annular aperture or length of semi-major axis of hyperbolic aperture or Y lower bound for rectangular-trapezoidal aperture.

AP4(I)                    Cols. 34 - 44 Outer radius of circular-annular aperture or length of semi-minor axis of hyperbolic aperture or Y upper bound for rectangular-trapezoidal aperture.

AP5(I)                    Cols. 45 - 55 Y lower bound for hyperbolic aperture or reciprocal slope of left hand side of rectangular-trapezoidal aperture. Enter 0<sub>1</sub> for circular-annular aperture.

AP6(L)                    Cols. 56 - 66    Y upper bound for hyperbolic aperture or reciprocal slope of right hand side of rectangular-trapezoidal aperture. Enter 0 for circular-annular aperture.

                          Cols. 67            Blank

NAP (I)                    Col. 68            Aperture code

    Blank: circular - annular

    1:        rectangular-trapezoidal

    2:        hyperbolic

                          Col. 69            Blank

NOUT(I)                    Col. 70            Output code

    Blank: no output at surface

    1:        output at surface

                          Cols. 71 - 77    Blank

RAY INPUT (Single ray per card, switch 1 on)

One record, either typed or punched, is used for each ray. NOTE: Coordinates with a bar over them are system coordinates. Those without a bar are local coordinates.

XA                        Cols. 1 - 11     $\bar{X}$  - coordinate of 1st point on ray.

YA                        Cols. 12 - 22    $\bar{Y}$  - coordinate of 1st point on ray.

ZA                        Cols. 23 - 33    $\bar{Z}$  - coordinate of 1st point on ray.

RPAR 1                    Cols. 34 - 44    $\bar{X}$  - coordinate of 2nd point on ray.  
    or  $\bar{X}$  direction cosine of ray at 1st point.

RPAR 2                    Cols. 45 - 55    $\bar{Y}$  - coordinate of 2nd point on ray  
    or  $\bar{Y}$  direction cosine of ray at 1st point.

RPAR 3                    Cols. 56 - 66    $\bar{Z}$  - coordinate of 2nd point on ray  
    or  $\bar{Z}$  direction cosine of ray at 1st point.

                          Col. 67            Blank





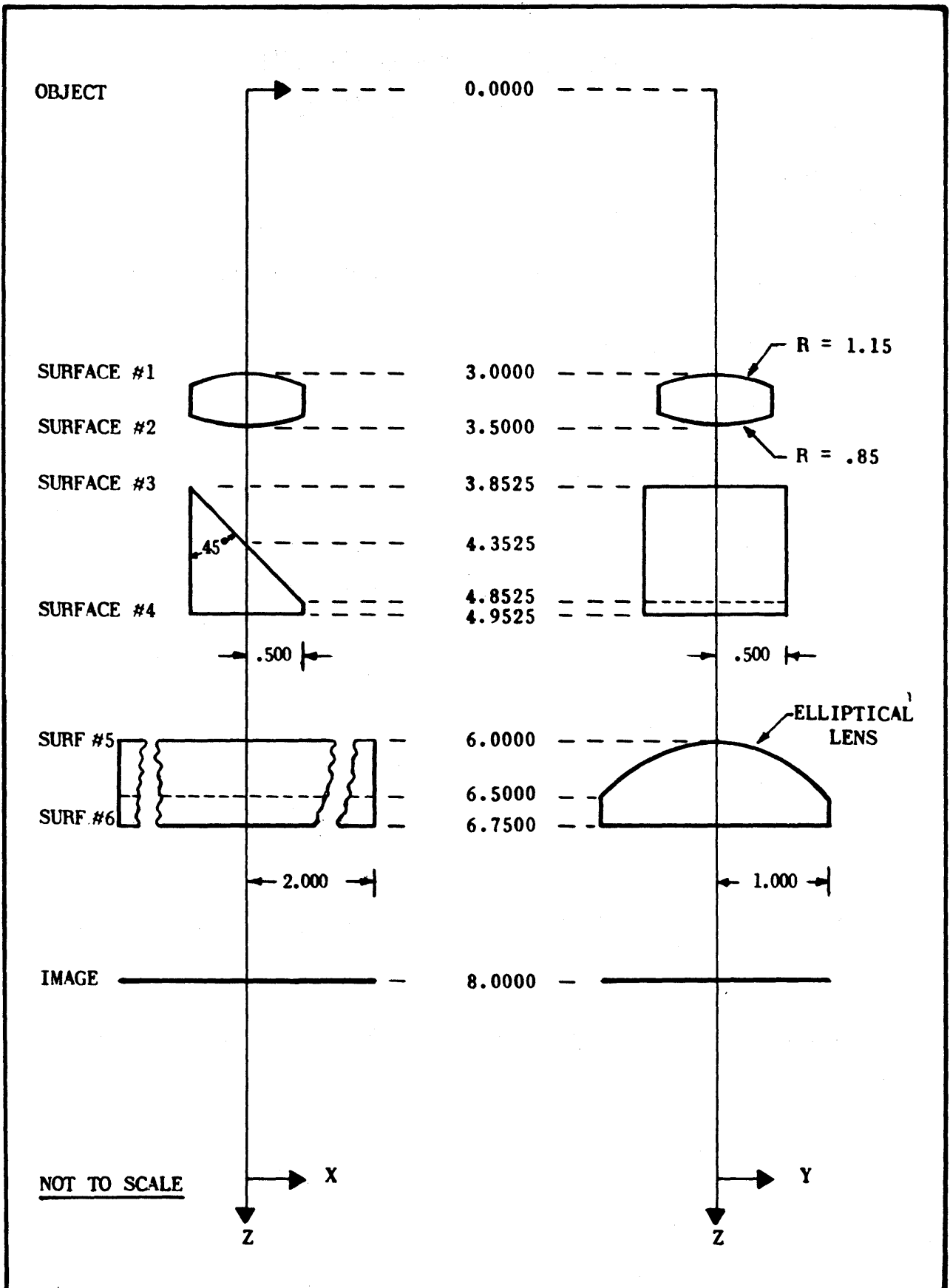


FIGURE - SAMPLE RAY TRACE SYSTEM

DATA INPUT TO GENERAL RAY TRACE

PROGRAM FOR SAMPLE PROBLEM

SURFACE DATA

SURFACES 1 & 2

These are both spherical surfaces. The surface equation becomes:

$$F = AX^2 + BY^2 + CZ^2 + Z = 0$$

$$= \frac{1}{2R} X^2 + \frac{1}{2R} Y^2 + \frac{1}{2R} Z^2 + Z = 0$$

$$= -.435 X^2 - .435 Y^2 - .435 Z^2 + Z = 0 \quad (\text{Surface \#1})$$

$$= .588 X^2 + .588 Y^2 + .588 Z^2 + Z = 0 \quad (\text{Surface \#2})$$

These are not cones. Therefore  $D = E = 0$

The vertex planes for these surfaces are perpendicular to the Z (optical) axis. Therefore  $\alpha$ ,  $\beta$ , and  $\gamma$  are zero.

The vertex coordinates are:

Surface #1           (0, 0, 3.00)

Surface #2           (0, 0, 3.50)

The index of refraction following Surface #1 is 1.523, that following Surface #2 is 1.00.

The apertature on both surfaces is a circle of 0.5 inches diameter.

Surfaces 3 & 4

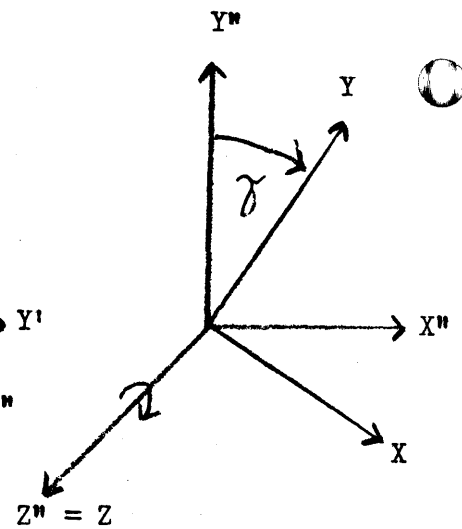
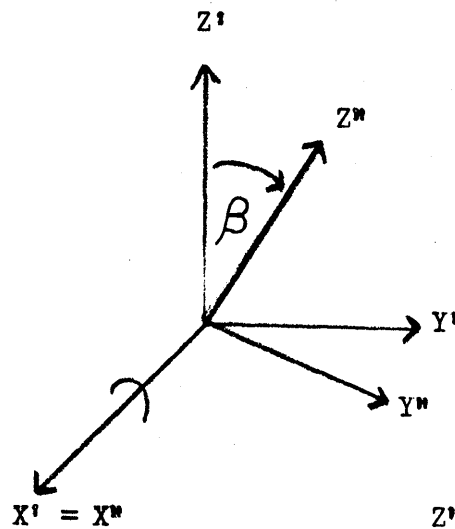
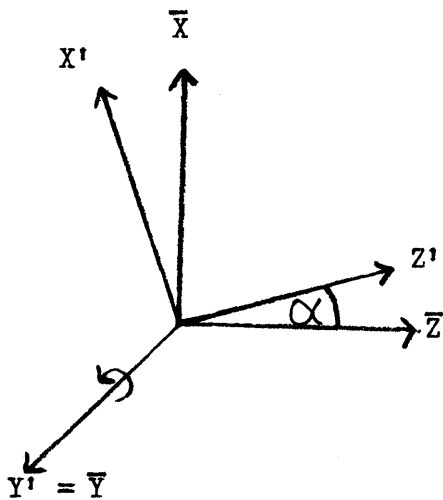
Both surfaces are planes arranged to constitute a prism. The first, No. 3, is arranged so that rays are deflected in the negative X direction. The second surface, No. 4, is perpendicular to the optic axis.

The surface equation for both surfaces is:

$$F = Z = 0$$

$$A = B = C = D = E = 0$$

Surface No. 3 is rotated 45 degrees when positioned in the system. To accomplish this rotation, it is necessary to specify the angle alpha equal to 45 degrees. The rotation angles, alpha, beta, and gamma are defined:



The specified apertature is a square 1" x 1". Therefore the minimum and maximum allowable x and y values are (0.5, 0.5) and (-0.5, 0.5) respectively. These values apply to the surfaces before rotation. After rotation of surface No. 3, the actual minimum and maximum x apertature values will be corrected by the program to  $(-0.5/\sqrt{2}, 0.5/\sqrt{2})$ . AP5 and AP6 are the inverse slopes of the left and right edges of the apertature as viewed from the object point (positive Y up). In the sample, these slopes are reciprocal infinity or zero. NAP is set equal to 1 to specify a rectangular apertature. NOUT is set equal to 1 to obtain output data at the surfaces. The index of refraction of the prism is 1.6.

#### Surface 5 & 6

Surface 5 is an elliptical cylinder. Surface 6 is a simple plane with input similar to surface No. 4. The equation of the ellipse in the YZ plane is:

$$\frac{Y^2}{a^2} + \frac{Z^2}{b^2} = 1 \quad a = 1 \quad b = 0.5$$

The surface equation, adjusted so that the origin is at the point nearest the object point, is:

$$F = -0.25 Y^2 - Z^2 + Z = 0$$

Therefore  $A = D = E = 0$ ,  $B = -0.25$ ,  $C = -1.00$ . Note that surface No. 5 is a cylinder with rulings in the X direction. If it had been desired to translate this surface off the optic axis in the Y direction, then YO would be specified accordingly.

This program was originally written by William Webb of Goodyear Aerospace Corporation, Akron, Ohio and to that gentlemen goes the credit for this significant contribution to the lens designer's kit of tools. The program has been altered slightly to suit current needs and is being maintained by the writer.

SAMPLE PROBLEM DATA INPUT

SAMPLE PROBLEM, PROGRAM NO. 143A-63, 12/02/63

1.0 7+0.10000000E-06

0.0	0.0	3.0	0.0	0.0	0.0		1A
-0.435	-0.435	-0.435	0.0	0.0	1.523		1B
0.0	0.0	0.0	0.5	0.0	0.0	0 1	1C
0.	0.	3.5	0.0	0.0	0.0		2A
0.588	0.588	0.588	0.0	0.0	1.0		2B
0.0	0.0	0.0	0.5	0.0	0.0	0 1	2C
0.0	0.0	4.3525	-45.0	0.0	0.0		3A
0.0	0.0	0.0	0.0	0.0	1.6		3B
-0.5	0.5	-0.5	0.5	0.0	0.0	1 1	3C
0.0	0.0	4.9525	0.0	0.0	0.0		4A
0.0	0.0	0.0	0.0	0.0	1.0		4B
-0.5	0.5	-0.5	0.5	0.0	0.0	1 1	4C
0.	0.0	6.0	0.0	0.0	0.0		5A
0.0	-0.25	-1.0	0.0	0.0	1.5		5B
-2.0	2.0	-1.0	1.0	0.0	0.0	1 1	5C
0.0	0.0	6.75	0.0	0.0	0.0		6A
0.0	0.0	0.0	0.0	0.0	1.0		6B
-2.0	2.0	-1.0	1.0	0.0	0.0	1 1	6C
0.0	0.0	8.0	0.0	0.0	0.0		7A
0.0	0.0	0.0	0.0	0.0	1.0		7B
0.0	0.0	0.0	10.0	0.0	0.0	0 1	7C
0.1	0.0	0.1	0.3	3.0	0.3		
0.3	0.3						

SAMPLE PROBLEM DATA OUTPUT

SAMPLE PROBLEM, PROGRAM NO. 143A-63, 12/02/63

RAY SURF		OBJ PT	2D PT/DC	INT PT	DIR COS
1 1	X OR K	.10000	.30000	.30568	-.05569
	Y OR L	.00000	.30000	.30852	-.03509
	Z OR M	.00000	3.00000	3.08521	.99781
1 2	X OR K	.10000	.30000	.28848	-.32813
	Y OR L	.00000	.30000	.29772	-.30423
	Z OR M	.00000	3.00000	3.39209	.89429
1 3	X OR K	.10000	.30000	-.04675	-.49435
	Y OR L	.00000	.30000	-.01308	-.19014
	Z OR M	.00000	3.00000	4.30574	.84820
1 4	X OR K	.10000	.30000	-.42369	-.79097
	Y OR L	.00000	.30000	-.15807	-.30423
	Z OR M	.00000	3.00000	4.95250	.53085
1 5	NO INCIDENCE				
2 1	X OR K	.10000	.00000	-.00136	-.02133
	Y OR L	.00000	.30000	.30409	-.02971
	Z OR M	.00000	3.00000	3.04095	.99933
2 2	X OR K	.10000	.00000	-.01006	-.02552
	Y OR L	.00000	.30000	.29198	-.24762
	Z OR M	.00000	3.00000	3.44823	.96851
2 3	X OR K	.10000	.00000	-.03302	-.33870
	Y OR L	.00000	.30000	.06923	-.15476
	Z OR M	.00000	3.00000	4.31947	.92807
2 4	X OR K	.10000	.00000	-.26404	-.54192
	Y OR L	.00000	.30000	-.03633	-.24762
	Z OR M	.00000	3.00000	4.95250	.80311
2 5	X OR K	.10000	.00000	-.99486	-.36128
	Y OR L	.00000	.30000	-.37026	-.08675
	Z OR M	.00000	3.00000	6.03553	.92840
2 6	X OR K	.10000	.00000	-1.27289	-.54192
	Y OR L	.00000	.30000	-.43702	-.13013
	Z OR M	.00000	3.00000	6.75000	.83028
2 7	X OR K	.10000	.00000	-2.08877	-.54192
	Y OR L	.00000	.30000	-.63295	-.13013
	Z OR M	.00000	3.00000	8.00000	.83028
2	NO INT W/OPT AX				
3 1	X OR K	.10000	-.30000	-.31159	.01673
	Y OR L	.00000	.30000	.30869	-.03736
	Z OR M	.00000	3.00000	3.08697	.99916
3 2	X OR K	.10000	-.30000	-.30660	.28045
	Y OR L	.00000	.30000	.29755	-.30434
	Z OR M	.00000	3.00000	3.38486	.91034
3 3	X OR K	.10000	-.30000	-.01229	-.21390
	Y OR L	.00000	.30000	-.02183	-.19021
	Z OR M	.00000	3.00000	4.34020	.95815
3 4	X OR K	.10000	-.30000	-.14898	-.34225
	Y OR L	.00000	.30000	-.14338	-.30434
	Z OR M	.00000	3.00000	4.95250	.88895
3 5	X OR K	.10000	-.30000	-.58127	-.22816
	Y OR L	.00000	.30000	-.52780	-.08569
	Z OR M	.00000	3.00000	6.07531	.96984
3 6	X OR K	.10000	-.30000	-.74000	-.34225

		Y OR L	.00000	.30000	-.58741	-.12854
		X OR M	.00000	3.00000	6.75000	.93077
3	7	X OR K	.10000	-.30000	-1.19964	-.30225
		Y OR L	.00000	.30000	-.76005	-.12854
		Z OR M	.00000	3.00000	8.00000	.93077
3		NO INT W/OPT AX				
4	1	X OR K	.00000	.30000	.30860	-.03614
		Y OR L	.00000	.30000	.30860	-.03614
		Z OR M	.00000	3.00000	3.08608	.99869
4	2	X OR K	.00000	.30000	.29766	-.30412
		Y OR L	.00000	.30000	.29766	-.30412
		Z OR M	.00000	3.00000	3.38849	.90278
4	3	X OR K	.00000	.30000	-.02026	-.40105
		Y OR L	.00000	.30000	-.02026	-.19007
		Z OR M	.00000	3.00000	4.33223	.85561
4	4	X OR K	.00000	.30000	-.36928	-.77032
		Y OR L	.00000	.30000	-.15005	-.30412
		Z OR M	.00000	3.00000	4.95250	.56045
4	5	APERTURE STOP				
5	1	X OR K	.00000	.00000	.00000	.00000
		Y OR L	.00000	.30000	.30409	-.02964
		Z OR M	.00000	3.00000	3.04095	.99956
5	2	X OR K	.00000	.00000	.00000	.00000
		Y OR L	.00000	.30000	.29201	-.24751
		Z OR M	.00000	3.00000	3.44820	.96080
5	3	X OR K	.00000	.00000	.00000	-.32679
		Y OR L	.00000	.30000	.06101	-.15449
		Z OR M	.00000	3.00000	4.35250	.93234
5	4	X OR K	.00000	.00000	-.21030	-.52207
		Y OR L	.00000	.30000	-.03853	-.24751
		Z OR M	.00000	3.00000	4.95250	.81568
5	5	X OR K	.00000	.00000	-.90414	-.34050
		Y OR L	.00000	.30000	-.36698	-.08021
		Z OR M	.00000	3.00000	6.03488	.93311
5	6	X OR K	.00000	.00000	-1.17120	-.52207
		Y OR L	.00000	.30000	-.43459	-.13232
		Z OR M	.00000	3.00000	6.75000	.84207
5	7	X OR K	.00000	.00000	-1.94745	-.52207
		Y OR L	.00000	.30000	-.63101	-.13232
		Z OR M	.00000	3.00000	8.00000	.84207
5		NO INT W/OPT AX				
6	1	X OR K	.00000	-.30000	-.30860	.03614
		Y OR L	.00000	.30000	.30860	-.03614
		Z OR M	.00000	3.00000	3.08608	.99869
6	2	X OR K	.00000	-.30000	-.29766	.30412
		Y OR L	.00000	.30000	.29766	-.30412
		Z OR M	.00000	3.00000	3.38849	.90278
6	3	X OR K	.00000	-.30000	.04085	-.20566
		Y OR L	.00000	.30000	-.04085	-.19007
		Z OR M	.00000	3.00000	4.39335	.95998
6	4	X OR K	.00000	-.30000	-.07093	-.32906
		Y OR L	.00000	.30000	-.15156	-.30412
		Z OR M	.00000	3.00000	4.95250	.89399
6	5	X OR K	.00000	-.30000	-.49297	-.21937
		Y OR L	.00000	.30000	-.53421	-.08393

6	6	Z OR M	.00000	3.00000	6.07732	.97202
		X OR K	.00000	-.30000	-.64479	-.32906
		Y OR L	.00000	.30000	-.59230	-.12590
		Z OR M	.00000	3.00000	6.75000	.93587
6	7	X OR K	.00000	-.30000	-1.08431	-.32906
		Y OR L	.00000	.30000	-.76046	-.12590
		Z OR M	.00000	3.00000	8.00000	.93587
6		NO INT W/OPT AX				
7	1	X OR K	-.10000	.30000	.31159	-.01673
		Y OR L	.00000	.30000	.30869	-.03736
		Z OR M	.00000	3.00000	3.08697	.99916
7	2	X OR K	-.10000	.30000	.30660	-.28045
		Y OR L	.00000	.30000	.29755	-.30434
		Z OR M	.00000	3.00000	3.38486	.91034
7	3	X OR K	-.10000	.30000	.00650	-.46886
		Y OR L	.00000	.30000	-.02811	-.19021
		Z OR M	.00000	3.00000	4.35900	.86254
7	4	X OR K	-.10000	.30000	-.31611	-.75010
		Y OR L	.00000	.30000	-.15900	-.30434
		Z OR M	.00000	3.00000	4.95250	.58702
7	5	X OR K	-.10000	.30000	-1.91805	-.50012
		Y OR L	.00000	.30000	-.68090	.11002
		Z OR M	.00000	3.00000	6.20602	.85707
7	6	APERTURE STOP				
8	1	X OR K	-.10000	.00000	.00136	.02133
		Y OR L	.00000	.30000	.30409	-.02971
		Z OR M	.00000	3.00000	3.04095	.99933
8	2	X OR K	-.10000	.00000	.01006	.02552
		Y OR L	.00000	.30000	.29198	-.24762
		Z OR M	.00000	3.00000	3.44023	.96051
8	3	X OR K	-.10000	.00000	.03481	-.31508
		Y OR L	.00000	.30000	.05180	-.15476
		Z OR M	.00000	3.00000	4.38731	.93636
8	4	X OR K	-.10000	.00000	-.15537	-.50413
		Y OR L	.00000	.30000	-.04152	-.24762
		Z OR M	.00000	3.00000	4.95250	.82736
8	5	X OR K	-.10000	.00000	-.81469	-.33600
		Y OR L	.00000	.30000	-.36538	-.08932
		Z OR M	.00000	3.00000	6.03457	.93758
8	6	X OR K	-.10000	.00000	-1.07115	-.50413
		Y OR L	.00000	.30000	-.43354	-.13300
		Z OR M	.00000	3.00000	6.75000	.85317
8	7	X OR K	-.10000	.00000	-1.00976	-.50413
		Y OR L	.00000	.30000	-.62984	-.13990
		Z OR M	.00000	3.00000	8.00000	.85317
8		NO INT W/OPT AX				
9	1	X OR K	-.10000	-.30000	-.30568	.05509
		Y OR L	.00000	.30000	.30852	-.03509
		Z OR M	.00000	3.00000	3.08521	.99781
9	2	X OR K	-.10000	-.30000	-.28048	.32013
		Y OR L	.00000	.30000	.29772	-.30423
		Z OR M	.00000	3.00000	3.39209	.89429
9	3	X OR K	-.10000	-.30000	.10094	-.19763
		Y OR L	.00000	.30000	-.06333	-.19014
		Z OR M	.00000	3.00000	4.45344	.96165



9	4	X OR K	-.10000	-.30000	-.00162	-.31621
		Y OR L	.00000	.30000	-.16200	-.30423
		Z OR M	.00000	3.00000	4.95250	.89858
9	5	X OR K	-.10000	-.30000	-.39854	-.21081
		Y OR L	.00000	.30000	-.54388	-.08126
		Z OR M	.00000	3.00000	6.08042	.97414
9	6	X OR K	-.10000	-.30000	-.54344	-.31621
		Y OR L	.00000	.30000	-.59974	-.12190
		Z OR M	.00000	3.00000	6.75000	.94082
9	7	X OR K	-.10000	-.30000	-.96357	-.31621
		Y OR L	.00000	.30000	-.76170	-.12190
		Z OR M	.00000	3.00000	8.00000	.94082
9		NO INT W/OPT AX				
10	1	X OR K	.10000	.30000	.30270	-.05012
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.04057	.99874
10	2	X OR K	.10000	.30000	.20206	-.27290
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.45185	.96204
10	3	X OR K	.10000	.30000	.02070	-.45816
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.37920	.88886
10	4	X OR K	.10000	.30000	-.27789	-.73305
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.95250	.68016
10	5	X OR K	.10000	.30000	-1.40684	-.48870
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.00000	.87244
10	6	X OR K	.10000	.30000	-1.82695	-.73305
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.75000	.68016
10	7	X OR K	.10000	.30000	-3.17415	-.73305
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	8.00000	.68016
10		OPT AX INTERSECTION Z = 5.05485				
11	1	X OR K	.10000	.00000	.00000	-.02107
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.00000	.99976
11	2	X OR K	.10000	.00000	-.01093	-.02658
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.49992	.99964
11	3	X OR K	.10000	.00000	-.03274	-.33430
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.31975	.94246
11	4	X OR K	.10000	.00000	-.25718	-.53488
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.95250	.84492
11	5	X OR K	.10000	.00000	-.92031	-.35659
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.00000	.93426
11	6	X OR K	.10000	.00000	-1.20657	-.53488
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.75000	.84492
11	7	X OR K	.10000	.00000	-1.99789	-.53488
		Y OR L	.00000	.00000	.00000	.00000

11	OPT AX INTERSECTION 2 =	4.84495	3.00000	8.00000	84492
12 1	X OR K	*10000	--.10000	-1.30531	*00947
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	3.04134	*99995
12 2	X OR K	*10000	--.10000	-1.30160	*22237
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	3.04440	*97496
12 3	X OR K	*10000	--.10000	-1.12259	--.22583
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	4.02290	*97416
12 4	X OR K	*10000	--.10000	-1.29010	-1.36135
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	4.09250	*93243
12 5	X OR K	*10000	--.10000	-1.69602	-1.24000
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	6.10000	*97059
12 6	X OR K	*10000	--.10000	-1.80216	-1.436153
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	6.175000	*99845
12 7	X OR K	*10000	--.10000	-1.36655	-1.36193
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	8.00000	*93243
12	OPT AX INTERSECTION 2 =	4.47350	3.00000	8.00000	93243
13 1	X OR K	*00000	*00000	3.00000	-1.02964
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	3.04095	*99956
13 2	X OR K	*00000	*00000	3.00000	-1.24751
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	3.44828	*98900
13 3	X OR K	*00000	*00000	3.00000	-1.44470
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	4.04110	*09563
13 4	X OR K	*00000	*00000	3.00000	-1.71165
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	4.95250	*70252
13 5	X OR K	*00000	*00000	3.00000	-1.47443
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	6.00000	*00028
13 6	X OR K	*00000	*00000	3.00000	-1.71165
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	6.75000	*70252
13 7	X OR K	*00000	*00000	3.00000	-1.71165
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	8.00000	*70252
13	OPT AX INTERSECTION 2 =	5.08112	3.00000	8.00000	70252
14 1	X OR K	*00000	*00000	3.00000	*00000
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	3.00000	*00000
14 2	X OR K	*00000	*00000	3.00000	*00000
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	3.50000	*00000
14 3	X OR K	*00000	*00000	3.00000	-1.32100
	Y OR L	*00000	*00000	1.00000	*00000
	Z OR M	*00000	3.00000	3.00000	*00000

14	4	Z OR M	.00000	3.00000	4.35250	.94688
		X OR K	.00000	.00000	-.20393	-.51488
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.95250	.85725
14	5	X OR K	.00000	.00000	-.83908	-.34325
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.00000	.93924
14	6	X OR K	.00000	.00000	-1.10718	-.51488
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.75000	.85725
14	7	X OR K	.00000	.00000	-1.85796	-.51488
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	8.00000	.85725
14		OPT AX INTERSECTION Z =	4.90661			
15	1	X OR K	.00000	-.30000	-.30409	.02964
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.04095	.99956
15	2	X OR K	.00000	-.30000	-.29201	.24751
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.44828	.96888
15	3	X OR K	.00000	-.30000	-.08195	-.21611
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.27054	.97636
15	4	X OR K	.00000	-.30000	-.23290	-.34578
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.95250	.93831
15	5	X OR K	.00000	-.30000	-.61892	-.23052
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.00000	.97506
15	6	X OR K	.00000	-.30000	-.79660	-.34578
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.75000	.93831
15	7	X OR K	.00000	-.30000	-1.25725	-.34578
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	8.00000	.93831
15		OPT AX INTERSECTION Z =	4.58836			
16	1	X OR K	-.10000	.30000	.30551	-.00947
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.04134	.99995
16	2	X OR K	-.10000	.30000	.30168	-.22237
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.44660	.97496
16	3	X OR K	-.10000	.30000	.07705	-.43166
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.42955	.90203
16	4	X OR K	-.10000	.30000	-.17919	-.69066
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.95250	.72317
16	5	X OR K	-.10000	.30000	-1.17361	-.46044
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.00000	.88768
16	6	X OR K	-.10000	.30000	-1.56263	-.69066
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.75000	.72317
16	7	X OR K	-.10000	.30000	-2.75644	-.69066

		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	0.00000	.72317
16		OPT AX INTERSECTION Z =	5.11381			
17	1	X OR K	-.10000	.00000	.00000	.02187
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.00000	.99976
17	2	X OR K	-.10000	.00000	.01093	.02658
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.49992	.99964
17	3	X OR K	-.10000	.00000	.03453	-.30950
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.38703	.95009
17	4	X OR K	-.10000	.00000	-.14951	-.49520
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.95250	.86877
17	5	X OR K	-.10000	.00000	-.74659	-.33013
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.00000	.94393
17	6	X OR K	-.10000	.00000	-1.00000	-.49520
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.75000	.86877
17	7	X OR K	-.10000	.00000	-1.72140	-.49520
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	8.00000	.86877
17		OPT AX INTERSECTION Z =	4.97999			
18	1	X OR K	-.10000	-.30000	-.30270	.05012
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.04057	.99874
18	2	X OR K	-.10000	-.30000	-.28206	.27290
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	3.45105	.96204
18	3	X OR K	-.10000	-.30000	-.03709	-.20658
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.31540	.97842
18	4	X OR K	-.10000	-.30000	-.17161	-.33055
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	4.95250	.94379
18	5	X OR K	-.10000	-.30000	-.53847	-.22035
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.00000	.97541
18	6	X OR K	-.10000	-.30000	-.70791	-.33055
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	6.75000	.94379
18	7	X OR K	-.10000	-.30000	-1.14569	-.33055
		Y OR L	.00000	.00000	.00000	.00000
		Z OR M	.00000	3.00000	8.00000	.94379
18		OPT AX INTERSECTION Z =	4.72060			
19	1	X OR K	.10000	.30000	.30560	-.05589
		Y OR L	.00000	-.30000	-.30852	.03509
		Z OR M	.00000	3.00000	3.08821	.99781
19	2	X OR K	.10000	.30000	.20048	-.32013
		Y OR L	.00000	-.30000	-.29772	.30429
		Z OR M	.00000	3.00000	3.39209	.89429
19	3	X OR K	.10000	.30000	-.04675	-.49435

		Y OR L	.00000	-.30000	.01308	.19014	
		Z OR M	.00000	3.00000	4.30574	.84820	
19	4	X OR K	.10000	.30000	-.42369	-.79097	
		Y OR L	.00000	-.30000	.15807	.30423	
		Z OR M	.00000	3.00000	4.95250	.53085	
19	5	NO INCIDENCE					
20	1	X OR K	.10000	.00000	-.00136	-.02133	
		Y OR L	.00000	-.30000	-.30409	.02971	
		Z OR M	.00000	3.00000	3.04095	.99933	
20	2	X OR K	.10000	.00000	-.01006	-.02552	
		Y OR L	.00000	-.30000	-.29198	.24762	
		Z OR M	.00000	3.00000	3.44823	.96851	
20	3	X OR K	.10000	.00000	-.03302	-.33870	
		Y OR L	.00000	-.30000	-.06923	.15476	
		Z OR M	.00000	3.00000	4.31947	.92807	
20	4	X OR K	.10000	.00000	-.26404	-.54192	
		Y OR L	.00000	-.30000	.03633	.24762	
		Z OR M	.00000	3.00000	4.95250	.80311	
20	5	X OR K	.10000	.00000	-.99486	-.36128	
		Y OR L	.00000	-.30000	.37026	.08675	
		Z OR M	.00000	3.00000	6.03553	.92840	
20	6	X OR K	.10000	.00000	-1.27289	-.54192	
		Y OR L	.00000	-.30000	.43702	.13013	
		Z OR M	.00000	3.00000	6.75000	.83028	
20	7	X OR K	.10000	.00000	-2.08877	-.54192	
		Y OR L	.00000	-.30000	.63295	.13013	
		Z OR M	.00000	3.00000	8.00000	.83028	
20	NO INT W/OPT AX						
21	1	X OR K	.10000	-.30000	-.31159	.01673	
		Y OR L	.00000	-.30000	-.30869	.03736	
		Z OR M	.00000	3.00000	3.08697	.99916	
21	2	X OR K	.10000	-.30000	-.30660	.28045	
		Y OR L	.00000	-.30000	-.29755	.30434	
		Z OR M	.00000	3.00000	3.38486	.91034	
21	3	X OR K	.10000	-.30000	-.01229	-.21390	
		Y OR L	.00000	-.30000	.02103	.19021	
		Z OR M	.00000	3.00000	4.34020	.95815	
21	4	X OR K	.10000	-.30000	-.14898	-.34225	
		Y OR L	.00000	-.30000	.14338	.30434	
		Z OR M	.00000	3.00000	4.95250	.88895	
21	5	X OR K	.10000	-.30000	-.50127	-.22816	
		Y OR L	.00000	-.30000	.52780	.08569	
		Z OR M	.00000	3.00000	6.07531	.96984	
21	6	X OR K	.10000	-.30000	-.74000	-.34225	
		Y OR L	.00000	-.30000	.58741	.12054	
		Z OR M	.00000	3.00000	6.75000	.93077	
21	7	X OR K	.10000	-.30000	-1.19964	-.34225	
		Y OR L	.00000	-.30000	.76005	.12054	
		Z OR M	.00000	3.00000	8.00000	.93077	
21	NO INT W/OPT AX						
22	1	X OR K	.00000	.30000	.30860	-.03614	
		Y OR L	.00000	-.30000	-.30860	.03614	
		Z OR M	.00000	3.00000	3.08608	.99869	
22	2	X OR K	.00000	.30000	.29766	-.30412	
		Y OR L	.00000	-.30000	-.29766	.30412	

22	3	Z OR M	.00000	3.00000	3.38849	.90278	
		X OR K	.00000	-.30000	-.02026	-.48143	
		Y OR L	.00000	-.30000	.02026	.19007	
22	4	Z OR M	.00000	3.00000	4.33223	.85561	
		X OR K	.00000	-.30000	-.36928	-.77032	
		Y OR L	.00000	-.30000	.15805	.30412	
		Z OR M	.00000	3.00000	4.95250	.56045	
22	5	APERTURE STOP					
23	1	X OR K	.00000	.00000	.00000	.00000	
		Y OR L	.00000	-.30000	-.30409	.02964	
		Z OR M	.00000	3.00000	3.04095	.99956	
23	2	X OR K	.00000	.00000	.00000	.00000	
		Y OR L	.00000	-.30000	-.29201	.24751	
		Z OR M	.00000	3.00000	3.44828	.96888	
23	3	X OR K	.00000	.00000	.00000	-.32679	
		Y OR L	.00000	-.30000	-.06101	.15469	
		Z OR M	.00000	3.00000	4.35250	.93234	
23	4	X OR K	.00000	.00000	-.21030	-.52287	
		Y OR L	.00000	-.30000	.03853	.24751	
		Z OR M	.00000	3.00000	4.95250	.81560	
23	5	X OR K	.00000	.00000	-.90414	-.34858	
		Y OR L	.00000	-.30000	.36698	.08821	
		Z OR M	.00000	3.00000	6.03488	.93311	
23	6	X OR K	.00000	.00000	-1.17120	-.52287	
		Y OR L	.00000	-.30000	.43459	.19232	
		Z OR M	.00000	3.00000	6.75000	.84207	
23	7	X OR K	.00000	.00000	-1.94745	-.52287	
		Y OR L	.00000	-.30000	.63101	.13232	
		Z OR M	.00000	3.00000	8.00000	.84207	
23	NO INT W/OPT AX						
24	1	X OR K	.00000	-.30000	-.30860	.03614	
		Y OR L	.00000	-.30000	-.30860	.03614	
		Z OR M	.00000	3.00000	3.08600	.99869	
24	2	X OR K	.00000	-.30000	-.29766	.30412	
		Y OR L	.00000	-.30000	-.29766	.30412	
		Z OR M	.00000	3.00000	3.38849	.90278	
24	3	X OR K	.00000	-.30000	.04085	-.20566	
		Y OR L	.00000	-.30000	.04085	.19007	
		Z OR M	.00000	3.00000	4.39335	.95998	
24	4	X OR K	.00000	-.30000	-.07893	-.32906	
		Y OR L	.00000	-.30000	.15156	.30412	
		Z OR M	.00000	3.00000	4.95250	.89399	
24	5	X OR K	.00000	-.30000	-.49297	-.21937	
		Y OR L	.00000	-.30000	.53421	.08393	
		Z OR M	.00000	3.00000	6.07732	.97202	
24	6	X OR K	.00000	-.30000	-.64479	-.32906	
		Y OR L	.00000	-.30000	.59230	.12590	
		Z OR M	.00000	3.00000	6.75000	.93587	
24	7	X OR K	.00000	-.30000	-1.08431	-.32906	
		Y OR L	.00000	-.30000	.76046	.12590	
		Z OR M	.00000	3.00000	8.00000	.93587	
24	NO INT W/OPT AX						
25	1	X OR K	.10000	.30000	.31159	-.01673	
		Y OR L	.00000	-.30000	-.30869	.03736	
		Z OR M	.00000	3.00000	3.08697	.99916	

25	2	X OR K	-.10000	.30000	430660	-.28045	
		Y OR L	.00000	-.30000	-.29755	.30434	
		Z OR M	.00000	3.00000	3.38486	.91034	
25	3	X OR K	-.10000	.30000	.00650	-.46886	
		Y OR L	.00000	-.30000	.02811	.19021	
		Z OR M	.00000	3.00000	4.35900	.86294	
25	4	X OR K	-.10000	.30000	-.43161	-.75018	
		Y OR L	.00000	-.30000	.15900	.30434	
		Z OR M	.00000	3.00000	4.95250	.58702	
25	5	X OR K	-.10000	.30000	-1.91805	-.50012	
		Y OR L	.00000	-.30000	.80890	-.11802	
		Z OR M	.00000	3.00000	6.20602	.85787	
25	6	APERTURE STOP					
26	1	X OR K	-.10000	.00000	.00136	.02133	
		Y OR L	.00000	-.30000	-.30409	.02971	
		Z OR M	.00000	3.00000	3.04095	.99933	
26	2	X OR K	-.10000	.00000	.01006	.02552	
		Y OR L	.00000	-.30000	-.29198	.24762	
		Z OR M	.00000	3.00000	3.44823	.96851	
26	3	X OR K	-.10000	.00000	.03401	-.31500	
		Y OR L	.00000	-.30000	-.05188	.15476	
		Z OR M	.00000	3.00000	4.38731	.99636	
26	4	X OR K	-.10000	.00000	-.15537	-.50413	
		Y OR L	.00000	-.30000	.04152	.24762	
		Z OR M	.00000	3.00000	4.95250	.82736	
26	5	X OR K	-.10000	.00000	-.01469	-.33600	
		Y OR L	.00000	-.30000	.36538	.08932	
		Z OR M	.00000	3.00000	6.03457	.93758	
26	6	X OR K	-.10000	.00000	-1.07115	-.50413	
		Y OR L	.00000	-.30000	.43354	.19398	
		Z OR M	.00000	3.00000	6.75000	.85317	
26	7	X OR K	-.10000	.00000	-1.00976	-.50413	
		Y OR L	.00000	-.30000	.62984	.13398	
		Z OR M	.00000	3.00000	8.00000	.85317	
26		NO INT W/OPT AX					
27	1	X OR K	-.10000	-.30000	-.30560	.05589	
		Y OR L	.00000	-.30000	-.30852	.03509	
		Z OR M	.00000	3.00000	3.08521	.99781	
27	2	X OR K	-.10000	-.30000	-.28048	.32013	
		Y OR L	.00000	-.30000	-.29772	.30423	
		Z OR M	.00000	3.00000	3.39209	.89429	
27	3	X OR K	-.10000	-.30000	.10094	-.19763	
		Y OR L	.00000	-.30000	.06333	.19014	
		Z OR M	.00000	3.00000	4.45344	.96165	
27	4	X OR K	-.10000	-.30000	-.00162	-.31621	
		Y OR L	.00000	-.30000	.16200	.30423	
		Z OR M	.00000	3.00000	4.95250	.89858	
27	5	X OR K	-.10000	-.30000	-.39054	-.21081	
		Y OR L	.00000	-.30000	.54388	.08126	
		Z OR M	.00000	3.00000	6.08042	.97414	
27	6	X OR K	-.10000	-.30000	-.54344	-.31621	
		Y OR L	.00000	-.30000	.59974	.12190	
		Z OR M	.00000	3.00000	6.75000	.94082	
27	7	X OR K	-.10000	-.30000	-.96357	-.31621	
		Y OR L	.00000	-.30000	.76170	.12190	
		Z OR M	.00000	3.00000	8.00000	.94082	

27 NO INT W/OPT AX



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C      GENERAL RAY TRACE PROGRAM                      CARD I/O
      DIMENSION XO(35),YO(35),ZO(35),ALPHA(35),BETA(35),GAMMA(35),A(35)
      DIMENSION REFR(36),AP1(35),AP2(35),AP3(35),AP4(35),AP5(35),AP6(35)
      DIMENSION NAP(35),NOUT(35),CKBAR(35),CLBAR(35),CMBAR(35),X0BAR(35)
      DIMENSION Y0BAR(35),Z0BAR(35),B(35),C(35),D(35),E(35)
      DIMENSION R(35),KSUR(35)

      KKK = 1
      GO TO (813, 803), KKK

      803 PRINT 115
      115 FORMAT (20HLOAD DATA,PUSH START)
      PAUSE
      813 KKK = 2
C      READ TITLE CARD
      READ 2
2FORMAT(49H                                     ,8H      )
      PUNCH 2
      PUNCH 100
      JB = 50
      PUNCH 1
      PUNCH 100
      PUNCH 100
      PI = 3.14159265
      CONV = PI/180.
C      READ HEADER CARD
      102 FORMAT (F11.5,I3,E15.8)
      READ 102, REFR(1),NOSUR,TOL
C      REFR = INITIAL INDEX OF REFR , NOSUR = NO. OF SURFACES
C      TOL = MIN. LIMIT ON ITERATION INCREMENT
      NOSUR = NOSUR+1

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DO 203 I = 2, NOSUR
C   READ SURFACE DATA
103 FORMAT (F11.5, F11.5, F11.5, F11.5, F11.5, F11.5, I2, I2, I3)
    READ 103, XO(I), YO(I), ZO(I), ALPHA(I), BETA(I), GAMMA(I)
C   XO, YO, AND ZO ARE LOCAL SYSTEM SURFACE VERTEX COORDINATES.
C   ALPHA, BETA, AND GAMMA ARE EULER ROTATIONAL ANGLES.
    READ 103, A(I), B(I), C(I), D(I), E(I), REFR(I)
C   A, B, C, D, AND E ARE SURFACE EQUATION COEFFICIENTS.
C   REFR = INDEX OF REFRACTION FOLLOWING SURFACE.
203 READ 103, AP1(I), AP2(I), AP3(I), AP4(I), AP5(I), AP6(I), NAP(I), NOUT(I)
C   AP1, AP2, AP3, AP4, AP5, AND AP6 SPECIFY APERTURE DIMENSIONS. SEE
C   INPUT WRITE UP.
C   NAP = 0 FOR CIRCULAR-ANNULAR APERTURE, = 1 FOR RECT-TRAPEZOIDAL
C   APERTURE, = 2 FOR HYPERBOLIC APERTURE.
C   NOUT = 0 FOR NO OUTPUT AT SURFACE, = 1 OTHERWISE.
C   READ RAY DATA
C   SWITCH 1 ON TO USE SINGLE RAY CARD INPUT
    IF (SENSE SWITCH 1) 205, 201
205 READ 103, XA, YA, ZA, RPAR1, RPAR2, RPAR3, NIN, NAXIN, IRAY
C   XA, YA, AND ZA ARE RAY COORDINATES AT FIRST POINT
C   RPAR1, RPAR2, RPAR3 ARE RAY COORDINATES AT SECOND POINT
C   OR RAY DIRECTION COSINES AT FIRST POINT
C   NIN = 0 IF RAYS ARE SPEC. BY 2 POINTS
C   NIN = 1 IF RAYS ARE SPEC. BY 1 POINT AND DIR. COSINES.
C   NAXIN = 1 FOR COMPUTATION OF INTERSECTION OF RAY WITH OPTICAL
C   AXIS, = 0 OTHERWISE IRAY = RAY NUMBER.
    GO TO 206
201 READ 116, FXU1, FZ1, XGAP1, FXU2, FZ2, XGAP2
    READ 116, FYU2, YGAP2

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C   RAY INPUT DATA, FXU1 = MAX.XR, FZ1 = ZR, FXU2 = MAX. X AT POINT 2
C   XGAP1 = X SPACING AT POINT 1, XGAP2 = X SPACING AT POINT 2
C   FZ2 = Z COORDINATE AT POINT 2
C   FYU2 = MAX Y COORDINATE AT 2ND POINT, YGAP2 = Y SPACING AT POINT 2
116 FORMAT(F11.5,F11.5,F11.5,F11.5,F11.5,F11.5)
    IRAY = 0
    GAP3=0.
    RPAR2=FYU2
812 RPAR2=RPAR2-GAP3
    IF(RPAR2+FYU2)803,814,814
814 GAP3=YGAP2
    GAP1 = 0.
    XA = FXU1
    4 XA = XA-GAP1
    IF(XA+FXU1)812,804,804
804 GAP1 = XGAP1
    GAP2 = 0.
    RPAR1 = FXU2
    3 RPAR1 = RPAR1-GAP2
    IF(RPAR1+FXU2)4,805,805
805 GAP2 = XGAP2
    YA = 0.0
    ZA = FZ1
    RPAR3 = FZ2
    NIN = 0
    NAXIN = 1
    IRAY = IRAY+1
206 NIN = NIN+1
    GO TO (207,208), NIN
```

```
208 CKBAR(1) = RPAR1
    CLBAR(1) = RPAR2
    CMBAR(1) = RPAR3
    GO TO 209

207 XD = RPAR1-XA
    YD = RPAR2-YA
    ZD = RPAR3-ZA
    RALEN = SQRT(XD*XD+YD*YD+ZD*ZD)
    CKBAR(1) = XD/RALEN
    CLBAR(1) = YD/RALEN
    CMBAR(1) = ZD/RALEN

209 X0BAR(1) = XA
    Y0BAR(1) = YA
    Z0BAR(1) = ZA

DO 15 I = 2, NOSUR
    IM1 = I-1
    SA = SIN(CONV*ALPHA(I))
    CA = COS(CONV*ALPHA(I))
    SB = SIN(CONV*BETA(I))
    CB = COS(CONV*BETA(I))
    SG = SIN(CONV*GAMMA(I))
    CG = COS(CONV*GAMMA(I))
    X1 = X0BAR(IM1)-XO(I)
    Y1 = Y0BAR(IM1)-YO(I)
    Z1 = Z0BAR(IM1)-ZO(I)
    R11 = CA*CG+SA*SB*SG
    R12 = -CB*SG
    R13 = CA*SB*SG-SA*CG
    R21 = CA*SG-SA*SB*CG
```

```

R22 = CB*CG
R23 = -(SA*SG+CA*SB*CG)
R31 = SA*CB
R32 = SB
R33 = CA*CB
X0 = X1*R11+Y1*R12+Z1*R13
Y0 = X1*R21+Y1*R22+Z1*R23
Z0 = X1*R31+Y1*R32+Z1*R33
CK = CKBAR(IM1)*R11+CLBAR(IM1)*R12+CMBAR(IM1)*R13
CL = CKBAR(IM1)*R21+CLBAR(IM1)*R22+CMBAR(IM1)*R23
CM = CKBAR(IM1)*R31+CLBAR(IM1)*R32+CMBAR(IM1)*R33
J = 0
IF(D(I)) 310,311,302
310 Z2 = -Z0
    IF(CM) 600,10,600
600 X0 = X0-2.*CK*Z0/CM
    Y0 = Y0-2.*CL*Z0/CM
    GO TO 211
311 IF(CM) 210,302,210
210 X0 = X0-CK*Z0/CM
    Y0 = Y0-CL*Z0/CM
    Z2 = 0.
    GO TO 211
302 Z2 = Z0
211 S = 0.
    5 J = J+1
    X = X0+CK*S
    Y = Y0+CL*S
    Z = Z2+CM*S

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```

      IF(D(I)) 212,213,214
212 IF (Z) 7,6,6
      7 Z = -Z
      IF (CM) 400 10, 400
400 X = X0+CK*(Z-Z2)/CM
      Y = Y0+CL*(Z-Z2)/CM
      6 IF(Y) 215,216,216
215 Y = -Y
      IF(CL) 401, 10, 401
401 X = X0+CK*(Y-Y0)/CL
      Z = Z2+CM*(Y-Y0)/CL
216 DD = D(I)*D(I)
      F = E(I)*DD*X*X+DD*Y*Y-Z*Z
      FX = 2.*X*E(I)*DD
      FY = 2.*Y*DD
      FZ = -2.*Z
      GO TO 8
214 IF(Z) 6,6,7
213 F = A(I)*X*X+B(I)*Y*Y+C(I)*Z*Z+Z
      FX = 2.*A(I)*X
      FY = 2.*B(I)*Y
      FZ = 2.*C(I)*Z+1.
      8 DETMT = CK*FX+CL*FY+CM*FZ
      IF(DETMT) 218,217,218
217 IF(F) 10,9,10
      10 PUNCH 104, IRAY, IM1
      11 CONTINUE
      IF (SENSE SWITCH 1)205,3
218 DELS = -F/DETMT

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```
DELS2 = DELS*DELS
TOL2 = TOL*TOL
IF(DELS2-TOL2) 9, 9, 219
219 IF(J-JB) 220,10,10
220 S = S+DELS
GO TO 5
9 IF(D(I)) 320,321,301
320 IF(S-2.*Z0/CM) 221, 222, 222
321 IF(CM) 300,301,300
301 IF(S) 221,222,222
300 IF(S-Z0/CM) 221,222,222
221 PUNCH 105, IRAY, IM1
GO TO 11
222 KAP = NAP(I)+1
GO TO (223,224,225),KAP
223 RHSQ = (X-AP1(I))*(X-AP1(I))+(Y-AP2(I))*(Y-AP2(I))
IF(RHSQ-AP3(I)*AP3(I)) 12, 226, 226
12 PUNCH 106, IRAY, IM1
GO TO 11
226 IF(RHSQ-AP4(I)*AP4(I)) 13, 12, 12
224 IF(Y-AP3(I)) 12,12,227
227 IF(Y-AP4(I)) 228,12,12
228 IF(X-(AP5(I)*(Y-AP3(I))+AP1(I))) 12,12,229
229 IF(X-(AP6(I)*(Y-AP3(I))+AP2(I))) 13,12,12
225 IF(Y-AP5(I)) 12,12,230
230 IF(Y-AP6(I)) 231,12,12
231 FFF = (X-AP1(I))/AP3(I)
FFF2 = FFF*FFF
GGG = (Y-AP2(I))/AP4(I)
```

```

GGG2 = GGG*GGG
IF(FFF2-GGG2-1.) 13, 12, 12
13 RAT = RFFR(IM1)/RFFR(I)
   ALC = (RAT*DETMT)/(FX*FX+FY*FY+FZ*FZ)
   IF (REFR(I)+REFR(IM1)) 232,233,232
233 GAMUC = 2.*ALC
   RAT = 1.
   GO TO 14
232 BLC = (RAT*RAT-1.)/(FX*FX+FY*FY+FZ*FZ)
   DISC = ALC*ALC-BLC
   IF(DISC) 234,235,235
234 PUNCH 107, IRAY, IM1
   GO TO 11
235 DISC = SQRT(DISC)
   IF(ALC) 236,10,237
236 DISC = -DISC
237 GAMUC = DISC-ALC
14 CK = RAT*CK+GAMUC*FX
   CL = RAT*CL+GAMUC*FY
   CM = RAT*CM+GAMUC*FZ
   X0BAR(I) = R11*X+R21*Y+R31*Z+XO(I)
   Y0BAR(I) = R12*X+R22*Y+R32*Z+YO(I)
   Z0BAR(I) = R13*X+R23*Y+R33*Z+ZO(I)
   CKBAR(I) = R11*CK+R21*CL+R31*CM
   CLBAR(I) = R12*CK+R22*CL+R32*CM
   CMBAR(I) = R13*CK+R23*CL+R33*CM
   KOUT = NOUT(I)+1
   GO TO (15,16),KOUT
16 PUNCH 108, IRAY, IM1, XA, RPAR1, X0BAR(I), CKBAR(I)

```



```

PUNCH 109, YA, RPAR2, Y0BAR(I), CLBAR(I)
PUNCH 110, ZA, RPAR3, Z0BAR(I), CMBAR(I)
15 CONTINUE
NAXIN = NAXIN+1
GO TO (17,239), NAXIN
239 IF(CKBAR(NOSUR)) 240,241,240
241 IF(X0BAR(NOSUR)) 18,242,18
18 PUNCH 111, IRAY
GO TO 17
242 S = -Y0BAR(NOSUR)
IF(CLBAR(NOSUR)) 243,244,243
244 IF(Y0BAR(NOSUR)) 18,245,18
243 S = S/CLBAR(NOSUR)
GO TO 245
240 S = -X0BAR(NOSUR)/CKBAR(NOSUR)
IF(CLBAR(NOSUR)) 246,244,246
246 IF(S+Y0BAR(NOSUR)/CLBAR(NOSUR)) 18,245,18
245 AXIN = Z0BAR(NOSUR)+S*CMBAR(NOSUR)
PUNCH 112, IRAY, AXIN
17 PUNCH 100
IF(SENSE SWITCH 1) 205, 3
1 FORMAT(8HRAY SURF,27X6HOBJ PT,4X8H2D PT/DC,5X6HINT PT,6X7HDIR COS)
100 FORMAT (1X)
104 FORMAT (I3,I3,13H NO INCIDENCE)
105 FORMAT (I3,I3,13H VIRTUAL PATH)
106 FORMAT (I3,I3,14H APERTURE STOP)
107 FORMAT (I3,I3,14H NO REFRACTION)
108 FORMAT (I3,I3,16X,8HX OR K ,F11.5,1X,F11.5,1X,F11.5,1X,F11.5)
109 FORMAT (22X,8HY OR L ,F11.5,1X,F11.5,1X,F11.5,1X,F11.5)

```

110 FORMAT (22X,8HZ OR M ,F11.5,1X,F11.5,1X,F11.5,1X,F11.5)

111 FORMAT (I3,4X,15HNO INT W/OPT AX)

112 FORMAT (I3,4X,23HOPT AX INTERSECTION Z =,F11.5)

END

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